Cybersecurity and Quantum Computation in Control of Cyberphysical Systems for Next-Generation Manufacturing

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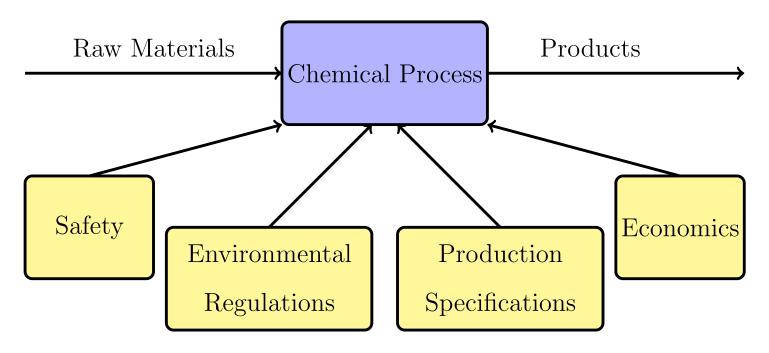


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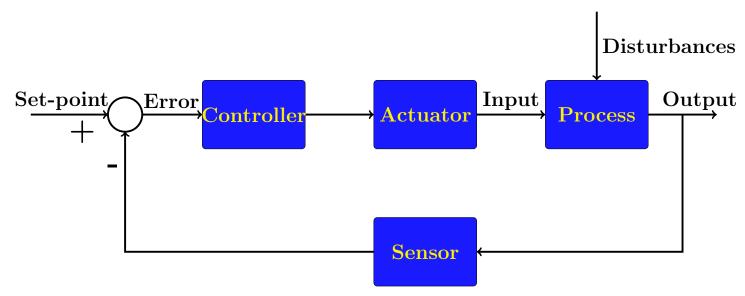
INTRODUCTION

• Incentives for chemical process control



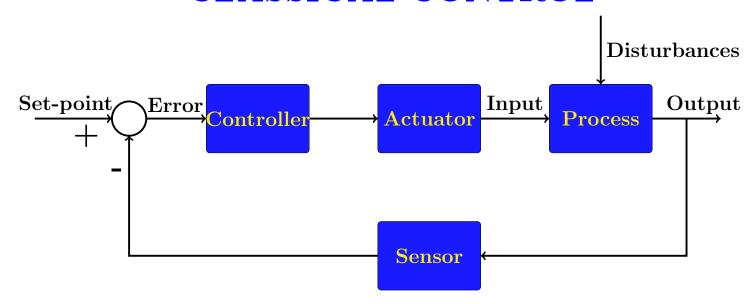
- Need for continuous monitoring and external intervention (process control)
- Objectives of a process control system
 - ♦ Ensuring stability of the process
 - ♦ Suppressing the influence of external disturbances
 - ♦ Optimizing process performance

FEEDBACK CONTROL LOOP



- How a feedback control loop (closed-loop system) works:
 - ♦ A variable describing the condition of a process (e.g., temperature, pressure, species concentration; known as an output) is measured by a sensor
 - ♦ The error between the measured output value and the desired value of this output (set-point) is calculated and fed to the controller
 - ♦ The controller computes a value of the manipulated input to the process to reduce the error
 - ♦ A control actuator (typically a valve) is used to apply the manipulated input value to the process

CLASSICAL CONTROL



- Classical control: single-input/single-output (SISO) control design
 - \diamond Proportional-integral-derivative (PID) control (error e(t))
 - ▶ Error reflects difference between measured output and set-point
 - \diamond Input/control action u(t)

$$u(t) = K_c e(t) + \underbrace{\frac{1}{\tau_I} \int_0^t e(\tau) d\tau}_{\mathbf{D}} + \underbrace{\tau_D \frac{de(t)}{dt}}_{\mathbf{D}}$$

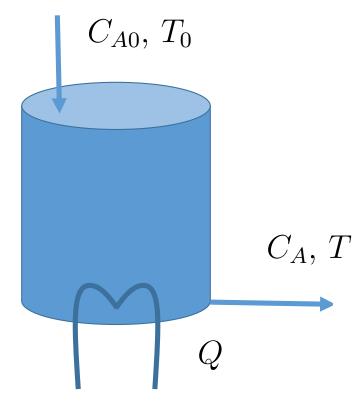
 $\diamond K_c, \tau_I, \tau_D$: scalar values that can be picked (tuned)

