Responses to Cyber Attacks in Distributed Systems

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Slide 1/27



Outline

- Problem Statement
- Solution Directions
- Some Promising Solutions
- Ongoing Challenges





Defending Distributed Systems



- Large-scale distributed systems to defend
 - Heterogeneous third-party services
- Lots of points for attacks
 - Lots of points to introduce cybersecurity mechanisms
- Interactions between the services allow for attack escalation

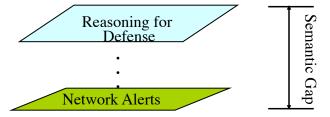


Slide 3/27



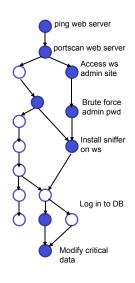
Drowning in a Sea of Alerts

- Large distributed systems get tons of alerts
 - Up to 20,000 per day
- Many of these are false alarms



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- Multi-stage attacks
 - Compromise outward facing services
 - Use transitive trust and privilege escalation
 - Compromise internal services
 - Access crown jewels
- Attack progresses in machine time, rather than human time
- Examples: Worms and other self-propagating malware

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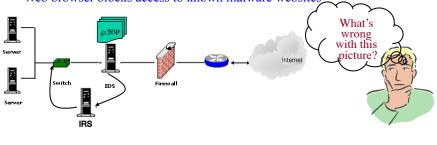
Slide 5/27

PURDUE

Signature-based Responses

- Intrusion Response Systems (IRS) take reports from IDS and carry out actions to counter the intrusion
- Many examples of IRS
 - Anti-virus software disables access to worm executables or files infected with virus
 - Iptables which terminates a session on matching a malware signature

Web browser blocks access to known malware websites



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Dealing with Zero-Day Attacks

- Zero-day attacks are difficult to deal with through signature-based mechanisms
 - They exploit unknown vulnerabilities
 - Their path of attack spread is not known a priori
- Challenges for zero-day attacks
 - Exact matching of mechanics of attack step does not work
 - A reactive approach to security allows devastating zero-day attacks to get through
 - Learning-based approaches are predicated on exact matches and therefore do not work well



Slide 7/27



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Solution Directions

- We want to perform secure configuration and intrusion response in the face of threats that are fast-changing and therefore unknown
- 1. We want to learn from past behavior
 - But not overlearn
- 2. We want to grow our knowledge structures with runtime information
 - But not learn untruths
- 3. We want to perform the learning at runtime
 - This implies expensive batch mode processing is out
- 4. We do not want to rely only on signature-based security
 - Abstractions of attack steps are useful



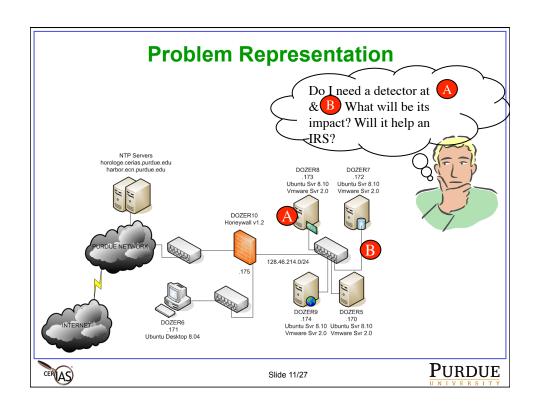
Slide 9/27

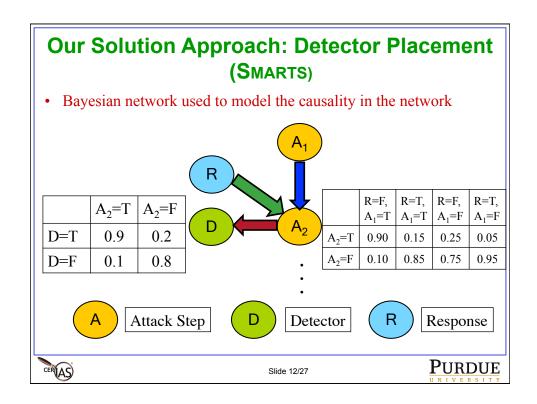


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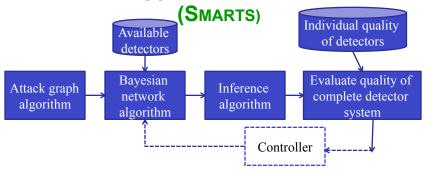
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Our Solution Approach: Detector Placement



- Inference on the Bayesian network performed through different choice and placements of detectors
- Heuristic-driven choice of one detector and its placement at a time
- Heuristic depends on individual detector quality and overlap with previously chosen detectors
- Controller to adjust detector setting when network changes



