

ECE595 / STAT598: Machine Learning I

Lecture 20 Support Vector Machine: Duality

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Outline

Support Vector Machine

- Lecture 19 SVM 1: The Concept of Max-Margin
- **Lecture 20 SVM 2: Dual SVM**
- Lecture 21 SVM 3: Kernel SVM

This lecture: Support Vector Machine: Duality

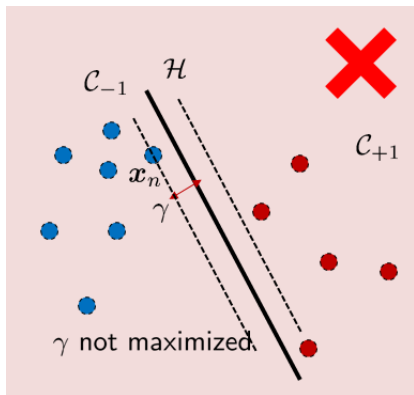
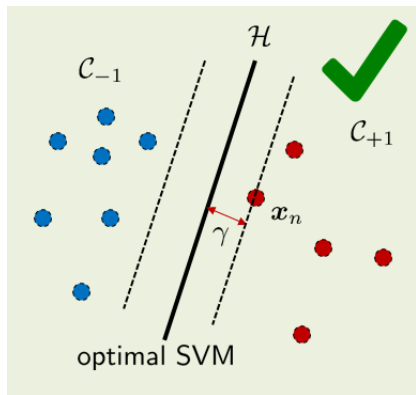
- Lagrange Duality
 - Maximize the dual variable
 - Minimax Problem
 - Toy Example
- Dual SVM
 - Formulation
 - Interpretation

Support Vector Machine

SVM is the solution of this optimization

$$\underset{\mathbf{w}, w_0}{\text{minimize}} \quad \frac{1}{2} \|\mathbf{w}\|_2^2$$

$$\text{subject to } y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \geq 1, \quad j = 1, \dots, N.$$



Lagrange Function

- **Goal:** Construct the dual problem of

$$\begin{aligned} & \underset{\mathbf{w}, w_0}{\text{minimize}} \quad \frac{1}{2} \|\mathbf{w}\|_2^2 \\ & \text{subject to} \quad y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \geq 1, \quad j = 1, \dots, N. \end{aligned}$$

- **Approach:** Consider the Lagrangian function

$$\mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) \stackrel{\text{def}}{=} \underbrace{\frac{1}{2} \|\mathbf{w}\|_2^2}_{\text{objective}} + \sum_{j=1}^N \lambda_j \underbrace{\left[1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0)\right]}_{\text{constraint}},$$

- Solution happens at the saddle point of \mathcal{L} :

$$\nabla_{(\mathbf{w}, w_0)} \mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) = \mathbf{0}, \quad \text{and} \quad \nabla_{\boldsymbol{\lambda}} \mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) = \mathbf{0}.$$

Lagrangian function

- The **Lagrangian function** is

$$\mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) \stackrel{\text{def}}{=} \frac{1}{2} \|\mathbf{w}\|_2^2 + \sum_{j=1}^N \lambda_j \underbrace{\left[1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \right]}_{\leq 0}$$

- Complementarity Condition** says

- $\lambda_j > 0$ and $[1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0)] = 0$
 - $\lambda_j = 0$ and $[1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0)] < 0$
- So, if we want $\nabla_{\boldsymbol{\lambda}} \mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) = \mathbf{0}$, then must be one of the two cases:

$$\sum_{j=1}^N \lambda_j [1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0)] \rightarrow \max \quad \text{or} \quad \min$$

- No saddle point because linear in λ .
- But $1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \leq 0$. So unbounded minimum. So must go with maximum.

Primal Problem

- Let λ^* as the maximizer

$$\lambda^* \stackrel{\text{def}}{=} \operatorname{argmax}_{\lambda \geq 0} \left\{ \sum_{j=1}^N \lambda_j \left[1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \right] \right\}$$

- Then the **primal problem** is

$$\underset{\mathbf{w}, w_0}{\text{minimize}} \quad \mathcal{L}(\mathbf{w}, w_0, \lambda^*)$$

$$= \underset{\mathbf{w}, w_0}{\text{minimize}} \left\{ \frac{1}{2} \|\mathbf{w}\|_2^2 + \max_{\lambda \geq 0} \left\{ \sum_{j=1}^N \lambda_j \left[1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \right] \right\} \right\}$$

$$= \underset{\mathbf{w}, w_0}{\text{minimize}} \left\{ \max_{\lambda \geq 0} \mathcal{L}(\mathbf{w}, w_0, \lambda) \right\}$$