ADVANCED MODEL-BASED PROCESS CONTROL

- Advanced process control utilizes a process dynamic model explicitly in the controller design
 - ♦ A mathematical process model is developed:
 - Constructed from first-principles
 - ▶ Identified from input-output process data
 - ♦ The model describes the process dynamics (variation of the process state variables in time due to disturbances, inputs, and interactions between variables)
 - Controllers are synthesized based on the process model
- Advantages of model-based control
 - ♦ Possibility of improved closed-loop performance
 - Model accounts for inherent process characteristics (e.g., nonlinear behavior, multivariable interactions)
 - Characterization of limitations on achievable closed-loop stability, performance and robustness

NONLINEAR MODEL-BASED PROCESS CONTROL

• Example: continuous stirred tank reactor (CSTR)



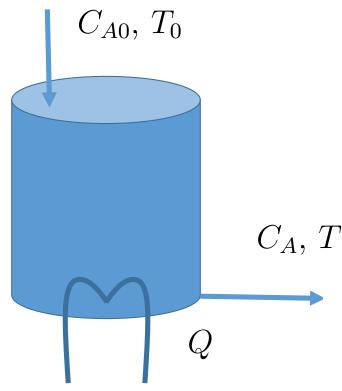
• Model: system of nonlinear ordinary differential equations (ODEs)

$$\frac{dT}{dt} = \frac{F}{V_r}(T_0 - T) + \frac{(-\Delta H)}{\rho C_p} k_0 e^{-E/RT} C_A + \frac{Q}{\rho C_p V_r} \Rightarrow x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} T - T_s \\ C_A - C_{As} \end{bmatrix}, \ \dot{x} = \frac{dx}{dt}$$

$$\frac{dC_A}{dt} = \frac{F}{V_r}(C_{A0} - C_A) - k_0 e^{-E/RT} C_A \qquad u = Q - Q_s, \ w = C_{A0} - C_{A0s}$$

NONLINEAR MODEL-BASED PROCESS CONTROL

• Example: continuous stirred tank reactor (CSTR)



• Model: system of nonlinear ordinary differential equations (ODEs)

$$\dot{x} = f(x, u, w)$$

- Techniques for nonlinear controller design for driving the process state to the operating steady-state
 - ♦ Lyapunov-based control

♦ Model predictive control

NONLINEAR PROCESS SYSTEMS

• State-space description

$$\dot{x} = f(x, u, w)$$

- $\diamond x \in X \subset \mathbb{R}^n$ is the state, $u \in U \subset \mathbb{R}^m$ is the manipulated input, $w \in W \subset \mathbb{R}^l$ is the disturbance, f is a vector function
- Explicit nonlinear feedback control law: u = h(x)
 - ♦ Control design technique: Lyapunov-based control (Y. Lin and E.D. Sontag, SCL, 1991; H. Khalil, Prentice Hall, 2002; P. D. Christofides and N. H. El-Farra, Springer-Verlag, 2005)
 - Renders the origin (steady-state) asymptotically stable
 - \diamond There exists a Lyapunov function V which satisfies

$$\dot{V} = \frac{\partial V(x)}{\partial x} f(x, h(x), 0) < 0, \ \forall \ x \in D$$
 $V : \text{energy of a physical system}$

- \diamond Typically, $V(x) = x^T P x$ (quadratic) and $\Omega_{\rho} \subseteq D$ is a level set of V where state constraints are met (i.e., $\Omega_{\rho} := \{x : V(x) \leq \rho\}$)
- $\diamond u = h(x)$ possesses a degree of robustness to disturbances and uncertainty
- Performance considerations and constraints are not directly/explicitly taken into account

MODEL PREDICTIVE CONTROL

• Model predictive control (MPC)

$$\min_{u \in S(\Delta)} \int_{t_k}^{t_{k+N}} l_T(\tilde{x}(\tau), u(\tau)) d\tau$$
s.t.
$$\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$$

