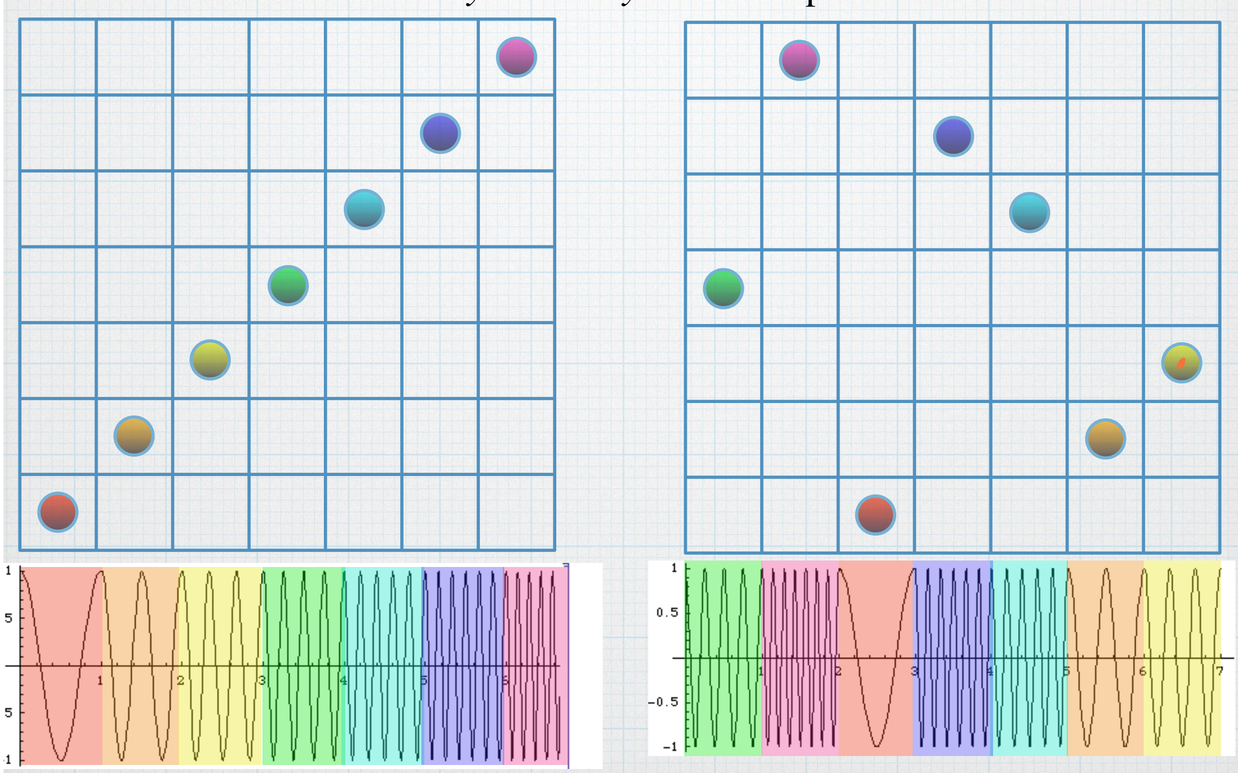
Session 34

34.1

Frequency-Coded Waveforms

Geometric Array or *Binary Matrix* Representation



The Ambiguity Function of Frequency-Coded Waveforms

The ambiguity function of $s(t) = \sum_{l=0}^{N-1} p(t - lT)e^{-j2\pi\Omega_l t}$ is

$$\chi_s(\tau,\nu) = \chi_s^{(1)}(\tau,\nu) + \chi_s^{(2)}(\tau,\nu),$$

where $\chi_s(\tau,\nu) \triangleq \int_{-\infty} s(t)\, s^*(t-\tau)\, e^{+i2\pi\nu t}\, dt$

$$\chi_s^{(1)}(\tau,\nu) = \sum_{m=0}^{N-1} e^{-j2\pi m\nu T} e^{-j2\pi\Omega_m \tau} \chi_p(\tau,\nu),$$

and

$$\chi_s^{(2)}(\tau,\nu) = \sum_{m=0}^{N-1} \sum_{n=0, n\neq m}^{N-1} e^{-j\pi(\Omega_m+\Omega_n)\tau} e^{-j\pi(m+n)T} \cdot \chi_p(\tau + (m-n)T, \nu + (\Omega_n - \Omega_m))$$

n.b. $\beta_s(\tau,\nu) = \chi_s(\tau,-\nu)$.