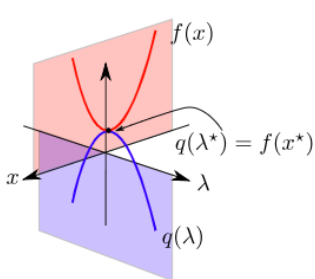
- This is a min-max problem:

$$\min_{\mathbf{w}, w_0} \max_{\lambda \geq 0} \mathcal{L}(\mathbf{w}, w_0, \lambda)$$

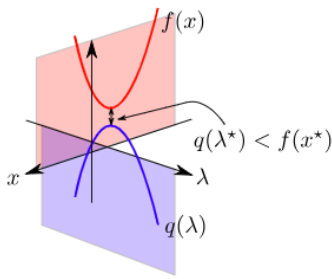
Strong Duality

- Recall that our problem is **quadratic programming** (QP).
- Strong Duality** holds for QP:

$$\underbrace{\min_{\mathbf{w}, w_0} \max_{\lambda \geq 0} \mathcal{L}(\mathbf{w}, w_0, \lambda)}_{\text{primal}} = \underbrace{\max_{\lambda \geq 0} \min_{\mathbf{w}, w_0} \mathcal{L}(\mathbf{w}, w_0, \lambda)}_{\text{dual}}$$



strong duality



weak duality

Toy Example

- The SVM problem

$$\begin{aligned} & \underset{\mathbf{w}, w_0}{\text{minimize}} && \frac{1}{2} \|\mathbf{w}\|_2^2 \\ & \text{subject to} && y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \geq 1, \quad j = 1, \dots, N. \end{aligned}$$

is in the form of

$$\underset{\mathbf{u}}{\text{minimize}} \quad \|\mathbf{u}\|^2, \quad \text{subject to} \quad \mathbf{a}_j^T \mathbf{u} \geq b_j, \quad j = 1, 2, \dots, N$$

- Example:**

$$\underset{u_1, u_2}{\text{minimize}} \quad u_1^2 + u_2^2, \quad \text{subject to} \quad \begin{bmatrix} 1 & 2 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} \geq \begin{bmatrix} 2 \\ 0 \\ 0 \end{bmatrix}$$

- Can we write down its dual problem?

Toy Example

- Lagrangian function is

$$\mathcal{L}(\mathbf{u}, \boldsymbol{\lambda}) \stackrel{\text{def}}{=} u_1^2 + u_2^2 + \lambda_1(2 - u_1 - 2u_2) - \lambda_2 u_1 - \lambda_3 u_3$$

- Minimize over \mathbf{u} :

$$\begin{aligned}\frac{\partial \mathcal{L}}{\partial u_1} = 0 &\Rightarrow u_1 = \frac{\lambda_1 + \lambda_2}{2} \\ \frac{\partial \mathcal{L}}{\partial u_2} = 0 &\Rightarrow u_2 = \frac{2\lambda_1 + \lambda_3}{2}.\end{aligned}$$

- Plugging into the Lagrangian function yields

$$\begin{aligned}\text{maximize}_{\boldsymbol{\lambda}} \quad & \mathcal{L}(\boldsymbol{\lambda}) = -\frac{5}{4}\lambda_1^2 - \frac{1}{4}\lambda_2^2 - \frac{1}{4}\lambda_3^2 - \frac{1}{2}\lambda_1\lambda_2 - \lambda_1\lambda_3 + 2\lambda_1 \\ \text{subject to} \quad & \boldsymbol{\lambda} \geq 0.\end{aligned}$$

- Primal is QP. Dual is also QP.

Have We Gained Anything?

- Here is the dual problem:

$$\begin{aligned} \underset{\lambda}{\text{maximize}} \quad \mathcal{L}(\lambda) &= -\frac{5}{4}\lambda_1^2 - \frac{1}{4}\lambda_2^2 - \frac{1}{4}\lambda_3^2 - \frac{1}{2}\lambda_1\lambda_2 - \lambda_1\lambda_3 + 2\lambda_1 \\ \text{subject to} \quad \lambda &\geq 0. \end{aligned}$$

- These terms are all **negative**! So we must have $\lambda_2 = \lambda_3 = 0$.
- This gives

$$\underset{\lambda_1 \geq 0}{\text{maximize}} \quad -\frac{5}{4}\lambda_1^2 + 2\lambda_1.$$

which is maximized at $\lambda_1 = \frac{4}{5}$.