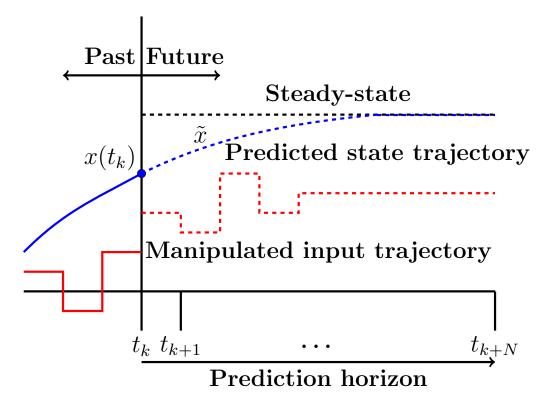
$$\tilde{x}(t_k) = x(t_k)$$

$$u(t) \in U, \ \tilde{x}(t) \in X, \ \forall \ t \in [t_k, t_{k+N})$$

• Quadratic tracking stage cost:

$$l_T(x, u) = x^T Q x + u^T R u$$

- $\diamond Q$, R are positive definite matrices
- Solve the optimization problem every Δ time units (sampling period)
 - \diamond At each sampling time t_k



- Solution is a piecewise-constant input trajectory
 - \diamond Each piece is held constant for a period Δ
 - \diamond Prediction horizon N

MODEL PREDICTIVE CONTROL

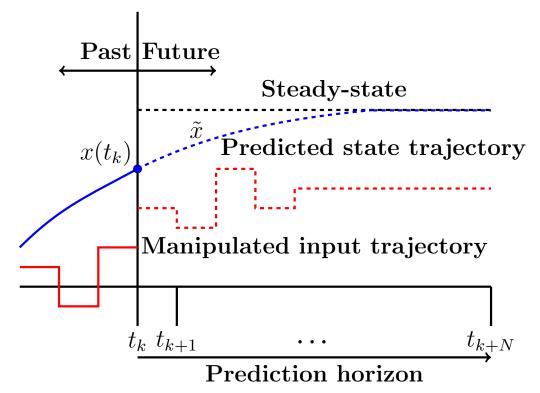
• Model predictive control (MPC)

$$\min_{u \in S(\Delta)} \int_{t_k}^{t_{k+N}} \left[\tilde{x}^T Q \tilde{x} + u^T R u \right] d\tau$$
s.t.
$$\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$$

$$\tilde{x}(t_k) = x(t_k)$$

$$u(t) \in U, \ \tilde{x}(t) \in X, \ \forall \ t \in [t_k, t_{k+N})$$

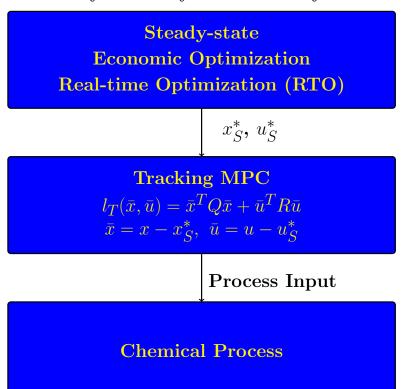
- Receding horizon implementation
 - Only the first piece of the input trajectory is applied
 - \triangleright Allows for feedback at every \triangle
 - Accounts for effects of disturbances and plant/model mismatch on the optimal solution
 - Longer prediction horizon may improve closed-loop performance



- Closed-loop stability is not guaranteed
- Approaches for closed-loop stability
 - Infinite/sufficiently long prediction horizon
 - ♦ Terminal cost/constraint
 - ♦ Contractive constraint

NEXT-GENERATION MANUFACTURING

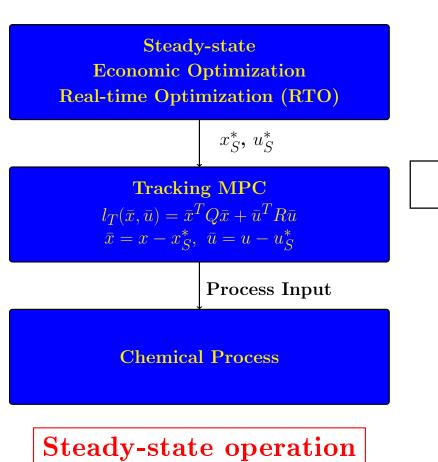
- Next-generation/smart manufacturing Objectives (J. Davis, T. Edgar, J. Porter, J. Bernaden and M. Sarli, Comput. Chem. Eng., 2012):
 - ♦ Profitability
 - ♦ Autonomy
 - ♦ Safety and cybersecurity



