

ECE608 APPENDIX 1 PROBLEMS

- 1) A.1-1
- 2) A.1-6
- 3) A.1-7
- 4) A.2-4
- 5) A.2-5
- 6) A-1

(1) CLR A.1-1

$$\sum_{k=1}^n (2k-1) = 2 \sum_{k=1}^n k - \sum_{k=1}^n 1 = \frac{2n(n+1)}{2} - n = n(n+1) - n = n^2$$

(2) CLR A.1-6

Prove that $\sum_{k=1}^n O(f_k(n)) = O(\sum_{k=1}^n f_k(n))$.

Proof

Let $A = \sum_{k=1}^n O(f_k(n))$, and $B = O(\sum_{k=1}^n f_k(n))$.

Because A and B are sets of functions, we have to show $A \subseteq B$ and $B \subseteq A$.

(1) $A \subseteq B$

Let $g_k(n) = O(f_k(n))$ for $k \in \{1, 2, \dots, n\}$. Then, there exist $c_k > 0$ and $n_k > 0$ such that

$$0 \leq g_k(n) \leq c_k f_k(n) \quad \forall n \geq n_k$$

Choose $c = \max_{k \in \{1, 2, \dots, n\}} \{c_k\}$ and $n_0 = \max_{k \in \{1, 2, \dots, n\}} \{n_k\}$. Then, we have

$$0 \leq g_k(n) \leq c_k f_k(n) \leq c f_k(n) \quad \forall n \geq n_0$$

By the linear property of the summation,

$$0 \leq \sum_{k=1}^n g_k(n) \leq c \sum_{k=1}^n f_k(n) \quad \forall n \geq n_0$$

Hence, $\sum_{k=1}^n g_k(n) = O(\sum_{k=1}^n f_k(n))$.

Because the above equation holds for an arbitrary $g_k(n)$, we have:

$$\sum_{k=1}^n O(f_k(n)) \subseteq O(\sum_{k=1}^n f_k(n))$$

.

(2) ($B \subseteq A$)

Let $g(n) = O(\sum_{k=1}^n f_k(n))$. Then, there exist $c > 0$ and $n_0 > 0$ such that

$$0 \leq g(n) \leq c \sum_{k=1}^n f_k(n) \quad \forall n \geq n_0$$

By the linear property of summations,

$$0 \leq g(n) \leq \sum_{k=1}^n [cf_k(n)] \quad \forall n \geq n_0$$

If we choose $c_k = c$ and $n_k = n_0$ for $k \in \{1, 2, \dots, n\}$ in the above inequality, we obtain that there exist $c_k > 0$ and $n_k > 0$ for $k \in \{1, 2, \dots, n\}$, such that $0 \leq g(n) \leq c_1 f_1(n) + c_2 f_2(n) + \dots + c_n f_n(n)$ whenever $n \geq n_1, n \geq n_2, \dots, n \geq n_n$. This means that $g(n) = \sum_{k=1}^n O(f_k(n))$.

Therefore, $O(\sum_{k=1}^n f_k(n)) \subseteq \sum_{k=1}^n O(f_k(n))$.

From (1) and (2), we have $A = B$.

(3) CLR A.1-7

One method of solution uses a logarithm to convert the product to a sum:

$$\begin{aligned} \prod_{k=1}^n 2 \cdot 4^k &= 2^n \prod_{k=1}^n 4^k \\ \lg(2^n \prod_{k=1}^n 4^k) &= \lg(2^n) + \lg(\prod_{k=1}^n 4^k) \\ &= \lg(2^n) + \sum_{k=1}^n \lg(4^k) \\ &= \lg(2^n) + \lg(4) \sum_{k=1}^n k \end{aligned}$$

$$\begin{aligned}
&= \lg(2^n) + 2\lg(2) \frac{n(n+1)}{2} \\
&= n\lg(2) + n(n+1)\lg(2) \\
&= n(n+2)\lg(2) \\
&= \lg(2^{n(n+2)})
\end{aligned}$$

$$\prod_{k=1}^n 2 \cdot 4^k = 2^{n(n+2)}.$$

An alternative method of solution uses product and exponent rules:

$$\begin{aligned}
\prod_{k=1}^n 2 \cdot 4^k &= 2^n \cdot 4^{\sum_{k=1}^n k} \\
&= 2^n \cdot 4^{\frac{n(n+1)}{2}} \\
&= 2^n \cdot 2^{n(n+1)} \\
&= 2^{n+n(n+1)} \\
&= 2^{n^2+2n} \\
&= 2^{n(n+2)}
\end{aligned}$$

