

40

**Mahdis Mosaheb**

Homework #3, 695C

This code can get any generating matrix and create the all possible codeword vectors. The output of a binary counter multiply to a generating matrix and create the code vector corresponding to the binary information vector.

I ran the code for G1 and G2 and the output matrices are below matrices such that every column is a codeword.

G1=

|   |   |   |   |
|---|---|---|---|
| 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 0 | 0 | 1 |

G2=

|   |   |   |   |
|---|---|---|---|
| 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |
| 1 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 |

X1 =

|   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 |   |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |

X2 =

|   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 |

### Question1

1-2) This implementation is a codeword-wise ML decoder, which gets the four probabilities (p00, p01, p10, p11) and also the Yobs vector and calculates the joint probability of Y obs with the all codewords and the codeword which has the maximum probability is going to be the decoded Yobs.

1-3) This part approximates the FER and BER of the codewise-decoder .

FER: In order to get the FER, at first we should calculate the error probability of any possible Y . That means to find the probability of decoded Y is not equal to x (all x vector).

$$\text{FER} = \sum_{y=000000}^{111111} \sum_0^{15} P_{\vec{x}, \vec{y}}(\vec{x}, \vec{y}) \{f(y) \neq x\}$$

We need to check the decoded y with all X's, in order to find which one is not equal to f(y), then get the summation of those probabilities. Also, we shouldn't forget to get the error probability for all possible y's .

BER: In order to calculate BER, we should at first find the Hamming distance between any decoded y and all possible codewords(X).Hamming distance gives us the number of bit differences of two vectors. Then we multiply the Hamming distance of two vectors(X, f(y)) by the joint probability of the same X and y . We repeat this algorithm for all possible y and then divide by the number of bits in one codeword to get the BER.

$$\text{BER} = 1/n \sum_{y=000000}^{111111} \sum_0^{15} P_{\vec{x}, \vec{y}}(\vec{x}, \vec{y}) |f(y) - x|$$

Where  $|f(y) - x|$  is hamming distance.

### Question2

2-2) To implement the Bitwise-decoder, at the first step we calculate the probability of first bit of X. for example, we calculate the probability of X(1)=1 given Yobs , if it was less than  $\frac{1}{2}$  the bitwise decoder, decides the first bit of decoded y is going to be 0 otherwise it is 1 . We repeat that for all bits of X.

The problem of this decoder is sometimes the decoded vector is not even a codeword. But this decoder minimizes the BER.

2-3) The functions of FER and BER in this question are the same as question 1 but the decoder that we use here is bit-wise decoder.,

According to the graphs, we can see the FER has the better performance by codewise decoder after the jump in the FERbitwise , before the jump, they show the equal performance. On the other hand, as we said before, BER has better performance with bitwise decoder. The BERbitwise and BERcodewise are the same at first but eventually the BERbitwise decreased from BERcodewise. So we can conclude, Bitwise decoder is more efficient for BER than Codewise decoder.

```
G = [0 1 1 1;1 0 1 1;1 1 0 1;1 0 0 0; 0 1 0 0; 0 0 1 0;0 0 0 1];
%G = [1 0 0 0;1 0 0 1;1 0 1 0;1 0 1 1;1 1 0 0; 1 1 0 1; 1 1 1 0; 1 1 1 1 ]
column = length(G(1,:));
row = length(G(:,1));
X = zeros(row,2^column);

temp=dec2bin((0:2^column-1));
for i=1:2^column
    for j=1:column
        I(j,i)=str2num(temp(i,j));
    end
end
tempy=dec2bin((0:2^row-1));
for i=1:2^row
    for j=1:row
        Y(j,i)=str2num(tempy(i,j));
    end
end
for x = 1:row;
    for z = 1:2^column;
        for y = 1:column;
            X(x, z) = xor(X(x, z) , and(G(x, y) , I(y, z)));
        end
    end
end
X1 =X
%X2 =X
```

```

G=input('G')
column = length(G(1,:));
row = length(G(:,1));
X = zeros(row,2^column);
% binary counter
temp=dec2bin((0:2^column-1));
for i=1:2^column
    for j=1:column
        I(j,i)=str2num(temp(i,j));
    end
end
for x = 1:row;
    for z = 1:2^column;
        for y = 1:column;
            X(x, z) = xor(X(x, z) , and(G(x, y) , I(y, z)));
        end
    end
end
p00=input('p00')
p01=input('p01')
p10=input('p10')
p11=input('p11')
Y1= input('Yobs')
Y= Y1.';

for x=1:2^column;
np00=0;
np01=0;
np10=0;
np11=0;
for y=1:row;
    if X(y,x)==0 & Y(y,1)==0
        np00=np00+1
    end
    if X(y,x)==0 & Y(y,1)==1
        np01=np01+1
    end
    if X(y,x)==1 & Y(y,1)==0
        np10=np10+1
    end
    if X(y,x)==1 & Y(y,1)==1
        np11=np11+1
    end
end
L(1,x)= (p00^np00)*(p01^np01)*(p10^np10)*(p11^np11);

end
[maxVal maxInd] = max(L)
m(1,1:column)=I(1:column,maxInd)
XML(1:row,1)=X(1:row,maxInd)