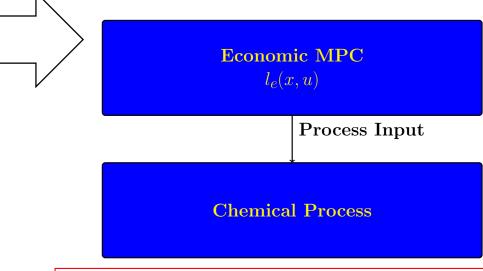
- Example: Moving away from a hierarchical approach to optimization and control
 - ♦ Upper layer:
 - Determine economically-optimal steady-state (real-time optimization (RTO)) (M. L. Darby,
 M. Nikolaou, J. Jones and D. Nicholson, JPC,
 2011)
 - ♦ Lower layer:
 - ▶ Feedback control drives the state of the process to the optimal steady-state
- Tighter integration of plant operation and process economic optimization

PROCESS ECONOMICS AND CONTROL

• Traditional Paradigm



- Integration of economic optimization and process control
- Generalization of MPC
 - ♦ General (economic) stage cost



Dynamic/time-varying operation

- Economic MPC (EMPC) potential use cases:
 - ♦ Time-varying objective function or constraints (M. Ellis and P. D. Christofides, AIChE J.,

2013; A. Gopalakrishnan and L. T. Biegler, CACE , 2013)

ECONOMIC MPC FORMULATION

• EMPC formulation:

$$\min_{u(\cdot) \in S(\Delta)} \int_{t_k}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) d\tau$$
s.t.
$$\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$$

$$\tilde{x}(t_k) = x(t_k)$$

$$u(t) \in U, \ \tilde{x}(t) \in X,$$

$$\forall \ t \in [t_k, t_{k+N})$$

$$|u(t_j) - u(t_{j-1})| \le \epsilon_d$$

$$j = k, \dots, k+N-1$$

- Components of EMPC:
 - ♦ Economic cost function
 - ♦ Dynamic model
 - ♦ State feedback measurement
 - ♦ Input and state magnitude constraints
 - ♦ Input rate of change constraints
- System equipped with a measure of instantaneous economics l_e
- Computes control actions that optimize economics
- Accounts for input and state constraints
 - ♦ Examples: temperature or flow rate bounds
- Prevents rapid variations in inputs which may damage actuators

LYAPUNOV-BASED ECONOMIC MPC

Boundedness / Time-varying Operation (Mode 1)

$$\min_{u(\cdot) \in S(\Delta)} \int_{t_k}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) \, d\tau$$

$$\text{s.t.} \qquad \dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$$

$$\tilde{x}(t_k) = x(t_k)$$

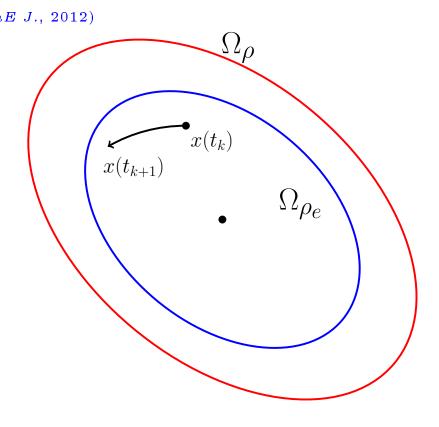
$$u(t) \in U, \ \tilde{x}(t) \in X, \ \forall \ t \in [t_k, t_{k+N})$$

$$|u_i(t_j) - h_i(\tilde{x}(t_j))| \le \epsilon_r, \ i = 1, \dots, m,$$

$$j = k, \dots, k+N-1$$

$$V(\tilde{x}(t)) \le \rho_e, \ \forall \ t \in [t_k, t_{k+N})$$

$$\text{if } V(x(t_k)) \le \rho_e \ \text{and} \ t_k < t_s$$



- Provable stability: boundedness of the closed-loop state in Ω_{ρ} ($\Omega_{\rho_e} \subset \Omega_{\rho}$)
- Provable feasibility: h(x) meets all state and input constraints