The sidelobes are given by

$$\chi_s^{(2)}(\tau,\nu) = \sum_{m=0}^{N-1} \sum_{n=0, n\neq m}^{N-1} e^{-j\pi(\Omega_m+\Omega_n)\tau} e^{-j\pi(m+n)T} \cdot \chi_p(\tau + (m-n)T, \nu + (\Omega_n - \Omega_m))$$

$$\chi_p(\tau + (m-n)T, \nu + (d_n - d_m)/T)$$

Large contribution when *these* equal zero!

$\tau = (n-m)T$ and $\nu = (d_n - d_m)/T)$

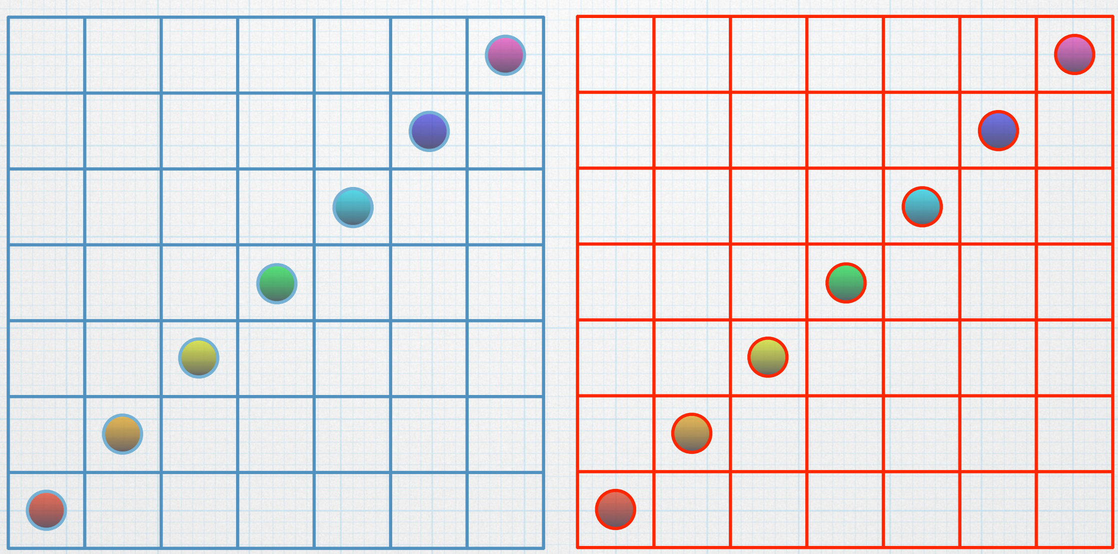
or taking $T = 1$ for simplicity...

