

* Error Decoding Via Table-based search.

Method # 1: Standard Array.

$$\mathcal{C} = \{\vec{0}, \vec{c}_1, \dots, \vec{c}_{2^k-1}\}$$

Step 1: List all codewords as columns

$$\begin{array}{ccccccc} \vec{0}, & \vec{c}_1, & \vec{c}_2, & \dots & & \vec{c}_{2^k-1} \\ \vec{e}, & \vec{e} + \vec{c}_1, & \dots & & & \vec{e} + \vec{c}_{2^k-1} \end{array}$$

Step 2: Set $\mathcal{V}_{\text{error}} = \{0, 1\}^n$

Step 3: Choose $\vec{e} \in \mathcal{V}_{\text{error}}$ with the smallest weight

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Step 4: Insert $\vec{e} + \vec{0}, \vec{e} + \vec{c}_1, \dots, \vec{e} + \vec{c}_{2^k-1}$,
as a new row.

Step 5: Remove 2^k vectors from $\mathcal{V}_{\text{error}}$.

$$\mathcal{V}'_{\text{error}} = \mathcal{V}_{\text{error}} \setminus \{\vec{e} + \vec{0}, \vec{e} + \vec{c}_1, \dots, \vec{e} + \vec{c}_{2^k-1}\}.$$

& Go back to Step 3 if the new $\mathcal{V}_{\text{error}}$ is non-empty.

Properties:

- ① Each entry of the table is distinct. (A non-trivial property follows from the linearity of the code.)

- ② We have exactly $\frac{2^n}{2^k} = 2^{n-k}$ rows

When decoding, the decoder observes \mathcal{Y}_{obs} . It searches the entire table to

jobs. It searches the entire table to find the entry that contains \vec{y}_{obs} .

- ② Output the column index m . i.e,
 \vec{c}_m is the decoded codeword.
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Method 2: A simplified version of
(more efficient)
Method 1.

Syndrome Decoding

$$\vec{y}_{obs} = \vec{c}_m + \vec{e}$$

$$H \cdot \vec{y}_{obs} = H \cdot \vec{c}_m + H \cdot \vec{e}$$

$$= 0 + H \cdot \vec{e} \triangleq \vec{s} \text{ syndrome}$$

Properties = ① Every row corresponds to the same \vec{c} , thus the same \vec{s}

$\sim \quad \dots \quad -n-k$

\xrightarrow{s}

② Totally 2^{n-k} rows

thus at most $2^{n-k} \xrightarrow{s}$

★ ③ Each row has a unique \vec{s} , and thus we have exactly 2^{n-k} distinct \vec{s}

(A non-trivial property due to the linearity of the code)

Improved Version of Method #1:

① Pre-compute the mapping

$$f(\vec{e}) = H \cdot \vec{e} = \vec{s} \text{ for all } e.$$

② For the received \vec{y}_0s , we compute the \vec{e} of the corresponding row by

$$f^{-1}(H \cdot \vec{y}_{obs}) = \vec{e}$$

③ Finally $\vec{c}_m = \vec{y}_{obs} - \vec{e}$

$$= \vec{y}_{obs} - f^{-1}(H \cdot \vec{y}_{obs})$$

Remark: We still have to have to do table look-up for f^{-1} where

$$f^{-1}: \{0,1\}^{n-k} \mapsto \{0,1\}^n$$

Comparison, the table look-up for the previous method is

$$\begin{aligned} T: \{0,1\}^n &\mapsto \{0,1\}^n \\ \vec{y}_{obs} &\mapsto \vec{c}_m \end{aligned}$$

Exercise: Apply the syndrome decoder to binary Hamming Codes and Gray Hamming Codes

Hamming Codes.