LYAPUNOV-BASED ECONOMIC MPC

Convergence to the Steady-State (Mode 2)

$$\min_{u(\cdot) \in S(\Delta)} \int_{t_k}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) d\tau$$
s.t.
$$\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t), 0)$$

$$\tilde{x}(t_k) = x(t_k)$$

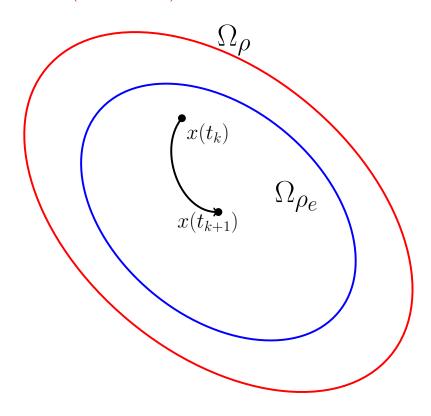
$$u(t) \in U, \ \tilde{x}(t) \in X, \ \forall \ t \in [t_k, t_{k+N})$$

$$|u_i(t_j) - h_i(\tilde{x}(t_j))| \le \epsilon_r, \ i = 1, \dots, m,$$

$$j = k, \dots, k + N - 1$$

$$\frac{\partial V(x(t_k))}{\partial x} f(x(t_k), u(t_k), 0)$$

$$\le \frac{\partial V(x(t_k))}{\partial x} f(x(t_k), h(x(t_k)), 0)$$
if $V(x(t_k)) > \rho_e$ or $t_k \ge t_s$



- Compute control actions that decrease the Lyapunov function
- Provable stability: convergence to a small neighborhood of the steady-state

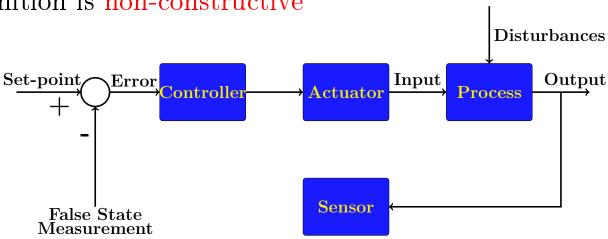
CYBERSECURITY AND PROCESS CONTROL SYSTEMS

- Cyberattacks on control systems seek to impact a physical process and can impact safety, profit, and production rates (A.A. Cárdenas et al., ASIACCS, 2011)
- Do cyberattackers care about attacking control and manufacturing systems?
 - ♦ 2010: Stuxnet (trellix.com)
 - > Attack on Iranian nuclear facilities
 - ▶ Worm entered systems via USB sticks and spread
 - > Searched for control system software
 - ▶ Ran centrifuges at conditions that cause breakdown
 - > Falsified information to main controller so that there was no indication of a problem
- How can we make it hard for cyberattackers to cause issues?
 - ♦ Design process to be cyberattack-resilient
 - ♦ Detect attacks and then shut down a process
 - ♦ Diagnose attacks to shut down only attacked components
 - ♦ Fight back against attackers
- How can we prevent reductions in agility while mitigating risk?

CYBERSECURITY AND PROCESS CONTROL SYSTEMS

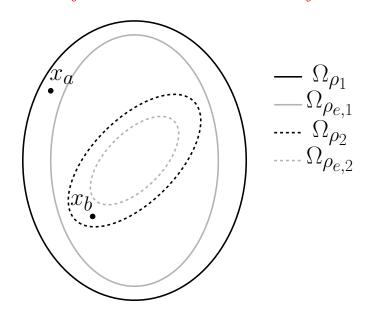
(H. Durand, Mathematics, 2018)