```

```

%G = [0 1 1 1;1 0 1 1;1 1 0 1;1 0 0 0; 0 1 0 0; 0 0 1 0;0 0 0 1];
G = [1 0 0 0;1 0 0 1;1 0 1 0;1 0 1 1;1 1 0 0; 1 1 0 1; 1 1 1 0; 1 1 1 1 ];
column = length(G(1,:));
row = length(G(:,1));
X = zeros(row,2^column);
XML = zeros(row,2^row);
%p=0.1
p=0.01:0.05:0.5;
% binary counter
temp=dec2bin((0:2^column-1));
for i=1:2^column
    for j=1:column
        I(j,i)=str2num(temp(i,j));
    end
end
tempy=dec2bin((0:2^row-1));
for i=1:2^row
    for j=1:row
        Y(j,i)=str2num(tempy(i,j));
    end
end
for x = 1:row;
    for z = 1:2^column;
        for y = 1:column;
            X(x, z) = xor(X(x, z) , and(G(x, y) , I(y, z)));
        end
    end
end
for P=1:length(p(1,:))
p00=1-p(1,P);
p01=p(1,P);
FERC = 0;
BERC =0;
BERC1=0;
for i=1:2^row;

    for x=1:2^column;

        np00=0;
        np01=0;

        for y=1:row;
            if (X(y,x)==Y(y,i))
                np00=np00+1;
            end
            if (X(y,x)~=Y(y,i))
                np01=np01+1;
            end
        end
    end
end

