

Securing Ad Hoc Wireless Networks

- Introduction
- Cryptography
 Example
- IDS
 Example1
 Example2

Types of attacks:

Passive (affects confidentiality)

eavesdropping

Active (integrity, non-repudiation, availability, freshness, and authentication)

Deleting, modifying and injecting messages

Jamming (DoS)

energy exhaustion

Sources of attacks:

Internal

External

Security mechanisms: (two lines of defense)

Cryptography and authentication:

prevention but no detection

reduces attacks but does not eliminate them

protects against some types of external attacks

does not work against malicious internal nodes

Intrusion Detection (ID):

a second wall of defense

detect and may prevent attacks

works against both internal and external attacks

Example one

(cryptography and redundancy)

**Secure Data Transmission in Mobile Ad Hoc networks
(SMT)**

Communication phases that need to be secured:

1. The route discovery
2. The data transmission

Two protocols compared for secure data transmission:

1. SMT (presented)
2. SSP (Secure Single Path) (for comparison)

SMT Assumptions:

1. Secure routing protocol
2. An initial trust between source and destination
may be via public key cryptography
3. A shared finite field for purposes of data dispersion
pre-computed set of columns

SMT description:

Source:

1. Discovers the Active Path Set (ASP) to the destination
2. Add redundancy to the message
3. Disperse the message into multiple pieces
4. Guard the pieces with a MAC
5. Transmit each piece across a different route to the destination

Destination:

1. Set a timer upon reception of the first piece
2. Verify the integrity and authenticity of each received piece
3. Send a disperse and redundant feed back for each received piece

SMT Description ...

Source:

1. A successful feedback means operational path
2. An unsuccessful feedback means broken or compromised path
3. Based on the feedback each path is given a rating

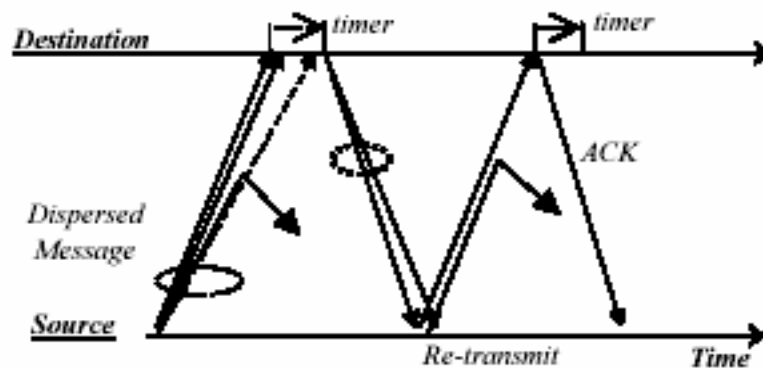
Destination:

1. Reconstruct the message if sufficient pieces are received or,
2. Ask for retransmission if timer runs out

Source:

1. Sets a counter for retransmission
2. Repeat retransmission trials until success or counter runs out

SMT description ...



SMT message dispersion:

Raw data (message) is viewed as a stream of m-bit characters

divide the message into L messages each of length M

add padding if necessary

establish an $M \times L$ matrix (B)

Establish an $N \times M$ matrix (A), $N > M$

any M combination of rows are independent

row k of $A = \{1, u_k, \dots, (u_k)^{M-1}\}$, u_k is random from field

Each piece sent $w_i = a_i \times B$, a_i the i th row of A

If we receive M pieces (v_1, \dots, v_M) each correspond to an a_i

establish an $M \times L$ matrix (W')

establish from them an $M \times M$ matrix (A')

$$B = [A']^{-1} \times W'$$

$$(i). \quad A = \begin{pmatrix} 1 & 45 & dc \\ 1 & 7d & f5 \\ 1 & b0 & 7a \\ 1 & 5b & 95 \end{pmatrix}$$