- Need for understanding control theory of attacks and mitigation strategies
- Cyberattacks on feedback controllers remove associations between state measurements and inputs
 - \diamond Undesired inputs $u \in U$ can be applied at a given state
- Attacks on sensors, actuators, or both
- Cyberattack-resilience for state measurement falsification requires:
 - \diamond There exist no possible input policies given the controllers used and their implementation strategies such that $x(t) \notin X$, for any allowable initial state $x_0 \in \bar{X}$ and $w(t) \in W$, $t \in [0, \infty)$
 - ▶ This definition is non-constructive



UNDERSTANDING RESILIENCE

- Attacks may be designed by reverse engineering known control laws
 - ♦ Suggests that randomly selecting the controller to be used at a given sampling time may make cyberattack design more difficult
 - ♦ Randomness in control design can only be considered if closed-loop stability is maintained under normal operation
 - Closed-loop stability and feasibility guarantees can be made with a randomized LEMPC implementation strategy
 - ▷ Cyberattack-resiliency is not guaranteed



- Implementation strategy:
 - \diamond Develop n_p LEMPC's and $h_1(x)$
 - \diamond At each t_k , randomly select one of the controllers until one is found for which:
 - $\triangleright x(t_k) \in \Omega_{\rho_i}, i = 1, \dots, n_p, \text{ for the } n_p th$ LEMPC

$$\triangleright x(t_k) \in \Omega_{\rho_1} \text{ for } h_1(x)$$

Process Description

• Continuous stirred tank reactor (CSTR) with second-order, exothermic, irreversible reaction of the form $A \to B$:

$$\frac{dC_A}{dt} = \frac{F}{V}(C_{A0} - C_A) - k_0 e^{\frac{-E}{RT}} C_A^2
\frac{dT}{dt} = \frac{F}{V}(T_0 - T) + \frac{-\Delta H}{\rho_L C_p} k_0 e^{\frac{-E}{RT}} C_A^2 + \frac{Q}{\rho_L C_p V}$$

- Control objective: regulate the process in an economically optimal time-varying fashion while maintaining closed-loop stability
 - ♦ Economic cost:

$$\int_{t_k}^{t_{k+N}} \left[k_0 e^{-\frac{E}{RT(\tau)}} C_A(\tau)^2 \right] d\tau$$

♦ Manipulated input constraints

$$0.5 \le C_{A0} \le 7.5 \text{ kmol/m}^3$$
 $-5.0 \times 10^5 \le Q \le 5.0 \times 10^5 \text{ kJ/h}$

Deviation variables:

$$x_1 = C_A - C_{As}, \quad x_2 = T - T_s$$

 \diamond Process model in input-affine form $\dot{x} = \tilde{f}(x) + gu$

Lyapunov-Based Controller Design

- Lyapunov-based controller for the inlet concentration: $h_{1,1}(x) = 0 \text{ kmol/m}^3$
 - ♦ Lyapunov-based controller for the heat rate input:
 - Sontag's Formula (Y. Lin and E.D. Sontag, SCL, 1991)

$$h_{2,1}(x) = \begin{cases} -\frac{L_{\tilde{f}}V_1 + \sqrt{L_{\tilde{f}}V_1^2 + L_{g_2}V_1^4}}{L_{g_2}V_1}, & \text{if } L_{g_2}V_1 \neq 0\\ 0, & \text{if } L_{g_2}V_1 = 0 \end{cases}$$

 \diamond A quadratic Lyapunov function of the form $V_1(x) = x^T P x$ with:

$$P = \left[\begin{array}{cc} 1200 & 5 \\ 5 & 0.1 \end{array} \right]$$

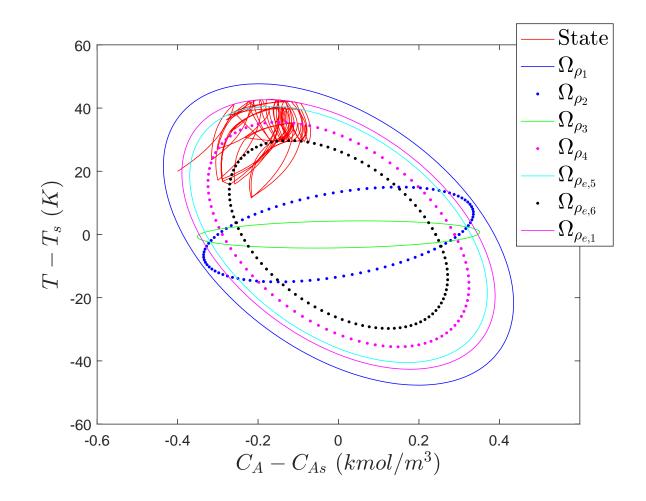
- \diamond Stability region $\rho_1 = 180$ (i.e., $\Omega_{\rho_1} = \{x \in \mathbb{R}^2 : V_1(x) \leq \rho_1\}$)
- Process state initialized at $x_{init} = [-0.4 \text{ kmol/m}^3 \text{ 20 K}]^T$
- LEMPC parameters: N = 10, $\Delta = 0.01$ h
- Process simulated with an integration step size of 10⁻⁴ h

