# Processing control structures

#### Statement lists

- So far we have discussed generating code for one assignment statement
- Generating code for multiple statements is easy

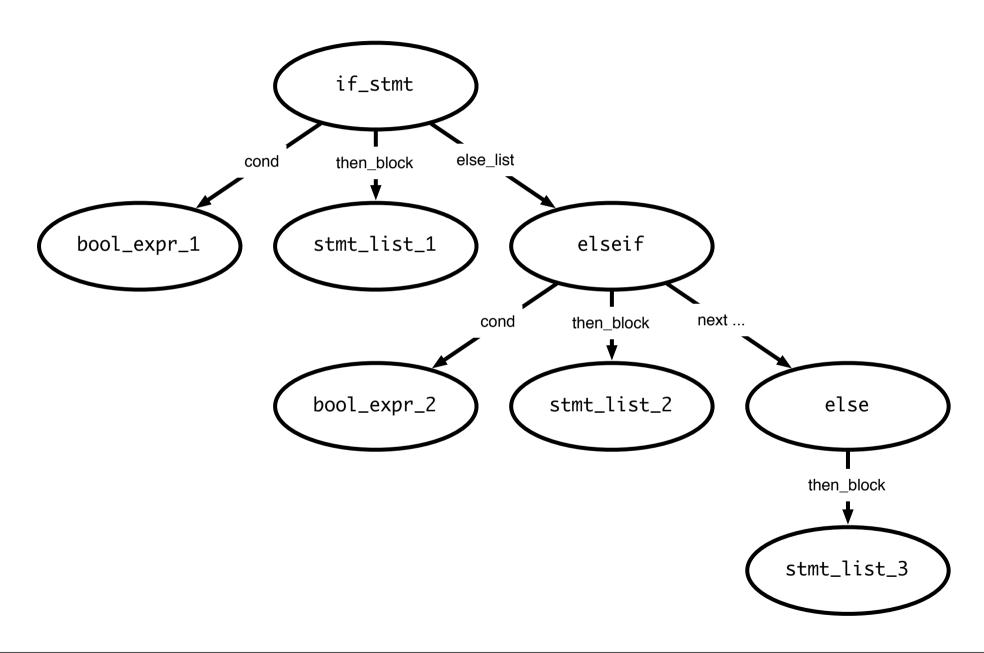
stmt\_list 
$$\rightarrow$$
 stmt\_list  $\mid \lambda$ 

- Keep appending (or prepending) the code generated by a single statement to the code generated by the rest of the statement list
- What if statement is not an assignment?

#### If statements

```
if <bool_expr_1> then
    <stmt_list_1>
elseif <bool_expr_2> then
    <stmt_list_2>
...
else
    <stmt_list_3>
endif
```

#### If statements



## Generating code for ifs

```
if <bool_expr_1> then
    <stmt_list_1>
elseif <bool_expr_2> then
    <stmt_list_2>
else
    <stmt_list_3>
endif
```

```
<code for bool_expr_1>
  j<!op> ELSE_1
  <code for stmt_list_1>
  jmp OUT

ELSE_1:
  <code for bool_expr_2>
  j<!op> ELSE
  <code for stmt_list_2>
  jmp OUT

ELSE:
  <code for stmt_list_3>
OUT:
```

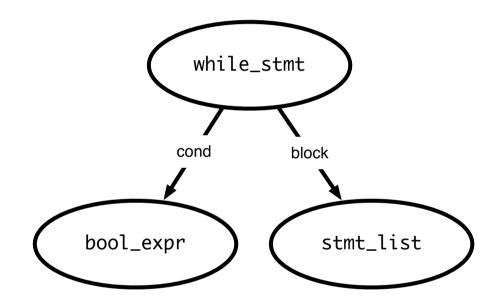
## Notes on code generation

- The <op> in j<!op> is dependent on the type of comparison you are doing in <bool\_expr>
- When you generate JUMP instructions, you should also generate the appropriate LABELs
  - But you may not put the LABEL into the code immediately
    - e.g., the OUT label (when should you create this? When should you put this in code?)
  - Instead, generate the labels when you first process the if statement (i.e., before you process the children) so that it's available when necessary
- Remember: labels have to be unique!

