ECE 573: Compilers & Translation Systems Engineering

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Tuesday, January 11, 2011

Announcements

• “Problem Set” 0 online
• Short questionnaire
• Project step 0 due this Friday
• Use subject line “ECE 573: Project Step 0”
• Let me know if you plan to work with a partner
• Send me an email even if you don’t!

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What is a compiler?

• Traditionally: Program that analyzes and translates from a high level language (e.g., C++) to low-level assembly language that can be executed by hardware

```plaintext
int a, b;
a = 3;
if (a < 4) {
b = 2;
} else {
b = 3;
}
```

```plaintext
var a
var b
mov 3 a
mov 4 r1
cmpl a r1
jge l_e
mov 2 b
jmp l_d
l_e:
mov 3 b
l_d:
; done
```

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Compilers are translators

• Fortran
• C
• C++
• Java
• Text processing language
• HTML/XML
• Command & Scripting Languages
• Natural language
• Domain specific languages

```plaintext
Machine code
Virtual machine code
Transformed source code
Augmented source code
Low-level commands
Semantic components
Another language
```

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Compilers are optimizers

• Can perform optimizations to make a program more efficient

```plaintext
int a, b, c;
b = a + 3;
c = a + 3;
```

```plaintext
var a
var b
var c
mov a r1
addi 3 r1
mov r1 b
mov a r2
addi 3 r2
mov r2 c
```

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Why do we need compilers?

• Compilers provide portability
• Old days: whenever a new machine was built, programs had to be rewritten to support new instruction sets
• IBM System/360 (1964): Common Instruction Set Architecture (ISA) — programs could be run on any machine which supported ISA
• Common ISA is a huge deal (note continued existence of x86)
• But still a problem: when new ISA is introduced (EPIC) or new extensions added (x86-64), programs would have to be rewritten
• Compilers bridge this gap: write new compiler for an ISA, and then simply recompile programs!
Why do we need compilers? (II)

- Compilers enable high performance and productivity
- Old: programmers wrote in assembly language, architectures were simple (no pipelines, caches, etc.)
- Close match between programs and machines — easier to achieve performance
- New: programmers write in high level languages (Ruby, Python), architectures are complex (superscalar, out-of-order execution, multicore)
- Compilers are needed to bridge this semantic gap
- Compilers let programmers write in high level languages and still get good performance on complex architectures

Semantic Gap

This would be impossible without compilers!

Some common compiler types

1. High level language \(\Rightarrow\) assembly language (e.g., gcc)
2. High level language \(\Rightarrow\) machine independent bytecode (e.g., javac)
3. Bytecode \(\Rightarrow\) native machine code (e.g., java's JIT compiler)
4. High level language \(\Rightarrow\) high level language (e.g., domain specific languages, many research languages—including mine!)

HLL to Assembly

- Compiler converts program into assembly
- Assembler is machine-specific translator which converts assembly into machine code
- Conversion is usually one-to-one with some exceptions
  - Program locations
  - Variable names

HLL to Bytecode

- Compiler converts program into machine independent bytecode
  - e.g., javac generates Java bytecode, C# compiler generates CIL
- Interpreter then executes bytecode “on-the-fly”
  - Bytecode instructions are “executed” by invoking methods of the interpreter, rather than directly executing on the machine
- Aside: what are the pros and cons of this approach?
**Bytecode to Assembly**

- Compiler converts program into machine independent bytecode
  - e.g., javac generates Java bytecode, C# compiler generates CIL
- Just-in-time compiler compiles code *while program executes* to produce machine code
- Is this better or worse than a compiler which generates machine code directly from the program?