Randomized LEMPC Development

• 6 LEMPC's were designed

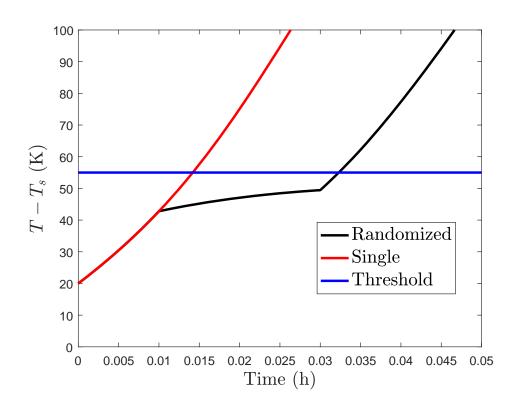
$$\diamond \Omega_{\rho_i} \subseteq \Omega_{\rho_1}, i = 1, \dots, 6$$

- $\diamond h_{i,1} = 0 \text{ kmol/m}^3$
- $\diamond h_{i,2}$ designed via Sontag's control law
- \diamond Closed-loop state is maintained within Ω_{ρ_1} throughout 1 h of operation in the absence of a cyberattack



Randomized LEMPC and LEMPC Under a Cyberattack

- Cyberattack with $x_f = [-0.0521 \text{ kmol/m}^3 8.3934 \text{ K}]^T$ is applied to a single LEMPC and the randomized LEMPC implementation strategy
- Randomized LEMPC results depend on seed to random number generator
- Randomized LEMPC barely delayed the time until $x_2 > 55$ K compared to the single LEMPC (0.0142 h)



Seed	Time $x_2 > 55$ (h)
5	0.0231
10	0.0144
15	0.0142
20	0.0323
25	0.0247
30	0.0142
35	0.0142
40	0.0146
45	0.0247
50	0.0142

CYBERATTACK DETECTION STRATEGIES

- Randomized LEMPC implementation strategy could not guarantee that no problematic inputs could be applied over time
 - ♦ Problem: No principle to the randomness besides luck
 - ♦ Demonstrates the need for principled design of cyberattack-handling strategies
- Three concepts for utilizing LEMPC to attempt to detect attacks were explored

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(H. Durand and M. Wegener, Mathematics, 2020; H. Oyama and H. Durand, AIChE J., 2020; H. Oyama et al., Front. Chem. Eng., 2022; K. Kasturi Rangan et al., DYCOPS, 2022)
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- ♦ LEMPC with random control law modifications to probe for cyberattacks
 - > Safety guaranteed under actuator attacks (not sensor measurement attacks)
- ♦ State feedback LEMPC with an attack detection strategy based on state predictions at each sampling time
 - ▷ Safety guaranteed under actuator attacks but only for a short time under undetected sensor measurement attacks
- ♦ Output feedback LEMPC (M. Ellis, J. Zhang, J. Liu and P. D. Christofides, SCL, 2014; L. Lao, M. Ellis, H. Durand and P. D. Christofides, AIChE J., 2015) with an attack detection strategy based on redundant state estimators
 - ▷ Safety guaranteed under limited sensor attacks (not actuator attacks)

OBSERVABILITY ASSUMPTION

• M sets of measurements are continuously available:

