Semantic actions for declarations and expressions

Semantic actions

- *Semantic actions* are routines called as productions (or parts of productions) are recognized
- Actions work together to build up intermediate representations

  <if-stmt> → IF <expr> #startif THEN <stmts> END #endif

  - Semantic action for #startif needs to pass a *semantic record* to #endif
- For LL parsers, semantic actions work easily, because they are predictive
- For LR parsers, do not know which production is used until reduce step; need to place semantic actions at end of production

  <if-stmt> → <begin-if> THEN <stmts> END #endif

  <begin-if> → IF <expr> #startif

Semantic Records

- Data structures produced by semantic actions
- Associated with both non-terminals (code structures) and terminals (tokens/symbols)
    - Do not have to exist (e.g., no action associated with ";")
- Control statements often require multiple actions (see <if-stmt> example on previous slide)
- Typically: semantic records are produced by actions associated with terminals, and are passed to actions associated with non-terminals
- Standard organization: *semantic stack*

Example of semantic stack

- Consider following grammar:

  assign → ID := expr
  expr → term addop term
  term → ID | LIT
  addop → + | –

- And now annotated with semantic actions:

  assign → ID `#process_id` := expr `#gen_assign`
  expr → term addop term `#gen_infix`
  term → ID `#process_id` | LIT `#process_lit`
  addop → + `#process_p` | – `#process_m`

Example of semantic stack

- Consider a := b + 1;
- Sequence of semantic actions invoked:

  process_id, process_id, process_op, process_lit, gen_infix, gen_assign

How do we manipulate stack?

- *Action-controlled*: actions directly manipulate stack (call push and pop)
- *Parser-controlled*: parser automatically manipulates stack

LR-parser controlled

- Shift operations push semantic records onto stack (describing the token)
- Reduce operations pop semantic records associated with symbols off stack, replace with semantic record associated with production
- Action routines do not see stack. Can refer to popped off records using handles
    - e.g., in yacc/bison, use $1, $2 etc. to refer to popped off records

LL-controlled

- Parse stack contains predicted productions, not matched productions
- Push empty semantic records onto stack when production is predicted
- Fill in records as symbols are matched
- When non-terminal is matched, pop off records associated with RHS, use to fill in the record associated with LHS (leave LHS record on stack)

Overview of declarations

- Symbol tables
- Action routines for simple declarations
- Action routines for advanced features
  - Constants
  - Enumerations
  - Arrays
  - Structs
  - Pointers

Symbol Tables

- Table of declarations, associated with each scope
- One entry for each variable declared
  - Store declaration *attributes* (e.g., name and type) – will discuss this in a few slides
- Table must be dynamic (why?)
- Possible implementations
  - Linear list (easy to implement, only good for small programs)
  - Binary search trees (better for large programs, but can still be slow)
  - Hash tables (best solution)
- BSTs and Hash tables can be difficult to implement, but languages like C++ and Java provide implementations for you

Managing symbol tables

- Maintain list of all symbol tables
- Maintain stack marking "current" symbol table
- Whenever you see a program block that allows declarations, create a new symbol table
  - Push onto stack as "current" symbol table
- When you see declaration, add to current symbol table
- When you exit a program block, pop current symbol table off stack

Handling declarations

- Declarations of variables, arrays, functions, etc.
  - Create entry in symbol table
  - Allocate space in *activation record*
    - Activation record stores information for a particular function call (arguments, return value, local variables, etc.)
      - Need to have space for all of this information
    - Activation record stored on program stack
    - We will discuss these in more detail when we get to functions

Simple declarations

- Declarations of simple types

  ```
  INT x;

  FLOAT f;
  ```

- Semantic action should
  - Get the type and name of identifier
  - Check to see if identifier is already in the symbol table
    - If it isn't, add it, if it is, error

Simple declarations (cont.)

- How do we get the type and name of an identifier?

```
var_decl → var_type id;

var_type → INT | FLOAT

id → IDENTIFIER
```

- Where do we put the semantic actions?

Simple declarations (cont.)

- How do we get the type and name of an identifier?

```
var_decl → var_type id; #decl_id

var_type → INT #int_type | FLOAT #float_type

id → IDENTIFIER #id
```

- Where do we put the semantic actions?
  - When we process #int_type and #id, can store the type and identifier name and pass them to #decl_id
- When creating activation record, allocate space based on type (why?)