```

```
end

L(1,x)= (p00.^np00).* (p01.^np01);
end
[maxVal maxInd] = max(L);

XML(:,i)=X(1:row,maxInd);% at this point we have the decoded Y(i)
for k =1:2^column
mp00=0;
mp01=0;

FERi=0;
BERi=0;
np00=0;
np01=0;

for z=1:row % Hamming distance
    if X(z,k)== XML(z,i)
        mp00=mp00+1;
    end
    if X(z,k)~= XML(z,i)
        mp01=mp01+1;
    end
end

for z=1:row
    if X(z,k)==Y(z,i)
        np00=np00+1;
    end
    if X(z,k)~=Y(z,i)
        np01=np01+1;
    end
end

FL(1,k)= (p00.^np00).* (p01.^np01);

BL(1,k)= (mp01).*FL(1,k);

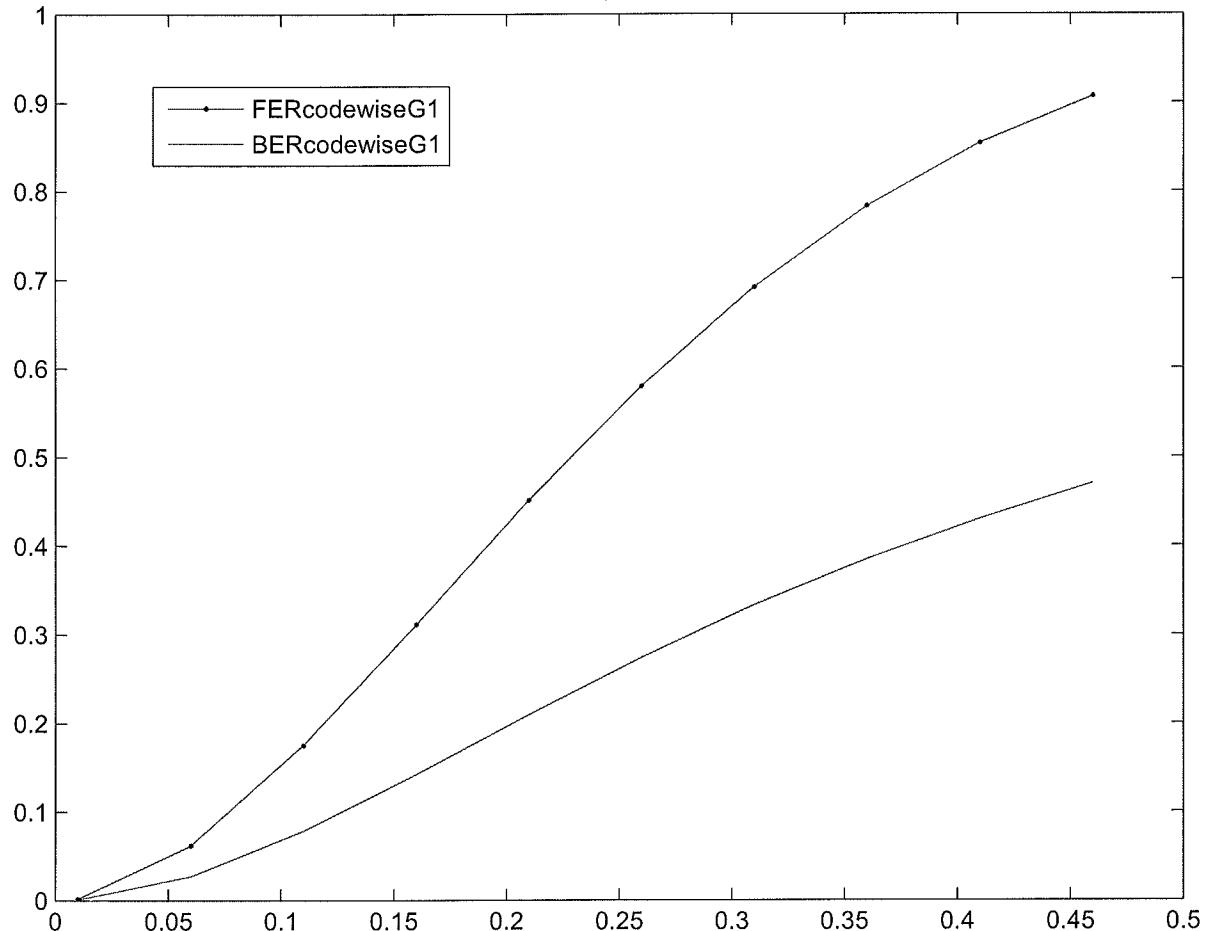
end

BERi= (1/2^column).* (sum(BL.'));
FERi=(1/2^column).* (sum(FL'))-(max(FL')));
BERC1=(BERC1+BERi);
FERC=FERC+FERi;
end
BERC = (1./row ).* BERCl;
BERcode(1,P)= BERC
FERcode(1,P)=FERC
end
```

