The Mystery of the Failing Jobs: Insights from Operational Data from Two University-Wide Computing Systems

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Slide 1/38



Overview

- Introduction
- System and Data Details
- Job Characteristics
- Analyses
 - Job Categories Based on Exit Statuses
 - Effect of Resource Usage on Job Failures
 - Predicting Job Failures and a Better Checkpointing Method
- Open Challenges
- Conclusion





Introduction

- · Job failure leads to resource wastage and user dissatisfaction
- University computing clusters are uniquely challenging:
 - Heterogeneity of jobs: Compute-Intensive, Memory-Intensive, IO-Intensive
 - Varied expertise level of the users
 - Relatively smaller size of the system administration staff
- The most comprehensive dataset publicly analyzed to date in terms of variety of data sources
 - Accounting logs, resource utilization stats, failure reports
 - System-A: Less expensive HW, 617 users, ∼3M jobs
 - System-B: More expensive HW, 467 users, ~2M jobs
- New insights and old insights in new environments
 - Recommendations to reduce job failure/resource wastage for both system user and system admin
 - Build an actionable failure prediction model based on resource usages



Slide 3/38



System and Data Details

- System A
 - 580 nodes, Intel Xeon E5-2670 processors, 64 GB/node, 100 MB/s local IO
 BW, 23 GB/s network IO BW
- System B
 - 26,868 nodes, AMD 6276 Interlagos processors, 64 GB/node, 1.1 TB/s network IO BW
- Data
 - Accounting logs
 - Resource Utilization Stats
 - 5-minute granularity for System A and 1-minute granularity for System B
 - Node Failure Reports



Summary of Data Analyzed

| Computing Cluster | | System A | System B | |
|-------------------|----------|-----------------------|----------------------|--|
| Duration | | Mar 2015-Jun 2017 | Feb-June 2017 | |
| # jobs | | 2,908k | 2,219k | |
| shared | # single | 1,125k (38.7%, 15.8%) | - | |
| | # multi | 28k (1.0%, 1.9%) | - | |
| | total | 1,153k (39.7%, 17.7%) | - | |
| | # single | 1,348k (46.3%, 18.4%) | 1,640k (73.9%, 5.4%) | |
| non-shared | # multi | 407k (14.0%, 63.9%) | 580k (26.1%, 94.6 %) | |
| | total | 1,755k (60.3%, 82.3%) | 2,219k (100%) | |
| # unique users | | 617 | 467 | |

- All production jobs
- Node-seconds = #nodes x execution time
- The percentages in parenthesis refers to the raw counts and node-seconds
- Sharing allows multiple jobs to run on the same node
 - System A: 39.7% by count and 17.7% by node-seconds

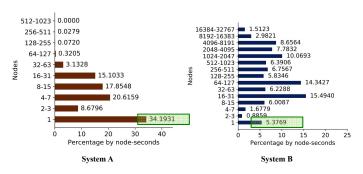


Slide 5/38



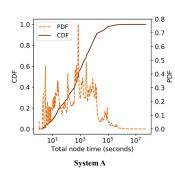
Job Characteristics

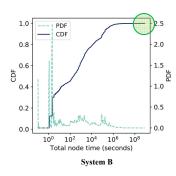
- · Job size
 - Single node jobs by count
 - System A: 85%, System B: 74%
 - Single node jobs by node-seconds
 - System A: 34%, System B: 5%



Job Characteristics

- Job node-seconds
 - System A and System B: 50% of the jobs run for less than ~10³ nodeseconds
 - System B: Jobs run up to $\sim 10^9$ node-seconds







Slide 7/38



Job Categories Based on Exit Statuses

Table: Job categories based on exit codes. Percentages in brackets are based on the total node-seconds

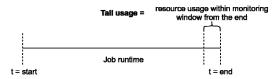
| | | Environment & Job Type | | | | | | | |
|----------|-------------|------------------------|--------|--------|--------|---------------|------------|-------|----------------|
| | | System A | | | | System B | | | |
| | | sha | shared | | hared | overall | non-shared | | overall |
| | | single | multi | single | multi | | single | multi | |
| | Success | 93.1% | 87.6% | 87.6% | 61.8 % | 86.1% (48.4%) | 91.6% | 64.0% | 84.4% (44.4 %) |
| Category | System | 2.7% | 6.5% | 6.5% | 8.8 % | 5.3% (4.0%) | 0.10% | 1.0% | 0.3% (1.4%) |
| eg | User | 1.6% | 2.2% | 3.5% | 7.2% | 3.3%(12.9%) | 3.8% | 3.0% | 3.6% (2.7%) |
| Cat | User/System | 0.6% | 0.2% | 0.4% | 6.1% | 1.3% (1.3%) | 1.2% | 0.8% | 3.6% (8.0%) |
| | Walltime | 2.0% | 3.5% | 2.0% | 16.1% | 4.0% (33.4%) | 3.7% | 20.4% | 8.0% (43.4%) |
| | Total | 1,125k | 28k | 1,348k | 407k | 2,908k | 1,640k | 579k | 2,219k |

- Failure categories System, User, User/System
- System related failures System A: 5.3%, System B: 0.3%
- Success category
 - Multi-node System A: 61.8% (non-shared), System B: 64.0% (non-shared)
 - Single-node System A: 93.1% (shared) vs 87.6% (non-shared), System B: 91.6%
 - · Sharing does not negatively impact the jobs failure probability.
- Walltime category by node-seconds System A: 33.4%, System B: 43.3%



Effect of Resource Usage on Job Failures

- Job failure rate is defined as the fraction of jobs that fail due to system related issues
- All analyses conducted using tail utilization values