$\tau = n - m$ and $\nu = d_n - d_m$

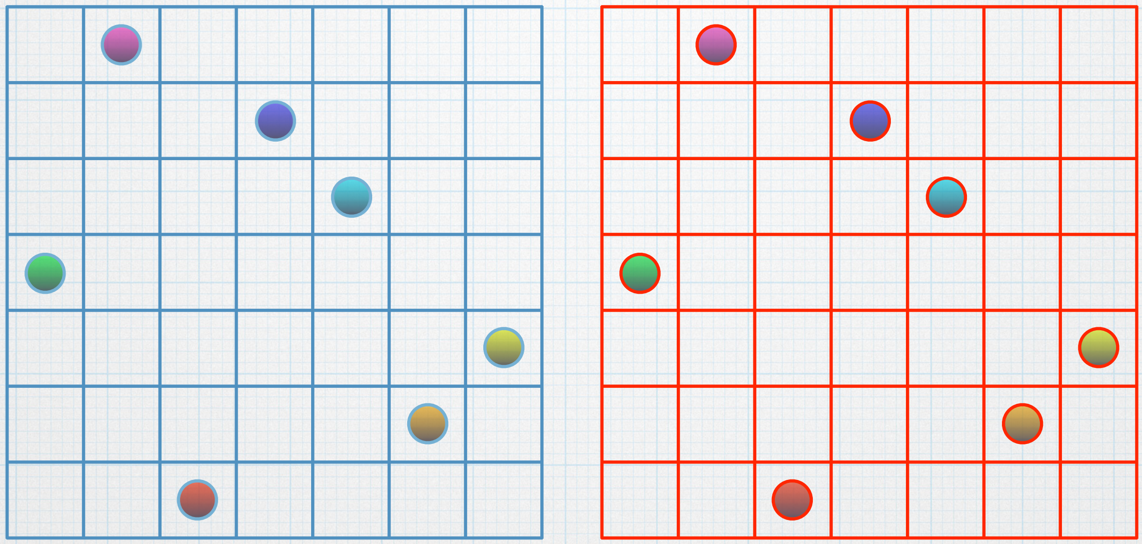
Coincident Sidelobe Approximation

- If we consider only the sidelobe contributions due to the situations where both arguments of the ambiguity function is zero, we want to minimize the number of situations where this occurs.
- We especially want to minimize multiple "hits" for any given delay and Doppler shift.
- While this approach only minimizes an approximation of the ambiguity function sidelobes, it is surprisingly effective.
- It is, in fact, the approach John Costas used in designing Costas sequences.

LFM Chirp Sidelobe Overlay Demo



Costas Sidelobe Overlay Demo

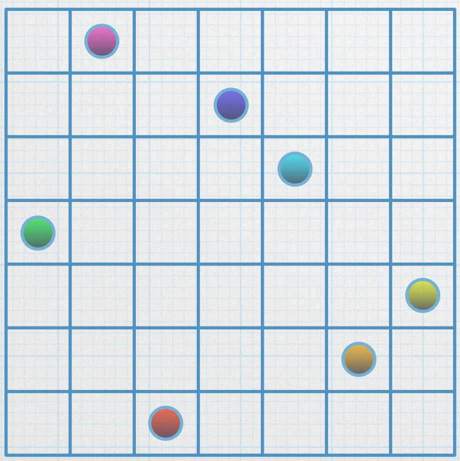


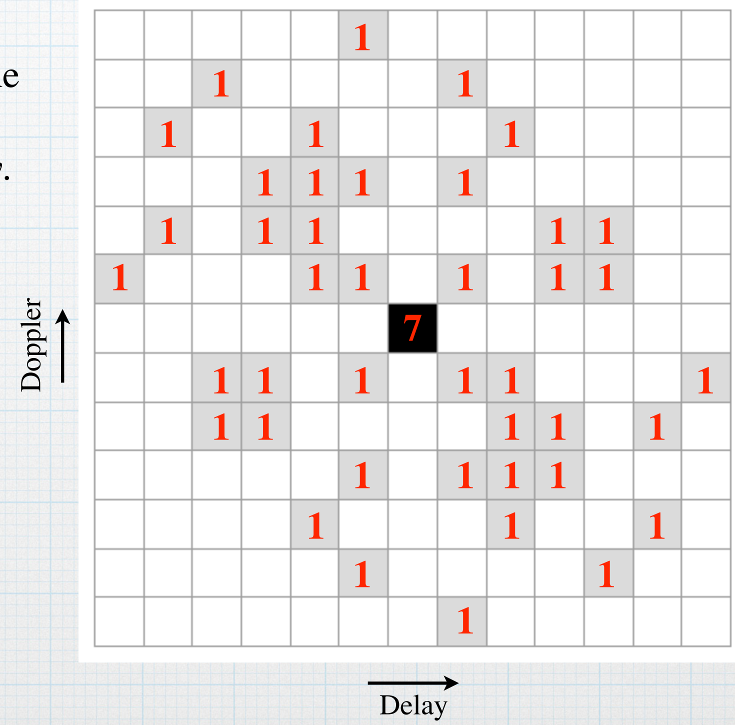
Characteristics of Stepped-Frequency Waveforms

- A wide variety of waveforms with different ambiguity functions can be generated.
- These waveforms can be easily generated and amplified for transmission.
- The ambiguity characteristics of these waveforms can be easily visualized because of their localization in time and frequency.
- Provides a straightforward approach to characterizing “ambiguity state” of a target environment.
- These characteristics make them ideal for adaptive waveform radar.

The Sidelobe Array

- If we count up the number of sidelobe coincidences for each combination of integer delay-Doppler shifts, we can tabulate the coincidences in an array called the *sidelobe array*.