- Plugging into the primal yields

$$u_1 = \frac{\lambda_1 + \lambda_2}{2} = \frac{2}{5}, \quad \text{and} \quad u_2 = \frac{2\lambda_1 + \lambda_3}{2} = \frac{4}{5}.$$

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Dual of SVM

- We want to find the dual problem of

$$\begin{aligned} & \underset{\mathbf{w}, w_0}{\text{minimize}} \quad \frac{1}{2} \|\mathbf{w}\|_2^2 \\ & \text{subject to} \quad y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \geq 1, \quad j = 1, \dots, N. \end{aligned}$$

- We start with the Lagrangian function

$$\mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) \stackrel{\text{def}}{=} \frac{1}{2} \|\mathbf{w}\|_2^2 + \sum_{j=1}^N \lambda_j \left[1 - y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \right].$$

- Let us **minimize** over (\mathbf{w}, w_0) :

$$\nabla_{\mathbf{w}} \mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) = \mathbf{w} - \sum_{j=1}^N \lambda_j y_j \mathbf{x}_j = \mathbf{0}$$

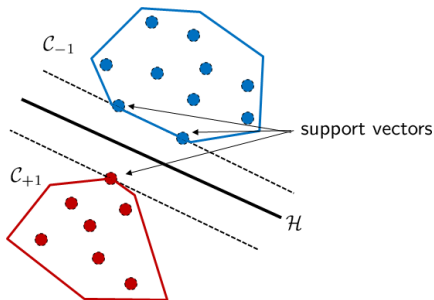
$$\nabla_{w_0} \mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) = \sum_{j=1}^N \lambda_j y_j = 0.$$

Interpreting $\nabla_{\mathbf{w}} \mathcal{L}(\mathbf{w}, \mathbf{w}_0, \boldsymbol{\lambda}) = 0$

- The first result suggests that

$$\mathbf{w} = \sum_{j=1}^N \lambda_j y_j \mathbf{x}_j.$$

- This is support vector: λ_j is either $\lambda_j = 0$ or $\lambda_j > 0$.



Interpreting $\nabla_{\mathbf{w}} \mathcal{L}(\mathbf{w}, w_0, \boldsymbol{\lambda}) = 0$

- The complementarity condition states that

$$\lambda_j^* \left[1 - y_j(\mathbf{w}^{*T} \mathbf{x}_j + w_0^*) \right] = 0, \quad \text{for } j = 1, \dots, N.$$

- If $1 - y_j(\mathbf{w}^{*T} \mathbf{x}_j + w_0^*) > 0$, then $\lambda_j^* = 0$
- If $\lambda_j^* > 0$, then $1 - y_j(\mathbf{w}^{*T} \mathbf{x}_j + w_0^*) = 0$
- So you can define the **support vector set**:

$$\mathcal{V} \stackrel{\text{def}}{=} \{j \mid \lambda_j^* > 0\}.$$

- So the optimal weight is

$$\mathbf{w}^* = \sum_{j \in \mathcal{V}} \lambda_j^* y_j \mathbf{x}_j.$$

Back to Duality

The Lagrangian function is

$$\begin{aligned}\mathcal{L}(\mathbf{w}^*, w_0^*, \boldsymbol{\lambda}) &= \frac{1}{2} \|\mathbf{w}^*\|_2^2 + \sum_{j=1}^N \lambda_j \left[1 - y_j ((\mathbf{w}^*)^T \mathbf{x}_j + w_0) \right] \\ &= \underbrace{\frac{1}{2} \left\| \sum_{j=1}^N \lambda_j y_j \mathbf{x}_j \right\|_2^2}_A \\ &\quad + \underbrace{\sum_{j=1}^N \lambda_j \left[1 - y_j \left(\left(\sum_{i=1}^n \lambda_i y_i \mathbf{x}_i \right)^T \mathbf{x}_j + w_0 \right) \right]}_B\end{aligned}$$

Back to Duality

- We can show that

$$A = \frac{1}{2} \sum_{i=1}^N \sum_{j=1}^N \lambda_i \lambda_j y_i y_j \mathbf{x}_i^T \mathbf{x}_j$$

$$B = \sum_{j=1}^N \lambda_j - \sum_{i=1}^N \sum_{j=1}^N \lambda_i \lambda_j y_i y_j \mathbf{x}_i^T \mathbf{x}_j - \underbrace{\left(\sum_{j=1}^N \lambda_j y_j \right)}_{=0} w_0$$