$$(ii). \quad B = \begin{pmatrix} c & e2 & 9f & 5 & c0 & 6 & 4c & 62 & c7 & 80 & 6f & 8a & 87 & 8c & 7a & 2c & 31 & 4a & 5a & ea & a5 & 5 \\ 70 & c6 & eb & 99 & b2 & 21 & 30 & 23 & 15 & 86 & 5d & 85 & 72 & 2a & 7f & 42 & b & 3f & cc & 9b & 89 & c5 \\ 8d & 46 & ba & 7f & f9 & 61 & 19 & 49 & a9 & 93 & d1 & e1 & 75 & 3 & 74 & 62 & bf & 77 & f4 & f4 & 4a & 8 \end{pmatrix}$$

$$(iii). \quad W = \begin{pmatrix} 67 & 27 & 5d & 11 & 81 & 4c & 13 & 8a & 6d & 34 & 33 & cb & a0 & 2c & df & 8 & 9d & 2 & 53 & 95 & 8a & ca \\ 54 & 24 & b1 & f9 & 60 & d5 & 31 & 7b & 1e & 1d & 1f & 23 & 7a & 49 & 2d & 57 & 58 & 2d & 75 & 1a & f & b0 \\ e8 & a2 & d5 & ca & 62 & 5f & 7a & 4a & ff & 87 & b1 & a0 & c1 & ef & e5 & ae & 5a & 44 & b9 & ee & d7 & 15 \\ 5f & 17 & a2 & af & d2 & 9a & d6 & c7 & 6b & 87 & e2 & aa & a3 & a4 & 3 & 17 & 9d & d6 & 7f & cd & b1 & b6 \end{pmatrix}$$

$$(iv). \quad W' = \begin{pmatrix} 67 & 27 & 5d & 11 & 81 & 4c & 13 & 8a & 6d & 34 & 33 & cb & a0 & 2c & df & 8 & 9d & 2 & 53 & 95 & 8a & ca \\ 54 & 24 & b1 & f9 & 60 & d5 & 31 & 7b & 1e & 1d & 1f & 23 & 7a & 49 & 2d & 57 & 58 & 2d & 75 & 1a & f & b0 \\ e8 & a2 & d5 & ca & 62 & 5f & 7a & 4a & ff & 87 & b1 & a0 & c1 & ef & e5 & ae & 5a & 44 & b9 & ee & d7 & 15 \end{pmatrix}$$

$$(v). \quad A' = \begin{pmatrix} 1 & 45 & dc \\ 1 & 7d & f5 \\ 1 & b0 & 7a \end{pmatrix}$$

$$(vi). \quad B = [A']^{-1} \times W' = \begin{pmatrix} c & e2 & 9f & 5 & c0 & 6 & 4c & 62 & c7 & 80 & 6f & 8a & 87 & 8c & 7a & 2c & 31 & 4a & 5a & ea & a5 & 5 \\ 70 & c6 & eb & 99 & b2 & 21 & 30 & 23 & 15 & 86 & 5d & 85 & 72 & 2a & 7f & 42 & b & 3f & cc & 9b & 89 & c5 \\ 8d & 46 & ba & 7f & f9 & 61 & 19 & 49 & a9 & 93 & d1 & e1 & 75 & 3 & 74 & 62 & bf & 77 & f4 & f4 & 4a & 8 \end{pmatrix}$$

APS adaptation:

Each path has short-term rating r_s and long-term rating r_l

r_s is increased by a constant β up to r_{\max} with each success and decreased by a constant α down to r_{th} with each failure

$r_l = \text{successfully received} / \text{total sent}$

If either r_s or $r_l < r_{\text{th}}$, the path discarded

BWL (bandwidth loss) = failed/(success+failed) = $f / (s+f)$

All the constants are protocol selectable

Simulation results compared to SSP:

SMT delivers more messages

SMT has much lower end-t-end delay

The more static the network the worse the adversary effect for both SMT and SSP

Discussion:

Each source is assigned a fixed destination throughout the simulation !

The size of the buffer at the source is infinite !

The attacker drops all data packets it receives (very simple attacker) !

The attacker does not interfere with the routing traffic !