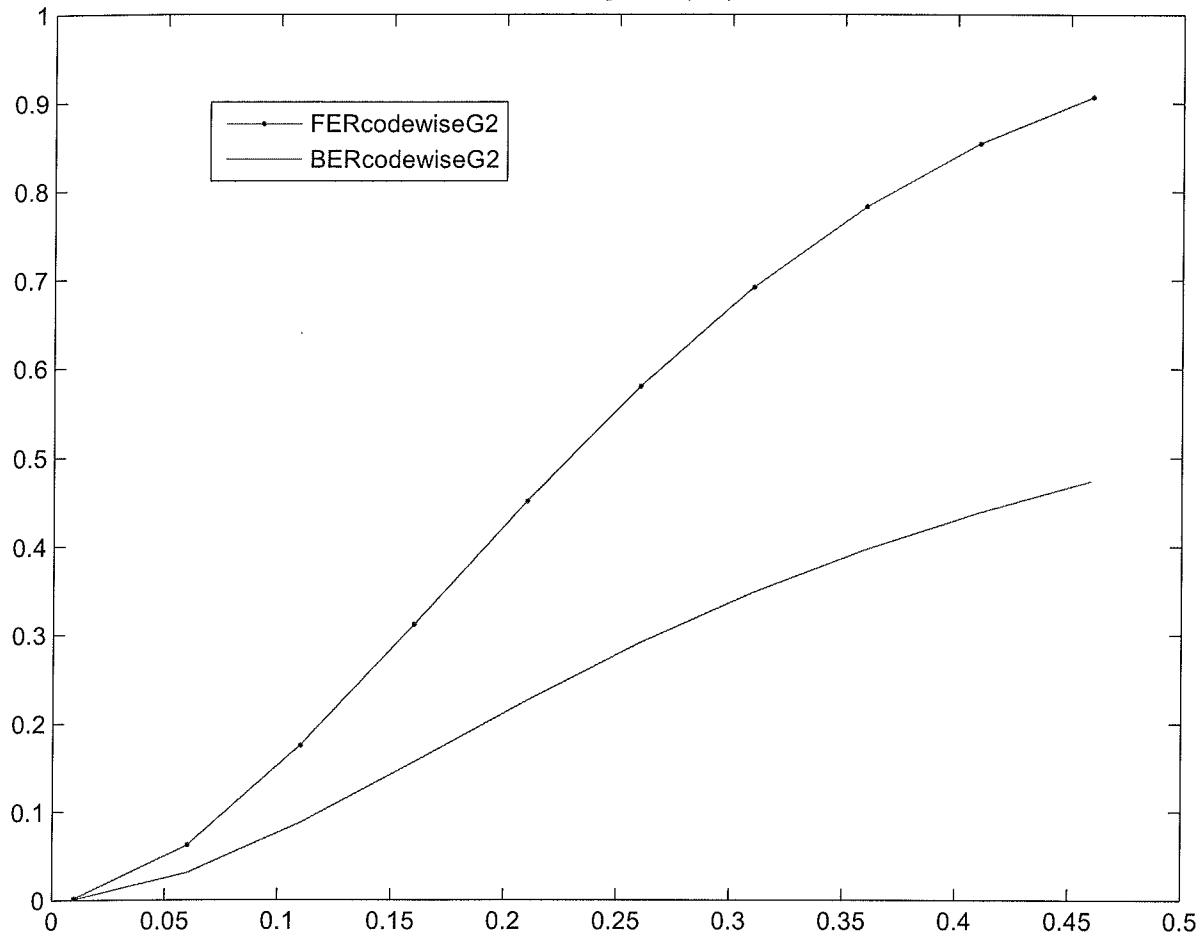
```
plot(p,FERcode,'.-')
axis([0 0.5 0 1])
hold on
plot(p,BERcode)
axis([0 0.5 0 1])
%legend('FERcodewiseG1','BERcodewiseG1');
legend('FERcodewiseG2','BERcodewiseG2');
%title('Hamming Code(7,4)')

title('Hamming Code(8,4)')
hold off
```

Hamming Code(7,4)



Hamming Code(8,4)



```
%G = [0 1 1 1;1 0 1 1;1 1 0 1;1 0 0 0; 0 1 0 0; 0 0 1 0;0 0 0 1];
G= [1 1 0 1;1 0 1 1;0 1 1 1; 1 0 0 0;0 1 0 0;0 0 1 0; 0 0 0 1];
column = length(G(1,:));
row = length(G(:,1));
X = zeros(row,2^column);
% binary counter
temp=dec2bin((0:2^column-1));
for i=1:2^column
    for j=1:column
        I(j,i)=str2num(temp(i,j));
    end
end
for x = 1:row;
    for z = 1:2^column;
        for y = 1:column;
            X(x, z) = xor(X(x, z) , and(G(x, y) , I(y, z)));
        end
    end
end
p00=input('p00')
p01=input('p01')
p10=input('p10')
p11=input('p11')
Y1= input('Yobs')
Y= Y1.';
for x=1:row;
    for y=1:2^column;
        np00=0;
        np01=0;
        np10=0;
        np11=0;
        for z=1:row;
            if X(z,y)==0 & Y(z,1)==0
                np00=np00+1;
            end
            if X(z,y)==0 & Y(z,1)==1
                np01=np01+1;
            end
            if X(z,y)==1 & Y(z,1)==0
                np10=np10+1;
            end
            if X(z,y)==1 & Y(z,1)==1
                np11=np11+1;
            end
        end
        if X(x,y)==1
            L(1,y)= (p00^np00)*(p01^np01)*(p10^np10)*(p11^np11);
            T(1,y)= (p00^np00)*(p01^np01)*(p10^np10)*(p11^np11);
        else
    end
end
```