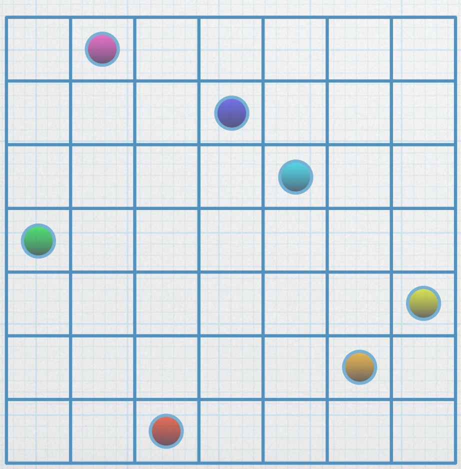


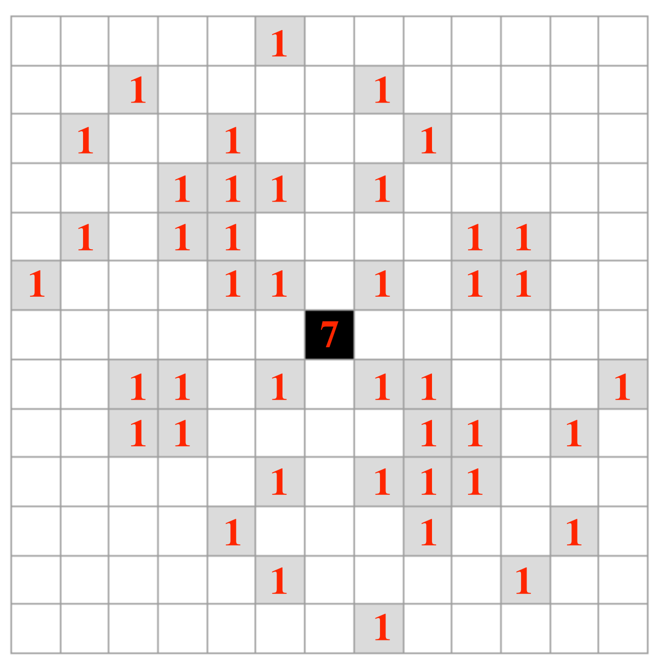


Costas Sequences

Definition: A **Costas sequence** of length N is a integer frequency firing sequence $\{d_1, \ldots, d_N\}$ (or $\{d_0, \ldots, d_{N-1}\}$ that is a permutation of the integers $1, \ldots, N$ (or $0, \ldots, N-1$) such that the maximum sidelobe height orcoincidence number in the sidelobe array is 1 for any nonzero integer delay-Doppler shift.

An Example ...





34.10

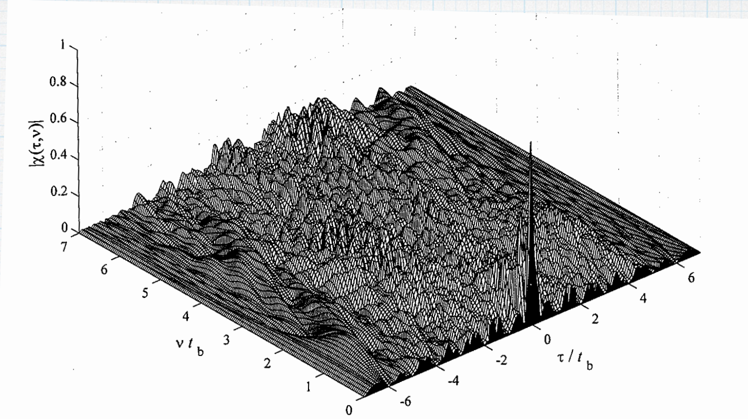


FIGURE 5.4 Partial ambiguity function of a Costas signal with code sequence {4 7 1 6 5 2 3}.

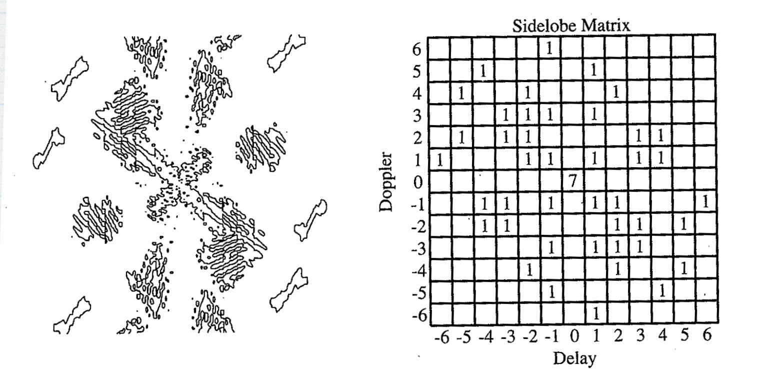


FIGURE 5.5 Ambiguity function contour at 0.125 (left) compared with the sidelobe matrix (right).