## Processing Loops

#### While loops

while <bool\_expr> do
 <stmt\_list>
end



## Generating code for do-while loops

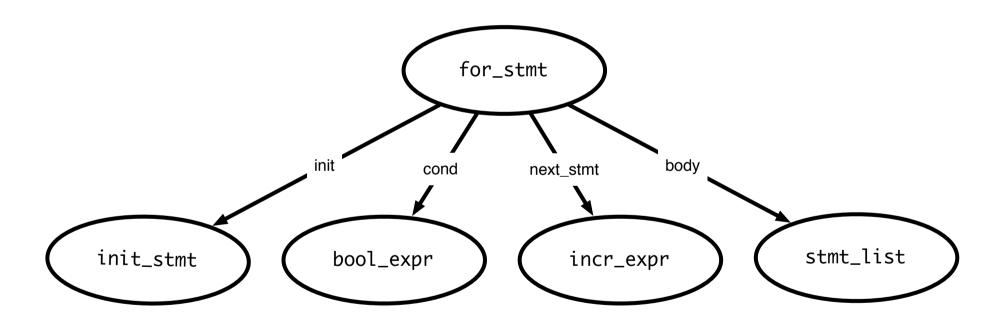
```
do
     <stmt_list>
while <bool_expr>;

LOOP:
     <stmt_list>
     <bool expr>
     j<op> LOOP
OUT:
```

- Note that we j<op> instead of j<!</li>
  - Jump when the expression is true
- Re-evaluate expression each time
- Question: what would code for "repeat until" loop look like?

#### For loops

```
for (<init_stmt>;<bool_expr>;<incr_stmt>)
    <stmt_list>
end
```



## Generating code: for loops

```
for (<init_stmt>;<bool_expr>;<incr_stmt>)
     <stmt_list>
end
```

```
<init_stmt>
LOOP:
    <bool_expr>
    j<!op> OUT
    <stmt_list>
INCR:
    <incr_stmt>
    jmp LOOP
OUT:
```

- Execute init\_stmt first
- Jump out of loop if bool\_expr is false
- Execute incr\_stmt after block, jump back to top of loop
- Question: Why do we have the INCR label?

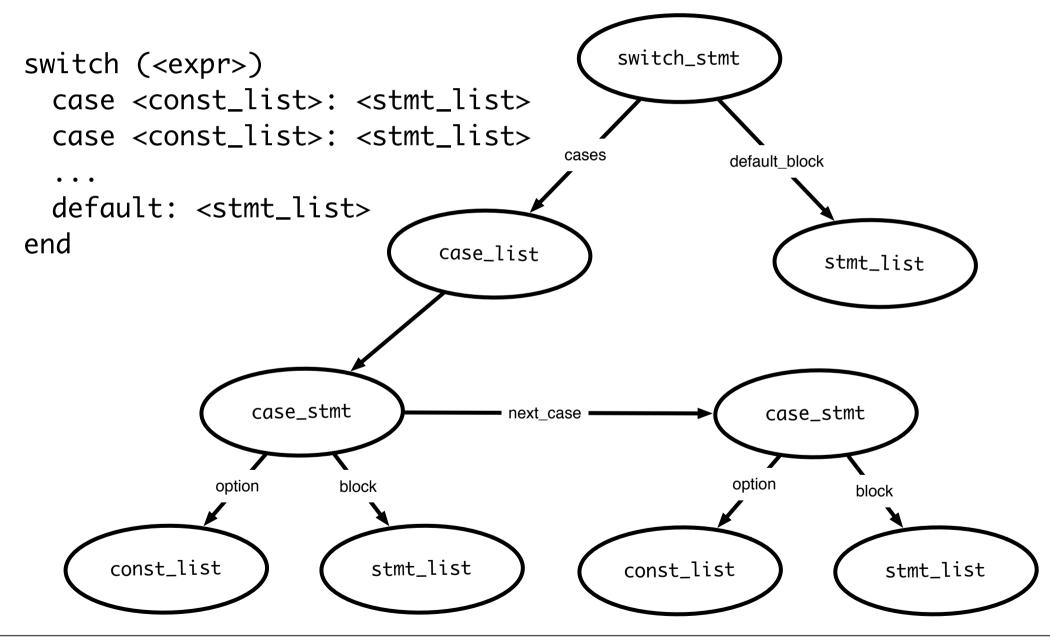
#### continue and break statements

```
for (<init_stmt>;<bool_expr>;<incr_stmt>)
     <stmt_list>
end
```

```
<init_stmt>
LOOP:
    <bool_expr>
    j<!op> OUT
    <stmt_list>
INCR:
    <incr_stmt>
    jmp LOOP
OUT:
```