- and we can show that

$$A + B = \sum_{j=1}^N \lambda_j + \frac{1}{2} \sum_{i=1}^N \sum_{j=1}^N \lambda_i \lambda_j y_i y_j \mathbf{x}_i^T \mathbf{x}_j$$

Back to Duality

- Therefore, the dual problem is

$$\begin{aligned} & \underset{\lambda \geq 0}{\text{maximize}} && -\frac{1}{2} \sum_{i=1}^N \sum_{j=1}^N \lambda_i \lambda_j y_i y_j \mathbf{x}_i^T \mathbf{x}_j + \sum_{j=1}^N \lambda_j \\ & \text{subject to} && \sum_{j=1}^N \lambda_j y_j = 0. \end{aligned}$$

- If you prefer matrix-vector:

$$\begin{aligned} & \underset{\lambda \geq 0}{\text{maximize}} && -\frac{1}{2} \boldsymbol{\lambda}^T \mathbf{Q} \boldsymbol{\lambda} + \mathbf{1}^T \boldsymbol{\lambda} \\ & \text{subject to} && \mathbf{y}^T \boldsymbol{\lambda} = 0. \end{aligned}$$

- We can combine the constraints $\boldsymbol{\lambda} \geq 0$ and $\mathbf{y}^T \boldsymbol{\lambda} = 0$ as

$$\mathbf{A} \boldsymbol{\lambda} \geq \mathbf{0}.$$

Back to Duality

- $\mathbf{y}^T \boldsymbol{\lambda} = 0$ means

$$\mathbf{y}^T \boldsymbol{\lambda} \geq 0, \quad \text{and} \quad \mathbf{y}^T \boldsymbol{\lambda} \leq 0.$$

- Thus, we can write $\mathbf{y}^T \boldsymbol{\lambda} = 0$ as

$$\begin{bmatrix} \mathbf{y}^T \\ -\mathbf{y}^T \end{bmatrix} \boldsymbol{\lambda} \geq \begin{bmatrix} 0 \\ 0 \end{bmatrix}.$$

- Therefore, the matrices \mathbf{Q} and \mathbf{A} are

$$\mathbf{Q} = \begin{bmatrix} y_1 y_1 \mathbf{x}_1^T \mathbf{x}_1 & \cdots & y_1 y_N \mathbf{x}_1^T \mathbf{x}_N \\ y_2 y_1 \mathbf{x}_2^T \mathbf{x}_1 & \cdots & y_2 y_N \mathbf{x}_2^T \mathbf{x}_N \\ \vdots & \vdots & \vdots \\ y_N y_1 \mathbf{x}_N^T \mathbf{x}_1 & \cdots & y_N y_N \mathbf{x}_N^T \mathbf{x}_N \end{bmatrix} \quad \text{and} \quad \mathbf{A} = \begin{bmatrix} \mathbf{y}^T \\ -\mathbf{y}^T \\ \mathbf{I} \end{bmatrix}$$

So How to Solve the SVM Problem?

- You look at the dual problem

$$\begin{aligned} & \underset{\lambda}{\text{maximize}} && -\frac{1}{2}\lambda^T \mathbf{Q}\lambda + \mathbf{1}^T \lambda \\ & \text{subject to} && \mathbf{A}\lambda \geq \mathbf{0}. \end{aligned}$$

- You get the solution λ^* .
- Then compute \mathbf{w}^* :

$$\mathbf{w}^* = \sum_{j \in \mathcal{V}} \lambda_j^* y_j \mathbf{x}_j.$$

- \mathcal{V} is the set of support vectors: $\lambda_j > 0$.

Are We Done Yet?

- Not quite.
- We still need to find out w_0^* .
- Pick any **support vector** $\mathbf{x}^+ \in \mathcal{C}_+$ and $\mathbf{x}^- \in \mathcal{C}_-$.
- Then we have

$$\mathbf{w}^T \mathbf{x}^+ + w_0 = +1, \quad \text{and} \quad \mathbf{w}^T \mathbf{x}^- + w_0 = -1.$$