2/16/12 10:12 AM \\\guardian.ics.purdue.edu\mm...\Bitwise.m 2 of

```
L(1,y)=0;
T(1,y)=(p00^np00)*(p01^np01)*(p10^np10)*(p11^np11);
end

end
LB=sum(L.');
LT=sum(T.');
Detector=(LB/LT);
if Detector < 0.5
    XML(1,x)=0;
else
    XML(1,x)=1;
end

end

Output=num2str(XML)
```

```
%G = [0 1 1 1;1 0 1 1;1 1 0 1;1 0 0 0; 0 1 0 0; 0 0 1 0;0 0 0 1];
G = [1 0 0 0;1 0 0 1;1 0 1 0;1 0 1 1;1 1 0 0; 1 1 0 1; 1 1 1 0; 1 1 1 1 ];
column = length(G(1,:));
row = length(G(:,1));
X = zeros(row,2^column);

p=0.05:0.02:0.45;
FERbit=zeros(1,length(p(1,:)));
% binary counter
temp=dec2bin((0:2^column-1));
for i=1:2^column
    for j=1:column
        I(j,i)=str2num(temp(i,j));
    end
end
tempy=dec2bin((0:2^row-1));
for i=1:2^row
    for j=1:row
        Y(j,i)=str2num(tempy(i,j));
    end
end
for x = 1:row;
    for z = 1:2^column;
        for y = 1:column;
            X(x, z) = xor(X(x, z) , and(G(x, y) , I(y, z)));
        end
    end
end
for P=1:length(p(1,:))
p00=1-p(1,P);
p01=p(1,P);

FERB1 = 0;
BERB =0;
BERB1=0;
for i=1:2^row;

    %Bitwise-decoder
    for x=1:row
        for y=1:2^column
            np00=0;
            np01=0;
            for z=1:row
                if X(z,y)==Y(z,i)
                    np00=np00+1;
                end
                if X(z,y)~=Y(z,i)
                    np01=np01+1;
                end
            end
        end
    end
    if np00>np01
        FERB1=FERB1+1;
    end
    if np01>np00
        BERB=BERB+1;
    end
    if np00==np01
        BERB1=BERB1+1;
    end
end
end
```

```
end
if X(x,y)==1
    L(1,y)= (p00.^np00).* (p01.^np01);
else

    L(1,y)=0;
end
T(1,y)=(p00.^np00).* (p01.^np01);

end
LB=sum(L.');
LT=sum(T.');
Detector=(LB./LT);
if Detector < 0.5
    XML(x,i)=0;
else
    XML(x,i)=1;
end
end % at this point XML (:,i) is decoded vector of Y(i)
for k =1:2^column
    mp00=0;
    mp01=0;
    FERBi=0;
    np00=0;
    np01=0;

    for z=1:row % Hamming distance
        if X(z,k)== XML(z,i)
            mp00=mp00+1;
        end
        if X(z,k) ~= XML(z,i)
            mp01=mp01+1;
        end
    end
    for z=1:row
        if X(z,k)==Y(z,i)
            np00=np00+1;
        end
        if X(z,k) ~=Y(z,i)
            np01=np01+1;
        end
    end
    if all(X(:,k)== XML(:,i))
        LEFR(1,k)=0;
    else
```