Reference: N. Levanon and E. Mozeson, ***Radar Signals***, Wiley, 2004 (ISBN 0-471-47378-2)

34.11

- A Length 40 Costas Sequence:

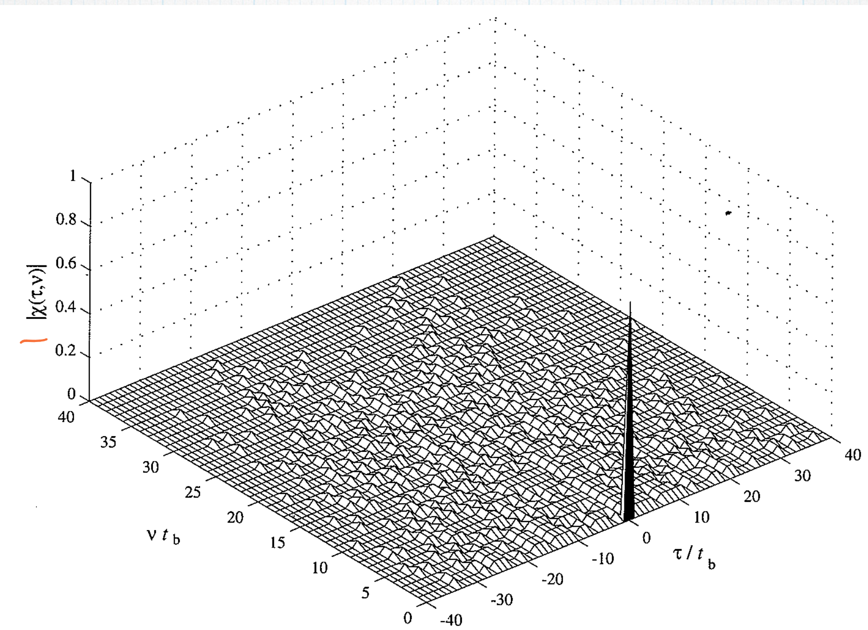


FIGURE 5.9 Ambiguity function of a Costas signal (length $M = 40$) at all relevant grid points.

Reference: N. Levanon and E. Mozeson, ***Radar Signals***, Wiley, 2004 (ISBN 0-471-47378-2)

34.12

- A Length 40 Costas Sequence:

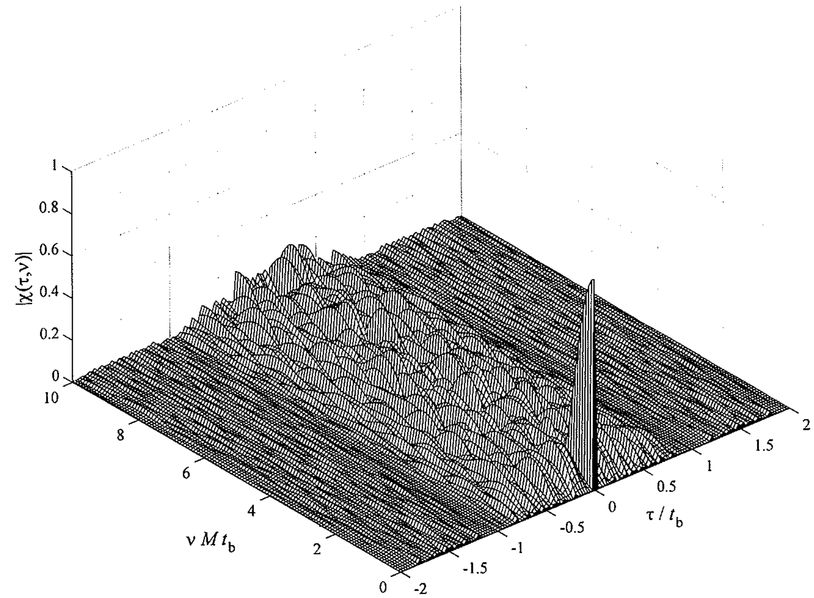


FIGURE 5.10 Ambiguity function of a Costas signal (length $M = 40$) zoom near the origin.