What about sensor networks ?

The overhead of multi-path discovery !

The overhead of the acknowledgment !

The overhead of pre-processing the data !

The overhead of redundant data !

The overhead of undetectable attacker !

Second Wall Of Defense

IDS

Classification of IDS:

Based on data collection mechanisms:

Host-based: OS audits and system and applications logs

Network-based: Packets captured from network traffic

Based on detection techniques:

Signature-based (misuse-detection) : Known attacks

Anomaly-based: deviation from normal behavior

Specification-Based: deviation from a set of predefined constraints

Classification of IDS ...

Based on decision-making mechanisms

Collaborative

- prone to DoS and spoofed intrusion attacks (any malicious node can trigger full-forced response affecting the entire network)
- + the amount of information obtained about each node participating in the network is sufficient.

Independent

- the amount of information obtained is limited.
- + far less prone to spoofing attacks

Primary assumptions in IDS:

User and program activities are observable

Normal and intrusion activities have distinct behavior

IDS for Wired networks:

Distributed Probing

(e.g. a router detects neighboring routers)

Conservation of Flow:

Watcher runs on each router

Statistical Anomaly detection

Protocol Analysis:

Finite State Machine models that register the transition states of the protocol and alarms when anomalous state detected

General problems of IDS:

High Cost due to local management

Failure to exhibit scalability

Fine tuning requirements (based on particular system needs)

Need for frequent Database updates

Passive behavior (no decisions about the actions to be undertaken)

Why Wireless ad-hoc networks is more difficult?

Wireless communication (open media)

Cooperation among nodes is necessary (lack of centralized author.)

Don't rely on existing infrastructure

Have many operational limitations:

Transmission Range and Bandwidth

Energy, CPU, and Memory

Autonomous units capable of roaming independently

easily captured and compromised without physical protection

very expensive and not scalable if physically protected

Why Wireless ad-hoc networks is more difficult?

Usually used in situations where rapid deployment is necessary

Usually deployed in hostile (not physically protected) places

Dynamic topology change (due to mobility)

Lack of key concentration points (e.g. switches and routers)

No firewalls or gateways

Difficult to distribute and update signatures (detection database)

Questions to develop a viable IDS in Ad-hoc networks:

What is the good architecture ?

What are the appropriate audit data sources ?

How to model the activities to distinguish anomaly from normalcy?

Architecture of IDS in Wireless Ad-Hoc Networks:

Stand alone:

- each node has an independent IDS
- no cooperation or sharing between different nodes

Distributed and cooperative IDS:

- each node has an independent IDS (local decisions)
- cooperatively participate in global intrusion detection

Hierarchical IDS:

- for multi-level ad-hoc networks
- cluster heads / control nodes / main stations

IDS for wireless Ad-Hoc Networks:

Watchdog: (e.g. on DSR to verify that the next node forwards the packet)

Control Messages:

- (e.g. adding two control messages to DSR; CREP and CREQ)

Neighborhood watch:

- either by listening to the transmission or observing route protocol behavior

Statistical Anomaly detection

IDS

Example1

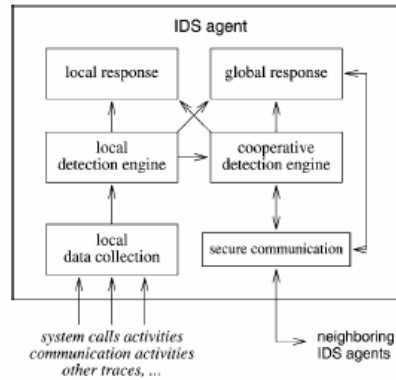
Intrusion Detection
Techniques for mobile Ad
Hoc Wireless Networks

ZHANG, LEE, HUANG

- The architecture
- The Model

Example...

The architecture



Architecture ...

Data collection:

- system and user activities within the mobile node
- communication activities by this node
- communication activities within the radio range

Local detection:

- analyses local data traces
- can use both misuse and anomaly detection
- can initiate anomaly response if it has strong evidence

Architecture ...