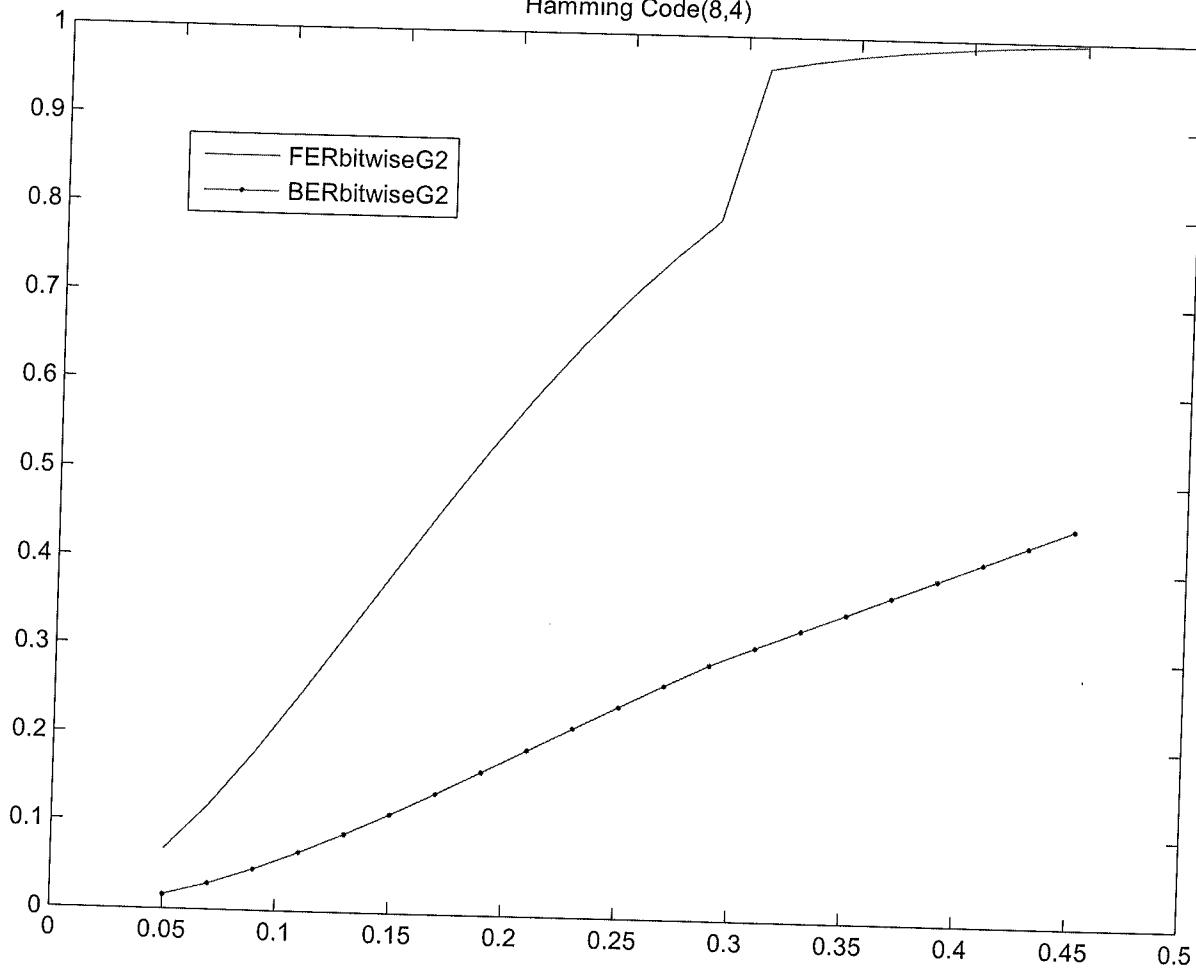
```
LFER(1,k) = (1/2^column).* (p00.^np00).* (p01.^np01);
end
BERtemp(1,k)= (mp01).* (p00.^np00).* (p01.^np01);
end
BERi=(1/2^ column).* (sum(BERtemp.'));
FERBi= sum(LFER.');
BERB1=(BERB1+BERi);
FERB1=FERB1+FERBi;
end
BERB= (1/row).* (BERB1);
BERbit(1,P)= BERB;
FERbit(1,P)=FERB1;

end

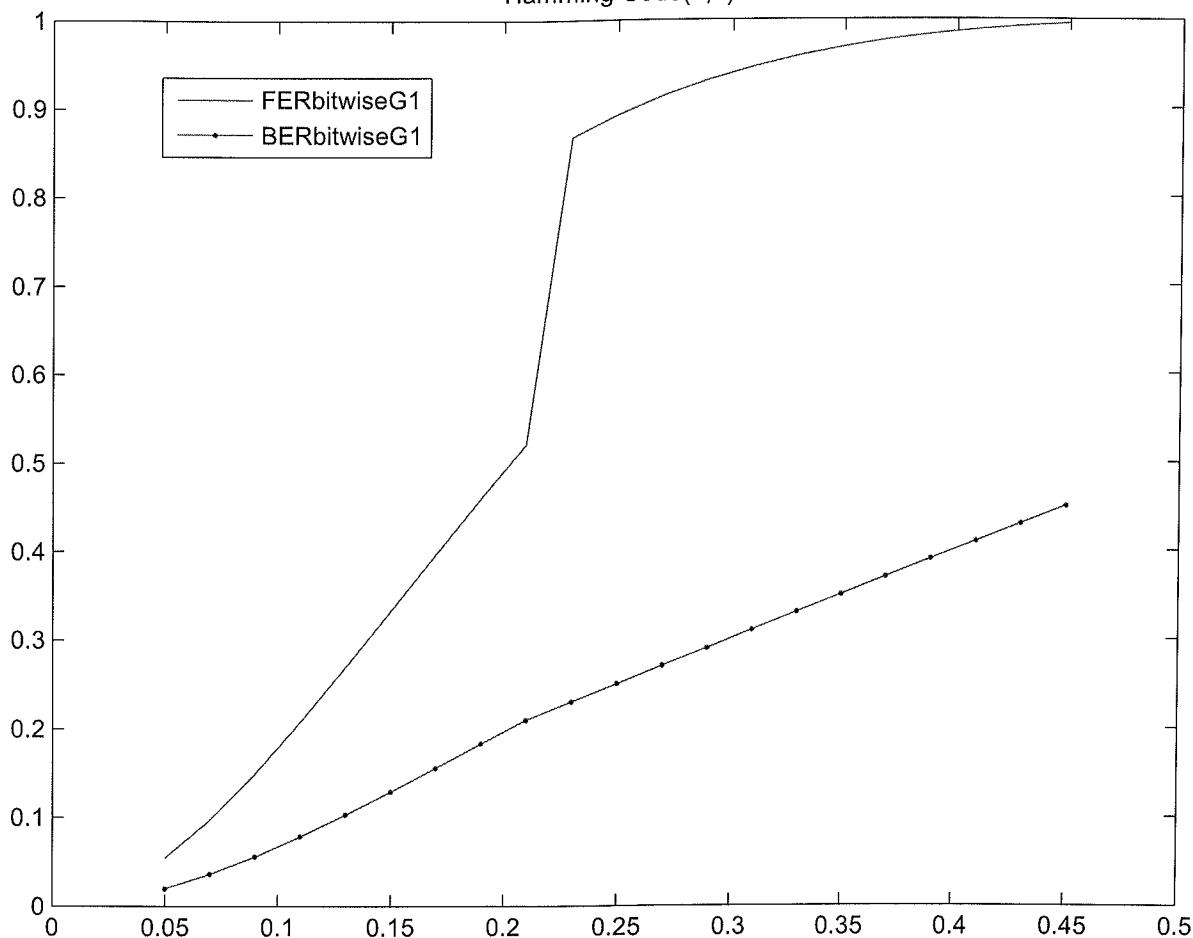
plot(p,FERbit)
axis([0 0.5 0 1])
hold on
plot(p,BERbit,'--')
axis([0 0.5 0 1])
%legend('FERbitwiseG1','BERbitwiseG1');
%legend('FERbitwiseG2','BERbitwiseG2');
%title('Hamming Code(7, 4)')

title('Hamming Code(8, 4)')
hold off
```