Reference: N. Levanon and E. Mozeson, ***Radar Signals***, Wiley, 2004 (ISBN 0-471-47378-2)

34.13

- A Length 40 Costas Sequence:

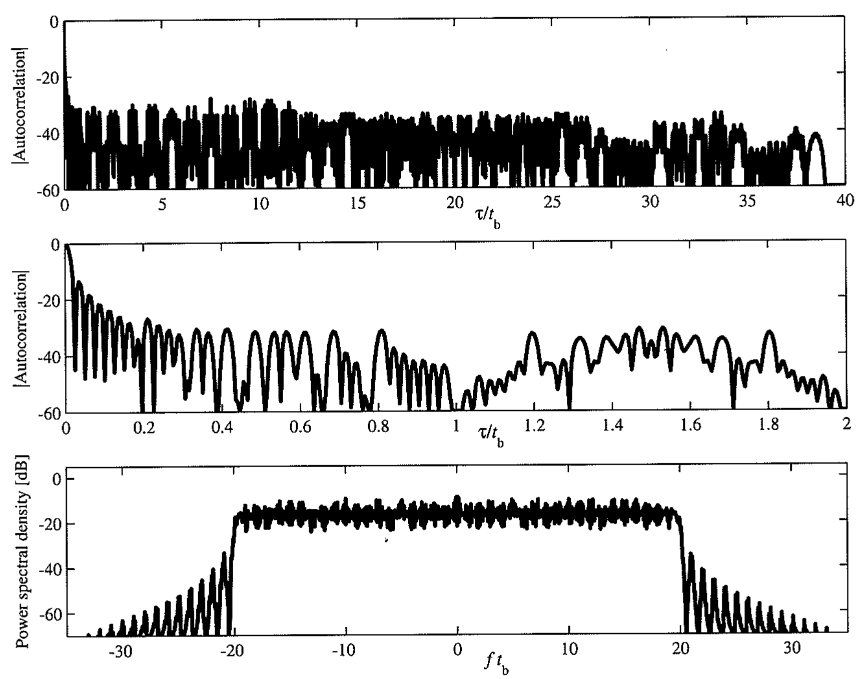


FIGURE 5.11 ACF (top and middle) and the spectrum (bottom) of a Costas signal (length 40).

Reference: N. Levanon and E. Mozeson, ***Radar Signals***, Wiley, 2004 (ISBN 0-471-47378-2)

Pushing Sequences:

A new class of Frequency-Coded Waveforms for Use in Adaptive Waveform Radar

Chieh-Fu Chang and Mark R. Bell, "Frequency-coded Waveforms for Enhanced Delay-Doppler Resolution," *IEEE Transactions on Information Theory,* vol. 49, no. 11, Nov. 2003, pp. 2960–2971.

The Ambiguity Function of Frequency-Coded Waveforms

The ambiguity function of $s(t) = \sum_{l=0}^{N-1} p(t - lT)e^{-j2\pi\Omega_l t}$ is

$$\chi_s(\tau, \nu) = \chi_s^{(1)}(\tau, \nu) + \chi_s^{(2)}(\tau, \nu),$$

where

$$\chi_s^{(1)}(\tau, \nu) = \sum_{m=0}^{N-1} e^{-j2\pi m\nu T} e^{-j2\pi\Omega_m \tau} \chi_p(\tau, \nu),$$

and

$$\chi_s^{(2)}(\tau, \nu) = \sum_{m=0}^{N-1} \sum_{n=0, n\neq m}^{N-1} e^{-j\pi(\Omega_m + \Omega_n)\tau} e^{-j\pi(m+n)T} \cdot \chi_p(\tau + (m-n)T, \nu + (\Omega_n - \Omega_m))$$

34.16

Characteristics of Stepped-Frequency Waveforms

- A wide variety of waveforms with different ambiguity functions can be generated.
- These waveforms can be easily generated and amplified for transmission.
- The ambiguity characteristics of these waveforms can be easily visualized because of their localization in time and frequency.
- Provides a straightforward approach to characterizing "ambiguity state" of a target environment.
- These characteristics make them ideal for adaptive waveform radar.