- Hypothesis testing for all correlation studies
- · Resource usage prediction models based on user profiling
 - Last: same resource usage as last finished job of a given user
 - Average: average resource usage of last 'n' finished job of a given user
 - Median: median resource usage of last 'n' finished job of a given user
 - Maximum cosine similarity: same resource usage as the most similar job based on cosine similarity



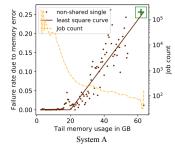
Slide 9/38

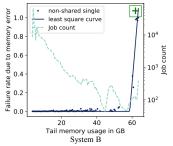


Effect of Resource Usage on Job Failures

- Memory
 - Single-node jobs: +ve correlation
 - Multi-node jobs: no correlation
 - 99th percentile value:
 - System A 11.7 GB
 - System B 45.6 GB

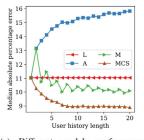
| Jocenn | ood Type | correlation coern, p variet | | |
|-------------|-------------------|-----------------------------|--|--|
| System A | Non-shared single | 0.83, 1.7e-28 | | |
| | Non-shared multi | 0.17, 0.4 | | |
| | Shared single | 0.84, 7.2e-32 | | |
| System B | Non-shared single | 0.57, 3.2e-9 | | |
| | Non-shared multi | 0.13, 0.2 | | |

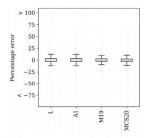




Resource Usage Prediction by User Profiling

- Memory (System A)
 - MAPE of all predictors are less than 12% (for at least one history length)
 - Maximum Cosine Similarity (MCS) outperforms others
 - Use case: Predict memory usage in advance
 - · Better scheduling for heterogeneous memory cluster
 - · Better scheduling when sharing is enabled





- training set
- (a) Different models performance (b) Percentage error distribution for with different history lengths on the different models on test set with best history length as per training set.



Slide 11/38

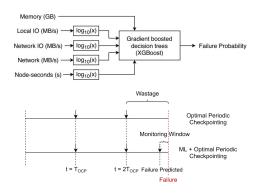


Summary: Effect of Resource Usage on Job Failure

- Random IO access requests lead to failure even at ~1% of BW
 - Local IO (System A)
 - BW of 100MB/s while failure rate starts rising with utilization as low as 3 MB/s (shared) - 6MB/s (non-shared)
 - Remote IO (System B)
 - BW of 1.1TB/s while failures are observed with a utilization of only 46MB/s for
- Contention at remote resources (outside node) dominant in non-shared environment, while the contention at local resources (at node) dominant in shared environment.
 - Use user-based resource usage prediction while making scheduling
 - Use dynamic reconfiguration of applications based on current resource availability, such as reconfiguring the number of threads or network timeout.



Predicting Job Failure



- ML model
 - Input: current resource usages, Output: failure probability within the next monitoring window
- Better checkpointing method
 - Combine our ML model with the optimal periodic checkpointing method



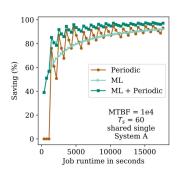
Slide 13/38



A Better Checkpointing System

Normalized area under the curve (normalized with respect to jobs with no wastage execution due to failures).

| | System | | Periodic | ML | ML+Periodic |
|-------------------------|--------|-------------------|----------|------|-------------|
| | | shared single | | 0.82 | 0.91 |
| MTBF=1e4, | A | non-shared single | 0.81 | 0.89 | 0.94 |
| T_S =60 sec | | non-shared multi | | 0.90 | 0.95 |
| | В | non-shared multi | | 0.91 | 0.94 |
| | | shared single | 0.66 | 0.82 | 0.88 |
| MTBF=1e5, T_S =60 sec | A | non-shared single | | 0.89 | 0.92 |
| | | non-shared multi | | 0.90 | 0.93 |
| | В | non-shared multi | | 0.91 | 0.93 |
| | A | shared single | 0.30 | 0.82 | 0.84 |
| MTBF=1e6, | | non-shared single | | 0.89 | 0.90 |
| T_S =60 sec | | non-shared multi | | 0.90 | 0.91 |
| | В | non-shared multi | | 0.91 | 0.92 |
| MTBF=1e6, T_S =10 sec | A | shared single | | 0.95 | 0.95 |
| | | non-shared single | 0.60 | 0.97 | 0.97 |
| | | non-shared multi | | 0.97 | 0.98 |
| | В | non-shared multi | | 0.97 | 0.98 |



- ML + periodic checkpointing method outperforms the base optimal checkpointing method by between 12.3% (unreliable system with MTBF =1e4, Ts=60s) and 2X (reliable system with MTBF = 1e6 and Ts=60s).
- Savings achieved by the optimal checkpointing method in case of failure decreases as a system becomes more reliable i.e., as the MTBF increases from 1*e*4 to 1e6.



Open Challenges

- Current optimal checkpointing estimation methods take only hardware reliability (such as MTBF) into account
 - This paper integrated it with job failure likelihood information
 - A better method is to consider in addition the rate of job progress
- Current contention-aware schedulers need to profile a job first to estimate job's interference and latency-sensitivity
 - Major limitation for clusters where majority of jobs are short running
 - Use user history-based resource usage predictions to profile a job profile



Slide 15/38



Conclusion

- The most comprehensive dataset publicly analyzed to date in terms of variety of data sources
- Publicly released the dataset on which the analyses are based
- Important insights into how the clusters behave and implications for how they can be managed more effectively.



Thank You!



Slide 17/38