Constants and ranges

- Constants
  - Symbol table needs a field to store constant value
  - In general, the constant value may not be known until runtime (`static final int i = 2 + j;`)
  - At compile time, we create code that allows the initialization expression to assign to the variable, then evaluate the expression at run-time
- Range types (like in Pascal)

  ```
  Type alpha = ‘a’ .. ‘z’
  ```

  - Need an entry for the type as well as the upper and lower bounds

Enums

- Enumeration types: `enum days {mon, tue, wed, thu, fri, sat, sun};`
  - Create an entry for the enumeration type itself, and an entry for each member of the enumeration
    - Entries are usually linked
  - Processing enum declaration sets the "enum counter" to lower bound (usually 0)
  - Each new member seen is assigned the next value and the counter is incremented
    - In some languages (*e.g.*, C), enum members may be assigned particular values. Should ensure that enum value isn't reused

Arrays

- Fixed size (static) arrays

  `int A[10];`

  - Store type and length of array
  - When creating activation record, allocate enough space on stack for array

- What about variable size arrays?

  `int A[M][N]`

  - Store information for a *dope vector*
    - Tracks dimensionality of array, size, location
    - Activation record stores dope vector
    - At runtime, allocate array at top of stack, fill in dope vector

Structs/classes

- Can have variables/methods declared inside, need extra symbol table
  - Need to store visibility of members
- Complication: can create multiple instances of a struct or class!
  - Need to store *offset* of each member in struct

Pointers

- Need to store type information and length of what it points to
  - Needed for pointer arithmetic

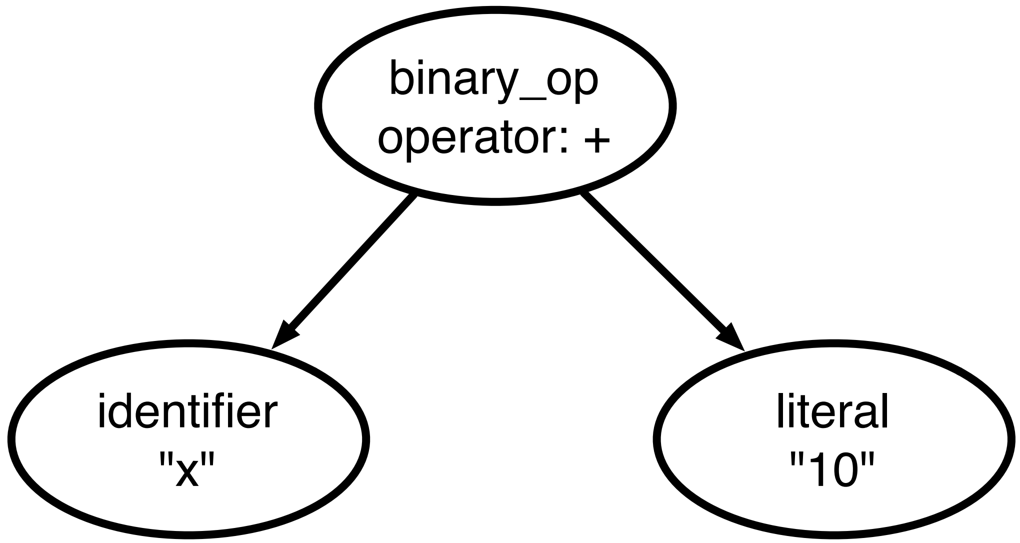
    ```
    int * a = &y;
    z = *(a + 1);
    ```

- Need to worry about forward declarations
  - The thing being pointed to may not have been declared yet

    ```
    Class Foo;
    Foo * head;
    Class Foo { ... };
    ```

Abstract syntax trees

- Tree representing structure of the program
  - Built by semantic actions
  - Some compilers skip this
- AST nodes
  - Represent program construct
  - Store important information about construct



ASTs for References

Referencing identifiers

- Different behavior if identifier is used in a declaration vs. expression
  - If used in declaration, treat as before
  - If in expression, need to:
    - Check if it is symbol table
    - Create new AST node with pointer to symbol table entry
    - Note: may want to directly store type information in AST (or could look up in symbol table each time)

Referencing Literals

- What about if we see a literal?

  primary → INTLITERAL | FLOATLITERAL

- Create AST node for literal
- Store string representation of literal
  - “155”,“2.45” etc.
- At some point, this will be converted into actual representation of literal
  - For integers, may want to convert early (to do *constant folding*)
  - For floats, may want to wait (for compilation to different machines). Why?