34.17

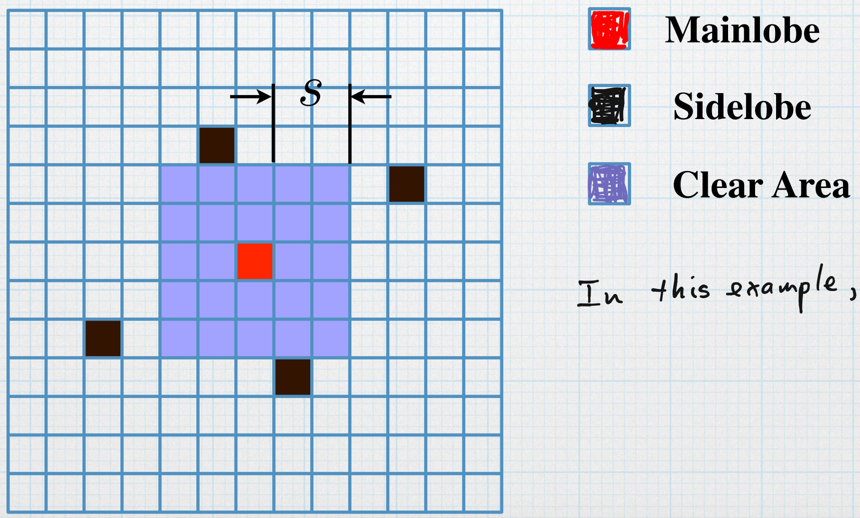
Pushing Sequences

- Pushing Sequences are frequency coded sequences that have a *clear region* clear of sidelobes surrounding the main lobe.
- *Costas sequences* approximate an ideal thumbtack ambiguity function globally. *Pushing sequences* do so locally.
- Pushing sequences are constructed using much the same intuition that is used for constructing Costas sequences (*difference matrix determination of sidelobes*.)
- Unlike Costas sequences, the frequency sequence need not be a permutation of 1, ... , *N*. Some of the frequencies may left out.
- Arbitrarily large clear areas can be achieved if arbitrarily long sequences are allowed.

34.18

Pushing Sequences

Definition: For the ambiguity function of a signal $s(t)$, a *clear area* of *size* s is a square area centered at the origin of the (τ, ν)-plane, where $|\tau| \leq sT_r$ and $|\nu| \leq s/T$, such that no sidelobe peaks are present in this area.



In this example, $s = 2$.

34.19

Pushing Sequences

Definition: A sequence having the ambiguity function with a clear area of size s is called a *pushing sequence with power* s, where $s \geq 1$.

Any sequence $\{\underline{d}_N\}$ satisfying either $|i-j| > s$ or $|d_i - d_j| > s$ for all i, j, where $0 \leq i, j \leq N-1$ and $i \neq j$, will have a clear area of size s and is thus a pushing sequence with power s. This property for a frequency coding sequence is called the *pushing property*.

We are interested in pushing sequences that efficiently fill the geometric array.

Constructing Pushing Sequences

Lemma: A Costas sequence derived from the Lempel T_4 construction is a pushing sequence of power 1.

Lee codes can be used to construct pushing sequences.

An *r*-error- correcting Lee code is a length 2 code having close-packed codewords in the geometric representation plane.

The *Lee metric* between codewords must be at least 2*r*+1.

Such codes exist for all positive *r*.

Constructing Pushing Sequences

Theorem: For every positive integer r, the codewords $\{(k, (2r \oplus 1)k)\}$ form a close-packed r-error correcting dictionary in the Lee metric, where $k = 0, 1, 2...N-1$, $N = 2r^2 + 2r + 1$ and $\oplus$ represents addition modulo N. In that case, the Lee metric between each pair of codewords is at least $2r+1$.

Theorem: If the hits exist at $(i, (2r \oplus 1)i)$ in the geometric array of $\{\underline{d}_N\}$, where $i = 0, 1, 2...N-1$, $N = 2r^2 + 2r + 1$, r is a positive integer and $\oplus$ represents addition modulo N, then $\{\underline{d}_N\}$ is a pushing sequence with power r.

So the geometric array of a pushing sequence of power *r* is given by the corresponding Lee Code and can be easily constructed.

Sidelobe Locations and Heights

Theorem: For a Lee pushing sequence with power r, the level of the sidelobe peak at

$$(\tau, \nu) = k_1 V_1 + k_2 V_2,$$

where k_1 and k_2 are integers, $V_1 = (r+1, r)$ and $V_2 = (r, -(r+1))$, is given by

$$l(k_1, k_2) = \left\lfloor \frac{(2r+1-|k_1+k_2|)(2r+1-|k_1-k_2|)}{2} \right\rfloor$$

when $|k_1|, |k_2| \leq (2r-1)$ and $|k_1| + |k_2| \leq 2r$, and 0 otherwise. Furthermore, these are the only sidelobes.