Hamming Code(8,4)



Hamming Code(7,4)



## HW3

Monday, October 10, 2022 9:27 AM

Q12: Suppose  $m = a \cdot b$

where  $a \geq 2$  and  $b \geq 2$  are the factors of  $m$ .

$$\begin{aligned} \text{then } a \cdot b &= a \boxtimes b \bmod m \\ &= 0. \end{aligned}$$

$\Rightarrow$   $\circ$  is not a commutative group

over  $\{0, 1, \dots, m-1\} \setminus \{0\}$

$\Rightarrow$  It is not a field.

Q13: Proof: ① Closedness:

$$(4 \cdot k_1) + (4 \cdot k_2) = (4k_1) \boxplus (4k_2) \bmod 32$$

$$= 4(k_1 \boxplus k_2) \bmod 32$$

$$= 4 \cdot ((k_1 \boxplus k_2) \bmod 8)$$

② Associativity.

② Associativity.

$$[(4k_1) + (4k_2)] + (4k_3)$$

$$= 4(k_1 \oplus k_2 \oplus k_3) \bmod 32$$

$$= (4k_1) + [(4k_2) + (4k_3)]$$

③ Identity : 0

④ Inverse :  $4k_1 \xrightarrow{\text{inverse}} 4 \cdot (8 \ominus k_1)$

$$\Rightarrow 4k_1 + 4(8 \ominus k_1)$$

$$= 4(8 \ominus k_1) + 4k_1 = 4(k_1 \oplus 8 \ominus k_1) \bmod 32$$

$$= 0.$$

---

Q14:  $X + Y + Z + W = 1$

$$X + Y + W = 1 \Rightarrow Z = 0.$$

$$X + Z + W = 0 \Rightarrow Y = 1$$

$$Y + Z + W = 0 \Rightarrow X = 1$$

$$\Rightarrow W = 1$$

\*