More complex references

- Arrays
  - A[i][j] is equivalent to

    `A + i*dim_1 + j`
  - Extract `dim_1` from symbol table or dope vector
- Structs
  - A.f is equivalent to

    `&A + offset(f)`
  - Find `offset(f)` in symbol table for declaration of record
- Strings
  - Complicated–depends on language

Expressions

- Three semantic actions needed
  - eval_binary (processes binary expressions)
    - Create AST node with two children, point to AST nodes created for left and right sides
  - eval_unary (processes unary expressions)
    - Create AST node with one child
  - process_op (determines type of operation)
    - Store operator in AST node

Expressions example

- x + y + 5

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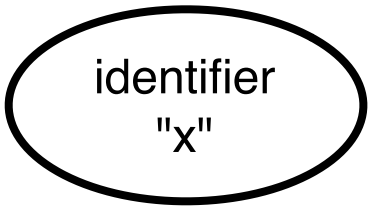

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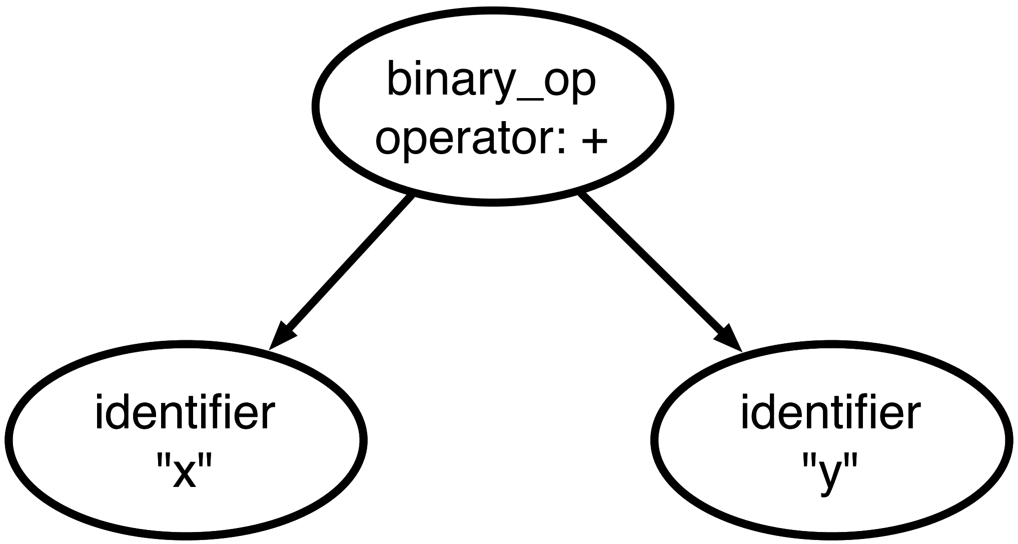

