## ECE 639, Homework #6 (CRN: 25576) Due date: Wednesday 11/16/2022 during the lecture

https://engineering.purdue.edu/~chihw/22ECE639F/22F\_ECE639.html

Question 27: [Lin, Costello Jr., Error Control Coding 2nd Ed., Problem 5.12] For a cyclic code, if an error pattern e(x) is detectable, show that its *i*-th cyclic shift  $e^{(i)}(x)$  is also detectable for all *i*.

Question 28: [Lin, Costello Jr., Error Control Coding 2nd Ed., Problem 5.14] For any cyclic code, let  $\vec{v}$  be an *n*-dimensional codeword of the given cyclic code, and  $\vec{v}^{(l)}$  is the *l*-cyclically-shifted version of  $\vec{v}$ . Prove that if we choose  $l_0 > 0$  as the smallest, non-zero *l* such that  $\vec{v} = \vec{v}^{(l)}$ , then we must have  $l_0|n$ , i.e.,  $l_0$  is a factor of *n*.

Question 29: [Lin, Costello Jr., Error Control Coding 2nd Ed., Problem 5.20] Let  $C_1$  and  $C_2$  be two cyclic codes of length n. Define  $C_3 = C_1 \cap C_2$  as the codewords that belong to both  $C_1$  and  $C_2$ . Prove that (1)  $C_3$  is cyclic; and (2) Find the generator polynomial of  $C_3$ .

*Question 30:* [Lin, Costello Jr., Error Control Coding 2nd Ed., Problem 4.7] Form the generator matrix of the RM(1,4) code of length 16. Answer the following questions:

- 1. What is the minimum distance of the code?
- 2. Decode the received vector r = (0, 0, 1, 1, 0, 0, 1, 0, 0, 1, 1, 1, 0, 0, 1, 1).

Question 31: [Lin, Costello Jr., Error Control Coding 2nd Ed., Problem 4.9] Prove that the dual code of RM(m - r - 1, m) is RM(r, m).