(4) CLR A.2-4

Evaluate $\sum_{k=1}^n k^3$

Since k^3 is a monotonically increasing function,

$$\int_0^n x^3 dx \leq \sum_{k=1}^n k^3 \leq \int_1^{n+1} x^3 dx \leq \int_0^{n+1} x^3 dx$$

Hence:

$$\frac{x^4}{4} \Big|_0^n \leq \sum_{k=1}^n k^3 \leq \frac{x^4}{4} \Big|_0^{n+1}$$

Giving:

$$\frac{n^4}{4} \leq \sum_{k=1}^n k^3 \leq \frac{(n+1)^4}{4}$$

Hence, $\sum_{k=1}^n k^3 = \Theta(n^4)$

(5) CLR A.2-5

If you bound the summation using an integral directly (the integral bounds are from 0 to n) the final solution becomes $[l n n - l n 0]$.

However, $l n 0$ is negative infinity.

So instead we break up the summation

$$\sum_{k=1}^n 1/k \text{ and rewrite it as } \sum_{k=2}^n 1/k + 1.$$

(6) CLR A-1

(a) Use the integral method of bounding the summation:

$$\begin{aligned} \int_0^n x^r dx &\leq \sum_{k=1}^n k^r \leq \int_1^{n+1} x^r dx \\ \frac{1}{r+1} x^{r+1} \Big|_0^n &\leq \sum_{k=1}^n k^r \leq \frac{1}{r+1} x^{r+1} \Big|_1^{n+1} \\ \frac{n^{r+1}}{r+1} &\leq \sum_{k=1}^n k^r \leq \frac{(n+1)^{r+1} - 1}{r+1} \leq \frac{(n+1)^{r+1}}{r+1} \end{aligned}$$

Therefore,

$$\sum_{k=1}^n k^r = \Theta(n^{r+1})$$

(b) On one hand,

$$\sum_{k=1}^n \lg^s k \leq \sum_{k=1}^n \lg^s n = n \lg^s n.$$

Clearly, $\sum_{k=1}^n \lg^s k = O(n \lg^s n)$.

On the other hand,

$$\begin{aligned} \sum_{k=1}^n \lg^s k &= \sum_{k=1}^{\lfloor n/2 \rfloor} \lg^s k + \sum_{k=\lfloor n/2 \rfloor + 1}^n \lg^s k \\ &\geq \sum_{k=1}^{\lfloor n/2 \rfloor} 0 + \sum_{k=\lfloor n/2 \rfloor + 1}^n \lg^s(n/2) \\ &\geq (n/2) \lg^s(n/2). \end{aligned}$$

We can easily prove that $(n/2) \lg^s(n/2) = \theta(n \lg^s n)$.

So $\sum_{k=1}^n \lg^s k = \theta(n \lg^s n)$.

(c)

On one hand,

$$\sum_{k=1}^n k^r \lg^s k \leq \sum_{k=1}^n k^r \lg^s n = \lg^s n \sum_{k=1}^n k^r$$

and from part (a)

$$\sum_{k=1}^n k^r = \Theta(n^{r+1})$$

so,

$$\sum_{k=1}^n k^r l g^s k = O(n^{r+1} l g^s n)$$

On the other hand,

$$\sum_{k=1}^n k^r l g^s k = \sum_{k=1}^{n/2} k^r l g^s k + \sum_{k=n/2}^n k^r l g^s k$$

$$\sum_{k=1}^n k^r l g^s k \geq \sum_{k=1}^{n/2} 0 + \sum_{k=n/2}^n (n/2)^r l g^s (n/2)$$

$$\sum_{k=1}^n k^r l g^s k \geq (n/2)^{r+1} l g^s (n/2)$$

And then combining the above two results it can be shown that,

$$\sum_{k=1}^n k^r l g^s k = \Theta(n^{r+1} l g^s n)$$