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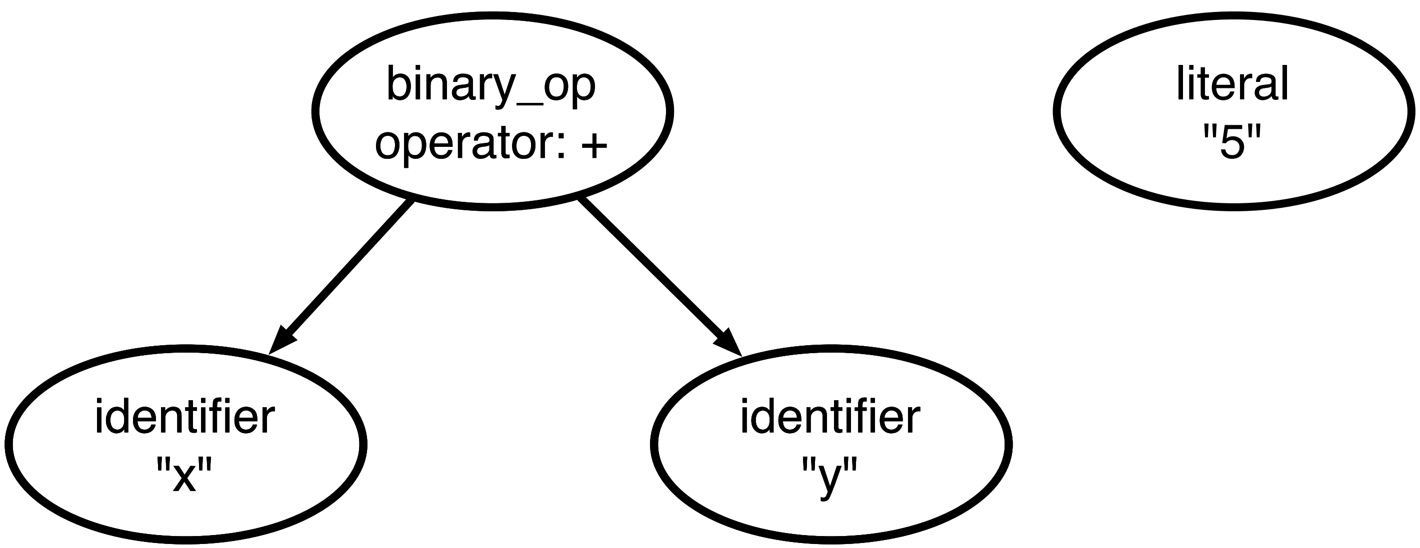

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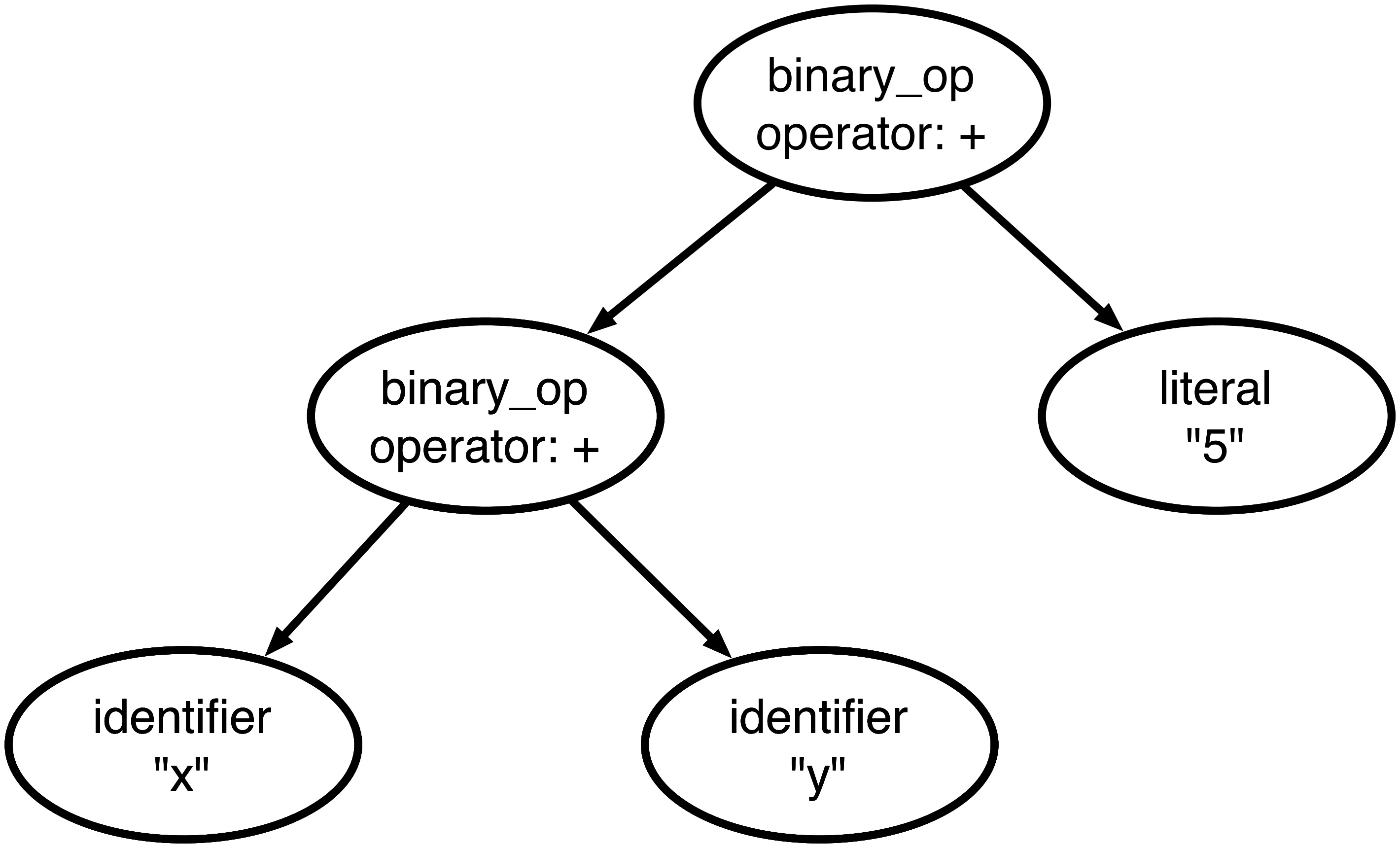