Slide 13/27



Adaptive to Current Threat Environment

- It is expensive to turn all sensor rules all the time
 - Example: Snort default rule set has > 9,000 single step attack rules, in 73 categories and takes > 5 sec to match all of them
- Approach:
 - Perform damage assessment currently through Bayesian inferencing
 - Damage assessment indicates
 - Which components are likely compromised but needs further evidence to determine with high confidence
 - Based on attack spread, which components are likely to be compromised
 - Sensor rules are activated based on results of damage assessment
- Responsive to changes in system
 - Incremental inferencing when some parts of system change



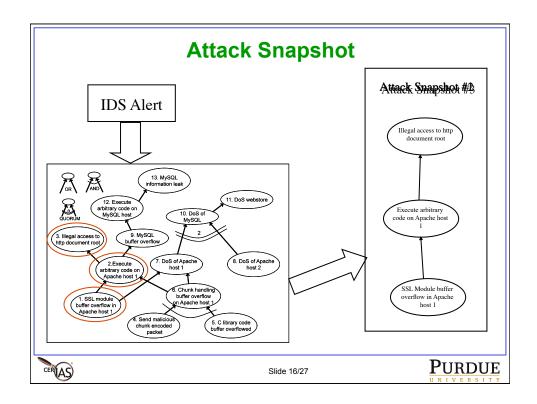
Our Solution Approach: Intrusion Response (ADEPTS)

- Short-term as well as long-term goals
 - Contain the current attack
 - Recover affected services to a functional state
 - Proactive defenses for future attacks
- Leverage distributed system's characteristics
 - Determine if the alert is false
 - Determine if the impact is worth responding to
- Learn from thy observations and mistakes
 - Calibrate prior responses
 - Learn characteristics of interactions in the system through past attacks
 - Quick customized responses to polymorphs of prior attacks



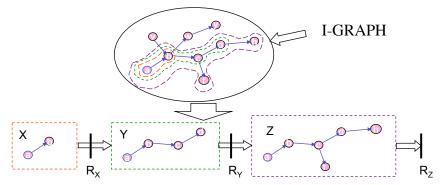
Slide 15/27





Dynamics between attack and responses

• Successive attack snapshots created for incoming IDS alerts



- · Assuming an attack includes three "snapshots" X, Y, and Z
- Each snapshot includes I-GRAPH nodes which have been achieved as part of the attack thus far
- Following each snapshot k, SWIFT determines a response combination R_k (a set of response actions) to deter the escalation

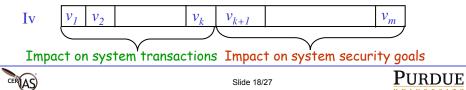


Slide 17/27



Impact Vector

- A system has transaction goals and security goals that it needs to meet through the time of operation
 - Example: provide authentication service & preserve privacy of sensitive data
- Attacks are meant to impact some of these goals
- Deployed responses also impact some of these goals
 - For example, by temporarily disabling some functionality for legitimate users as well
- Assume the impact can be quantified through a vector Iv
 - − Each element in the Iv corresponds to the impact on each transaction/ security goal $\in [0, \infty]$



Optimality of Response Actions

• We formally define the cost for a response combination (a set of response actions) RC_i as:

CostRCOMPT SS

 $Iv(n_k)$: Impact from reaching an attack step node n_k

 $Pr(n_k)$: Probability of reaching node n_k

 $Iv(r_k)$: Impact from deploying the response r_k

- The response combination RC_i is said to be optimal for a given attack if it achieves the minimal Cost(RC_i)
 - In ADEPTS, optimality achieved "per node and per out-going edge"

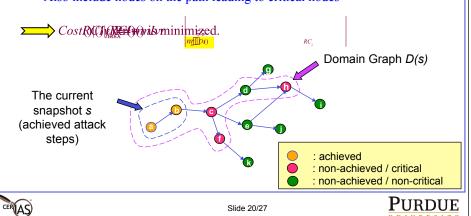


Slide 19/27



Domain Graph

- Limit the response search space for a snapshot s to a subset of I-GRAPH, namely the **Domain Graph** D(s)
- D(s) includes critical nodes from I-GRAPH
 - A node n is critical if |Prob(n)*Iv(n)| is greater than a given threshold
 - Also include nodes on the path leading to critical nodes



Responding to the Unknown

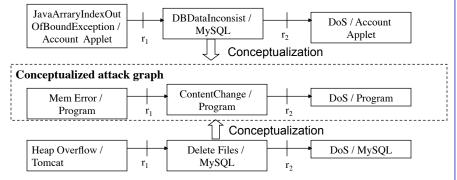
- Zero-day attack
 - Knowledge of the steps in the attack does not exist in the IRS
- Current solution: Take a drastic response, such as disconnecting the service
- Problem:
 - May be reacting to spurious alarms
 - Cannot learn from the spread of the attack
- Our solution approach:
 - Abstract the specifics of the attack
 - At a higher level of abstraction, map the attack to a previously seen attack
 - Use the learning on the previous attack to guide the responses for the current zero-day attack