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**Structure of a Compiler**

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**Scanner**

- Compiler starts by seeing only program text
  - Scanner converts program text into string of tokens
  
  ```
  if (a < 4) {
    b = 5
  }
  ```

---

**Scanner**

- Compiler starts by seeing only program text
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  ```
  'i' 'f' '(' 'a' '<' '4' ')' '{' 'b' '=' '5' '}'
  ```

---

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  if (ID(a) OP(<=) LIT(4)) {
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  if (ID(a) OP(<=) LIT(4)) {
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  }
  ```

- But we still don't know what the *syntactic structure* of the program is
Parser

- Converts string of tokens into a parse tree or an abstract syntax tree.
- Captures syntactic structure of code (i.e., “this is an if statement, with a then-block”)

Semantic actions

- Interpret the semantics of syntactic constructs
- Note that up until now we have only been concerned with what the syntax of the code is
- What’s the difference?

Syntax vs. Semantics

- Syntax: “grammatical” structure of language
  - What symbols, in what order, is a legal part of the language?
  - But something that is syntactically correct may mean nothing!
  - “colorless green ideas sleep furiously”
- Semantics: meaning of language
  - What does a particular set of symbols, in a particular order, mean?
  - What does it mean to be an if statement?
  - “evaluate the conditional, if the conditional is true, execute the then clause, otherwise execute the else clause”

A note on semantics

- How do you define semantics?
  - Static semantics: properties of programs
    - All variables must have a type
    - Expressions must use consistent types
    - Can define using attribute grammars
    - Execution semantics: how does a program execute?
    - Can define an operational or denotational semantics for a language
    - Not covered in this class (take CS 565)
  - For many languages, “the compiler is the specification”
Symbol tables

- A list of every declaration in the program, along with other information
  - Variable declarations: types, scope
  - Function declarations: return types, # and type of arguments

Program Example

```plaintext
Integer ii;
...
ii = 3.5;
...
print ii;
```

Symbol Table

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Scope</th>
</tr>
</thead>
<tbody>
<tr>
<td>ii</td>
<td>int</td>
<td>global</td>
</tr>
</tbody>
</table>

Intermediate representation

- Also called IR
- A (relatively) low level representation of the program
  - But not machine-specific!
- One example: three address code

```plaintext
bge a, 4, done
mov 5, b
done: //done!
```

- Each instruction can take at most three operands (variables, literals, or labels)
- Note: no registers!

Optimizer

- Transforms code to make it more efficient
- Different kinds, operating at different levels
  - High-level optimizations
    - Loop interchange, parallelization
    - Operates at level of AST, or even source code
  - Scalar optimizations
    - Dead code elimination, common sub-expression elimination
    - Operates on IR
  - Local optimizations
    - Strength reduction, constant folding
    - Operates on small sequences of instructions

Code generation

- Generate assembly from intermediate representation
  - Select which instructions to use
  - Schedule instructions
  - Decide which registers to use

```plaintext
bge a, 4, done
mov 5, b
done: //done!
```

```plaintext
ld a, r1
mov 4, r2
cmp r1, r2
bge done
mov 5, r3
st r3, b
done:
```

Overall structure of a compiler

Use regular expressions to define tokens. Can then use scanner generators such as lex or flex.

Define language using context free grammars. Can then use parser generators such as yacc or bison.

Semantic routines done by hand. But can be formalized.

Written manually. Automation is an active research area (e.g., dataflow analysis frameworks)

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Use regular expressions to define tokens. Can then use scanner generators such as lex or flex.

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Many of these passes can be combined

Design considerations (I)

- Compiler and language designs influence each other
- Higher level languages are harder to compile
  - More work to bridge the gap between language and assembly
  - Flexible languages are often harder to compile
- Dynamic typing (Ruby, Python) makes a language very flexible, but it is hard for a compiler to catch errors (in fact, many simply won’t)

Design considerations (II)

- Compiler design is influenced by architectures
  - CISC vs. RISC
    - CISC designed for days when programmers wrote in assembly
      - For a compiler to take advantage of string manipulation instructions, it must be able to recognize them
    - RISC has a much simpler instruction model
  - Easier to compile for