Cooperative detection

If the evidence is weak or inconclusive

Initiates a global detection procedure,

This procedure propagates intrusion detection state information among neighboring nodes

The passed state may be:

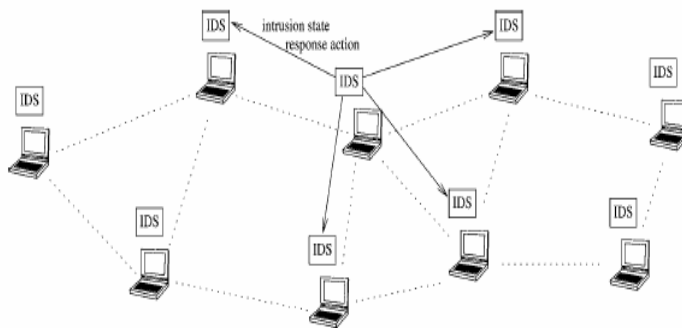
a mere level-of-confidence value $p\%$ that there is anomaly or, a specific state that node A is compromised, or a complicated record including the complete evidence

Decision at each node is build based on majority opinion

Any node that detects anomaly can initiate a response

Architecture ...

cooperative detection



Architecture ...

Intrusion response:

Depends on:

- the type of intrusion,
- the type of network protocols and applications,
- and the confidence of the evidence

Could be (examples)

- reinitializing communication channels
- Identifying the compromised nodes and reorganize the network to exclude them

Anomaly detection model in mobile Ad-hoc networks:

Framework:

Uses entropy and conditional entropy to:

- describe the characteristics of normal and abnormal data flows

Uses classification algorithms to build the model

Details:

- select audit data with low entropy
- transform the data to construct high information gain features
- build classifier using training data
- apply classifier to test data
- post-process alarms to produce intrusion reports

Attack models:

Route logic compromise by manipulating routing information

misrouting

false route update

Traffic pattern distortion

packet dropping

generate packets with fake source address

corruption on packet contents

denial-of-service

Audit data:

Local routing information

Position locator or GPS

Classifier:

RIPPER

applied directly to the feature space

SVM Light (more accurate)

creates new features

Examples of used features:

PCR : the % of changed routes

PCH : the % of changes in the sum of hops of all the routes

VELOCITY: velocity

DISTANCE: distance from last log

RDC : relative distance change.

...

Protocols studied:

DSR

AODV

DSDV

Issues and Discussion:

Efficient Host-Based monitoring requires

 large amount of CPU processing power

 large amount of energy

Efficient Network-Based monitoring requires:

 Traffic bottlenecks

 does not exist in wireless media

 Thus monitoring at every node is required!

 which is not bandwidth nor energy efficient

Anomaly based model

 built on a long-term monitoring

Issues and Discussion ...

Ad-hoc may not have sufficient life time to do so
mobile networks are very dynamic in structure (hard to
build a reliable normal behavior)

Misuse requires

extensive database of attack signatures
must be replicated among all nodes

Another Example on IDS

Effective Intrusion Detection
Using Multiple Sensors in
Ad Hoc Wireless Networks

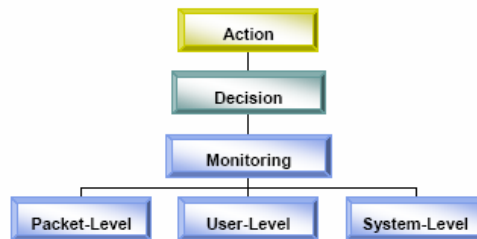
Paper contributions:

Restrict the computation intensive Analysis to a few nodes

A protocol to dynamically select cluster heads

A Selected cluster collect packets within its communication Range

Modular architecture



Modular Architecture:

Network monitoring

only certain nodes have these agents

network packet monitoring

Host monitoring

every node has this agent

monitors the node itself (system and application activities)

Decision Making

every node decides locally on intrusion levels it detect

certain nodes collect intrusion alarms and make collective decision

Action : every node will have an action model

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