- Continue statements: skip past rest of block, perform incr\_stmt and restart loop
- Break statements: jump out of loop (do not execute incr\_stmt)
- Caveats:
  - Code for stmt\_list is generated earlier—where do we jump?
  - Keep track of "loop depth" as you descend through AST

#### Switch statements



#### Switch statements

```
switch (<expr>)
  case <const_list>: <stmt_list>
  case <const_list>: <stmt_list>
  ...
  default: <stmt_list>
end
```

- Generated code should evaluate <expr> and make sure that some case matches the result
- Question: how to decide where to jump?

### Deciding where to jump

- Problem: do not know which label to jump to until switch expression is evaluated
- Use a jump table: an array indexed by case values, contains address to jump to
  - If table is not full (i.e., some possible values are skipped),
     can point to a default clause
    - If default clause does not exist, this can point to error code
  - Problems
    - If table is sparse, wastes a lot of space
    - If many choices, table will be very large

### Jump table example

Consider the code: ((xxxx) is address of code)

Case x is

(0010) When 0: stmts

(0017) When 1: stmts

(0192) When 2: stmts

(0198) When 3 stmts;

(1000) When 5 stmts;

(1050) Else stmts;

Table only has one Unnecessary row (for choice 4)

Jump table has 6 entries:

0	JUMP 0010
	JUMP 0017
2	JUMP 0192
3	JUMP 0198
4	JUMP 1050
5	JUMP 1000

### Jump table example

Consider the code: ((xxxx) Is address of code)

Case x is

(0010) When 0: stmts0

(0017) When I: stmts I

(0192) When 2: stmts2

(0198) When 3 stmts3

(1000) When 987 stmts4

(1050) When others stmts5

Table only has 983 unnecessary rows. Doesn't appear to be the right thing to do! NOTE: table size is proportional to range of choice clauses, not number of clauses!

Jump table has 6 entries:

0	JUMP 0010
I	JUMP 0017
2	JUMP 0192
3	JUMP 0198
4	JUMP 1050
• • •	JUMP 1050
986	JUMP 1050
987	JUMP 1000

#### Do a binary search

Consider the code: ((xxxx) Is address of code)

Case x is

(0010) When 0: stmts0

(0017) When I: stmts I

(0192) When 2: stmts2

(0198) When 3 stmts3

(1000) When 987 stmts4

(1050) When others stmts5

Jump table has 6 entries:

0	JUMP 0010
	JUMP 0017
2	JUMP 0192
3	JUMP 0198
987	JUMP 1000

Perform a binary search on the table. If the entry is found, then jump to that offset. If the entry isn't found, jump to others clause. O(log n) time, n is the size of the table, for each jump.

#### Linear search example

Consider the code:
(xxxx) Is offset of local
Code start from the
Jump instruction

Case x is
(0010) When 0: stmts
(0017) When 1: stmts
(0192) When 2: stmts
(1050) When others stmts;

If there are a small number of choices, then do an in-line linear search. A straightforward way to do this is generate code analogous to an IFTHEN ELSE.

```
If (x == 0) then stmts I;
Elseif (x = 1) then stmts 2;
Elseif (x = 2) then stmts 3;
Else stmts 4;
```

O(n) time, n is the size of the table, for each jump.

## Dealing with jump tables

```
switch (<expr>)
  case <const_list>: <stmt_list>
  case <const_list>: <stmt_list>
  default: <stmt list>
end
      <expr>
      <code for jump table>
    IABFI0:
      <stmt list>
    IABFI1:
      <stmt_list>
    DEFAULT:
      <stmt list>
    OUT:
```

- Generate labels, code, then build jump table
  - Put jump table after generated code
- Why do we need the OUT label?
  - In case of break statements