Generating three-address code

- For project, will need to generate three-address code
  - op A, B, C //C = A op B
- Can do this directly or after building AST

Generating code from an AST

- Do a post-order walk of AST to generate code, pass generated code up

```
data_object generate_code() {
  data_object lcode = left.generate_code();
  data_object rcode = right.generate_code();
  return generate_self(lcode, rcode);
}
```

- Important things to note:
  - A node generates code for its children before generating code for itself
  - Data object can contain code or other information
  - Code generation is *context free*
    - What does this mean?

Generating code directly

- Generating code directly using semantic routines is very similar to generating code from the AST
  - Why?
  - Because post-order traversal is essentially what happens when you evaluate semantic actions as you pop them off stack
    - LL parser: evaluate left child before right child
    - LR parser: evaluate right child before left child
  - AST nodes are just semantic records

Data objects

- Records various important info
  - The temporary storing the result of the current expression
  - Flags describing value in temporary
    - Constant, L-value, R-value
  - Code for expression

L-values vs. R-values

- L-values: addresses which can be stored to or loaded from
- R-values: data (often loaded from addresses)
  - Expressions operate on R-values
- Assignment statements:

  `L-value := R-value`

- Consider the statement `a := a`
  - the `a` on LHS refers to the memory location referred to by `a` and we store to that location
  - the `a` on RHS refers to data *stored in* memory location referred to by `a` so we will load from that location to produce the R-value

Temporaries

- Can be thought of as an unlimited pool of registers (with memory to be allocated at a later time)
- Need to declare them like variables
- Name should be something that cannot appear in the program (*e.g.*, use illegal character as prefix)
- Memory must be allocated if address of temporary can be taken (e.g. $a := \&b$)
- Temporaries can hold either L-values or R-values

Simple cases

- Generating code for constants/literals
  - Store constant in temporary
  - Optional: pass up flag specifying this is a constant
- Generating code for identifiers
  - Generated code depends on whether identifier is used as L-value or R-value
    - Do we load from it? Or store to it?
    - One solution: just pass variable up to next level
      - Set flag specifying this is an L-value

Generating code for expressions

- Create a new temporary for result of expression
- Examine data-objects from subtrees
- If temporaries are L-values, load data from them into new temporaries
    - Generate code to perform operation
- If temporaries are constant, can perform operation immediately
    - No need to perform code generation!
- Store result in new temporary
    - Is this an L-value or an R-value?
- Return code for entire expression

Generating code for assignment

- Store value of temporary from RHS into address specified by temporary from LHS
  - Why does this work?
  - Because temporary for LHS holds an address
    - If LHS is an identifier, we just stored the address of it in temporary
    - If LHS is complex expression

      ```
      int *p = &x
      *(p + 1) = 7;
      ```

      it *still* holds an address, even though the address was computed by an expression