- Sum them, we have $\mathbf{w}^T (\mathbf{x}^+ + \mathbf{x}^-) + 2w_0 = 0$, which means

$$w_0^* = -\frac{(\mathbf{x}^+ + \mathbf{x}^-)^T \mathbf{w}^*}{2}$$

Summary of Dual SVM

- Primal

$$\begin{aligned} & \underset{\mathbf{w}, w_0}{\text{minimize}} \quad \frac{1}{2} \|\mathbf{w}\|_2^2 \\ & \text{subject to} \quad y_j(\mathbf{w}^T \mathbf{x}_j + w_0) \geq 1, \quad j = 1, \dots, N. \end{aligned}$$

- Strong Duality

$$\underbrace{\min_{\mathbf{w}, w_0} \max_{\lambda \geq 0} \mathcal{L}(\mathbf{w}, w_0, \lambda)}_{\text{primal}} = \underbrace{\max_{\lambda \geq 0} \min_{\mathbf{w}, w_0} \mathcal{L}(\mathbf{w}, w_0, \lambda)}_{\text{dual}}$$

- Dual

$$\begin{aligned} & \underset{\lambda \geq 0}{\text{maximize}} \quad -\frac{1}{2} \sum_{i=1}^N \sum_{j=1}^N \lambda_i \lambda_j y_i y_j \mathbf{x}_i^T \mathbf{x}_j + \sum_{j=1}^N \lambda_j \\ & \text{subject to} \quad \sum_{j=1}^N \lambda_j y_j = 0. \end{aligned}$$

Summary of Dual SVM

- The weights are computed as

$$\mathbf{w}^* = \sum_{j=1}^N \lambda_j^* y_j \mathbf{x}_j.$$

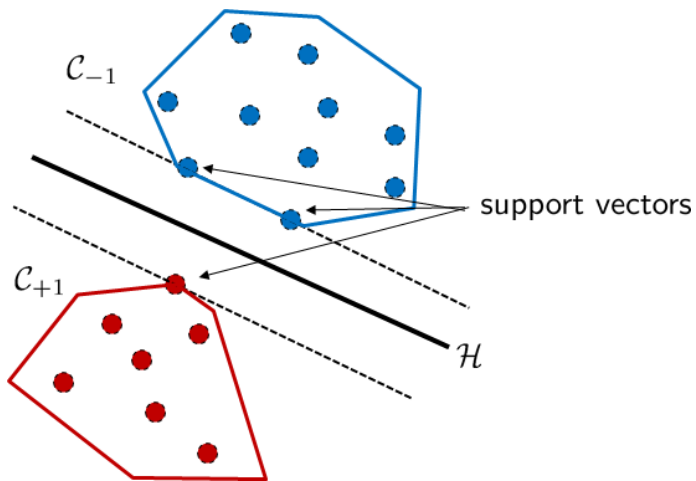
- This is support vector: λ_j is either $\lambda_j = 0$ or $\lambda_j > 0$.
- Pick any **support vector** $\mathbf{x}^+ \in \mathcal{C}_+$ and $\mathbf{x}^- \in \mathcal{C}_-$.
- Then we have

$$\mathbf{w}^T \mathbf{x}^+ + w_0 = +1, \quad \text{and} \quad \mathbf{w}^T \mathbf{x}^- + w_0 = -1.$$

- Sum them, we have $\mathbf{w}^T (\mathbf{x}^+ + \mathbf{x}^-) + 2w_0 = 0$, which means

$$w_0^* = -\frac{(\mathbf{x}^+ + \mathbf{x}^-)^T \mathbf{w}^*}{2}$$

Summary of Dual SVM



Support Vector Machine

- Mustafa, *Learning from Data*, e-Chapter
- Duda-Hart-Stork, *Pattern Classification*, Chapter 5.5
- Chris Bishop, *Pattern Recognition*, Chapter 7.1
- UCSD Statistical Learning
<http://www.svcl.ucsd.edu/courses/ece271B-F09/>

Appendix

Inequality Constrained Optimization

Inequality constrained optimization:

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && f(\mathbf{x}) \\ & \text{subject to} && g_i(\mathbf{x}) \geq 0, \quad i = 1, \dots, m \\ & && h_j(\mathbf{x}) = 0, \quad j = 1, \dots, k. \end{aligned}$$

Requires a function: Lagrangian function

$$\mathcal{L}(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\nu}) \stackrel{\text{def}}{=} f(\mathbf{x}) - \sum_{i=1}^m \mu_i g_i(\mathbf{x}) - \sum_{j=1}^k \nu_j h_j(\mathbf{x}).$$

$\boldsymbol{\mu} \in \mathbb{R}^m$ and $\boldsymbol{\nu} \in \mathbb{R}^k$ are called the **Lagrange multipliers** or the **dual variables**.