Slide 21/27



Responding to the Unknown: Example



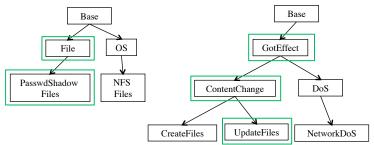
Responses: r_1 : Disable connection from tomcat/applet to MySQL; r_2 : Rollback to last data files checkpoint

• Challenges: (1) High similarity does not necessarily give you the best response; (2) To what level should each node be conceptualized



Utilize History from Similar Attack

- How to calculate similarity between two attacks?
 - Inheritance hierarchy for components, detector alerts, and connections



- Calculate distance for each node and each connection
- Compute graph edit distance
 - Conceptually, the sequence of steps to convert one graph to another
 - Through addition, deletion, or modification of nodes and connections



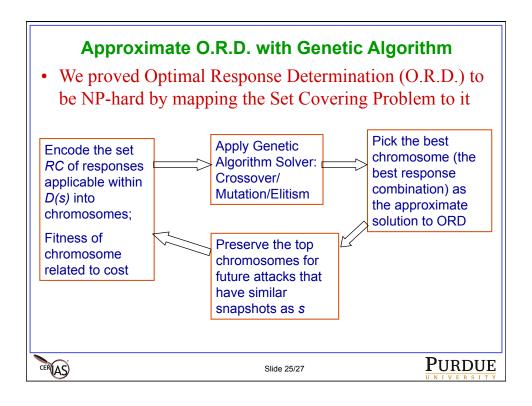
Slide 23/27

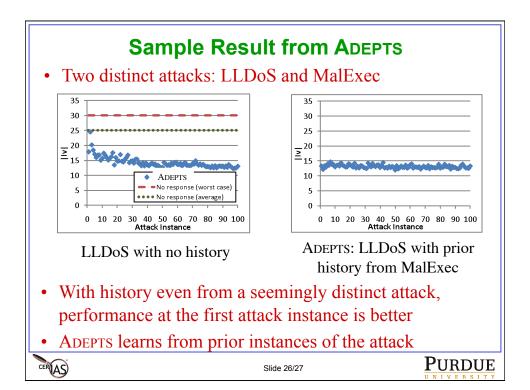


Utilize History from Similar Attack

- Acquire from the similar attack
 - Effectiveness Index (EI) values of responses
 - Edge Propagation Factor (EPF) values of edges
 - Effective Response Combinations
- Efficient search through space of prior attacks
 - Attack similarity is defined to follow metric space conditions: d(x, x) = 0; d(x, y) = d(y, x); $d(x, y) + d(y, z) \ge d(x, z)$
 - Prior work allows for efficient storage and search through attack template library
 - Disjoint parts of multiple attacks can be used in responding to the current attack







Goals of Ongoing Work

- Secure Configuration Management
 - Detector placement is a specific example of security configuration
 - Tool should detect (when insecure configuration is introduced) and diagnose (which component has been mis-configured)
 - Tradeoffs exist between security of configuration and usability
 - Tool must not make arbitrary decisions on this spectrum
- Automated Intrusion Response
 - Resilience to zero-day attacks through more effective responses (i.e., less drastic than rebooting the servers)
 - Correlation of multiple detectors to increase confidence that an attack is underway before responding

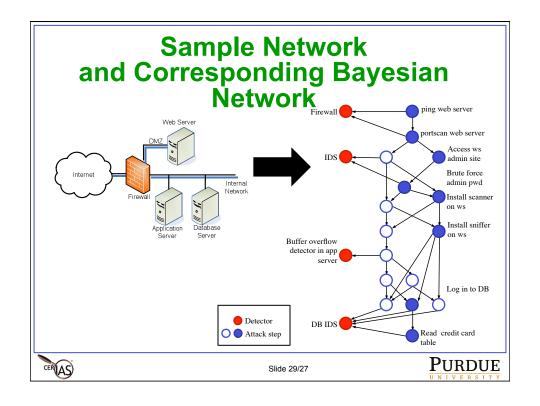


Slide 27/27



Backup Slides





Impact on Choice and Placement of Detectors • System: Three-tier webbased online service 0.8 • Objective: determine impact True Positive Rate of selecting detectors and corresponding locations $(d_{20}, x_{19}), (d_{20}, x_{16})$ Performance of detector pair 0.2 $(d_{13},x_{12}),(d_7,x_6)$ (selected from algorithm) is compared against randomly 0.2 0.4 0.6 False Positive Rate selected pairs Selection made by our algorithm **PURDUE** CER (AS) Slide 30/27