Karush-Kahn-Tucker Conditions

If $(\mathbf{x}^*, \boldsymbol{\mu}^*, \boldsymbol{\nu}^*)$ is the solution to the constrained optimization, then all the following conditions should hold:

(i) $\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}^*, \boldsymbol{\mu}^*, \boldsymbol{\nu}^*) = \mathbf{0}$.

- **Stationarity.**

- The primal variables should be stationary.

(ii) $g_i(\mathbf{x}^*) \geq 0$ and $h_j(\mathbf{x}^*) = 0$ for all i and j .

- **Primal Feasibility.**

- Ensures that constraints are satisfied.

(iii) $\mu_i^* \geq 0$ for all i and j .

- **Dual Feasibility.**

- Require $\mu_i^* \geq 0$; but no constraint on ν_i^* .

(iv) $\mu_i^* g_i(\mathbf{x}^*) = 0$ for all i and j .

- **Complementary Slackness**

- Either $\mu_i^* = 0$ or $g_i(\mathbf{x}^*) = 0$ (or both).

KKT Condition is a first order **necessary** condition.

Example: ℓ_2 -minimization with two constraints

Solve the following least squares over positive quadrant problem.

$$\begin{aligned} & \underset{\mathbf{x} \in \mathbb{R}^n}{\text{minimize}} && \frac{1}{2} \|\mathbf{x} - \mathbf{b}\|^2, \\ & \text{subject to} && \mathbf{x}^T \mathbf{1} = 1, \quad \text{and} \quad \mathbf{x} \geq \mathbf{0}. \end{aligned} \tag{1}$$

```
%MATLAB code: Use CVX to solve min ||x-b|| s.t. sum(x) = 1, x >= 0.
cvx_begin
    variable x(n)
    minimize( norm(x-b, 2) )
    subject to
        sum(x) == 1;
        x      >= 0;
cvx_end
```

Analytic Solution

$$\mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}, \gamma) = \frac{1}{2} \|\mathbf{x} - \mathbf{y}\|^2 - \boldsymbol{\lambda}^T \mathbf{x} - \gamma(1 - \mathbf{x}^T \mathbf{1}).$$

Stationarity suggests that:

$$\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}, \gamma) = \mathbf{x} - \mathbf{b} - \boldsymbol{\lambda} + \gamma \mathbf{1} = \mathbf{0}$$

This means

$$\mathbf{x} = \mathbf{b} + \boldsymbol{\lambda} - \gamma \mathbf{1} \quad \text{or} \quad x_i = b_i + \lambda_i - \gamma.$$

The **complementary slackness** implies $\lambda_i x_i = 0$.

- Case 1: If $\lambda_i = 0$, then
 - $x_i = b_i + \overset{0}{\cancel{\lambda_i}} - \gamma = b_i - \gamma$.
 - Since constraint requires $x_i \geq 0$, this means $b_i \geq \gamma$.
- Case 2: If $\lambda_i > 0$, then $x_i = 0$.
 - $\overset{0}{\cancel{x_i}} = b_i + \lambda_i - \gamma$.
 - This implies $b_i + \lambda_i = \gamma$.
 - Since $\lambda_i > 0$, this implies $b_i < \gamma$.

These three cases can be re-written as:

- If $b_i > \gamma$, then $x_i = b_i - \gamma$;
- If $b_i = \gamma$, then $x_i = 0$;
- If $b_i < \gamma$, then $x_i = 0$.

Compactly written as

$$\mathbf{x} = \max(\mathbf{b} - \gamma \mathbf{1}, 0).$$

Primal feasibility implies that

$$\mathbf{x}^T \mathbf{1} = 1, \quad \Leftrightarrow \quad \sum_{i=1}^n x_i = 1.$$

Therefore, γ needs to satisfy the equation

$$\sum_{i=1}^n \max(b_i - \gamma, 0) = 1,$$

which can be found by doing a root-finding of

$$g(\gamma) = \sum_{i=1}^n \max(b_i - \gamma, 0) - 1.$$

Non-CVX Implementation

```
%MATLAB code: solve min ||x-b|| s.t. sum(x) = 1, x >= 0.  
n = 10;  
b = randn(n,1);  
fun = @(gamma) myfun(gamma,b);  
gamma = fzero(fun,0);  
x = max(b-gamma,0);
```

where the function myfun is defined as

```
function y = myfun(gamma,b)  
y = sum(max(b-gamma,0))-1;
```