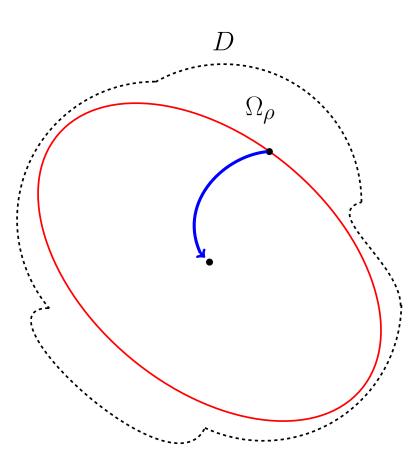
$$y_i(t) = k_i(x(t)) + v_i(t)$$

- \diamond k_i is vector-valued function, and v_i represents the measurement noise associated with the measurements y_i
- $\diamond v_i \in V_i \subset \mathbb{R}_i^q \ (|v_i| \leq \theta_{v,i}), \ i = 1, \dots, M$
- A deterministic observer exists for each of the M sets of measurements:

$$\dot{z}_i = F_i(\epsilon_i, z_i, y_i)$$

- \diamond Observer estimate z_i ; $\epsilon_i > 0$
- Assumptions:
 - \diamond For an initial state estimate with sufficiently low error between z_i and x, $h(z_i)$ maintains the closed-loop state in Ω_{ρ}
 - \diamond There exists a time t_{bi} such that:

$$|z_i(t) - x(t)| \le \epsilon_{mi}$$



CYBERATTACK-RESILIENT OUTPUT FEEDBACK LEMPC

- Cyberattacks on state measurements could impact the state estimate used by the LEMPC
- If the estimate is sufficiently incorrect, the closed-loop state may exit Ω_{ρ}
- Estimator properties suggest an attack detection methodology
 - $\diamond |z_i(t) x(t)| \le \max\{e_{mi}\}, i = 1, \dots, M$
 - \diamond Implies $|z_i(t) z_j(t)| \leq \epsilon_{\max}$, $i, j = 1, \ldots, M$, when no attack occurs
 - Condition can be used with redundant estimators to attempt to flag falsified sensor measurements

$$\min_{u(t) \in S(\Delta)} \int_{t_k}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) d\tau$$
s.t. $\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t))$

$$\tilde{x}(t_k) = z_1(t_k)$$

$$\tilde{x}(t) \in X, \, \forall \, t \in [t_k, t_{k+N})$$

$$u(t) \in U, \, \forall \, t \in [t_k, t_{k+N})$$

$$V(\tilde{x}(t)) \leq \rho_{e,1}, \, \, \forall \, t \in [t_k, t_{k+N}),$$

$$\text{if } \tilde{x}(t_k) \in \Omega_{\rho_{e,1}}$$

$$\frac{\partial V(\tilde{x}(t_k))}{\partial x} (f(\tilde{x}(t_k), u(t_k)))$$

$$\leq \frac{\partial V(\tilde{x}(t_k))}{\partial x} (f(\tilde{x}(t_k), h(x(t_k))))$$

$$\text{if } \tilde{x}(t_k) \in \Omega_{\rho}/\Omega_{\rho_{e,1}}$$

CYBERATTACK-RESILIENT OUTPUT FEEDBACK LEMPC

- Consider that at least one state estimate is not impacted by an attacker
- If $|z_i(t) z_j(t)| > \epsilon_{\max}$, $i, j = 1, \dots, M$, flag an attack
- If $|z_i(t) z_j(t)| \le \epsilon_{\max}$, i, j = 1, ..., M, but an attack occurred:
 - \diamond Closed-loop state will be maintained in Ω_{ρ} over the subsequent sampling period under sufficient conditions
 - \triangleright Examples: sufficiently small $\rho_{e,1}$, θ , and Δ

$$\min_{u(t) \in S(\Delta)} \int_{t_k}^{t_{k+N}} l_e(\tilde{x}(\tau), u(\tau)) d\tau$$
s.t. $\dot{\tilde{x}}(t) = f(\tilde{x}(t), u(t))$

$$\tilde{x}(t_k) = z_1(t_k)$$

$$\tilde{x}(t) \in X, \ \forall \ t \in [t_k, t_{k+N})$$

$$u(t) \in U, \ \forall \ t \in [t_k, t_{k+N})$$

$$V(\tilde{x}(t)) \leq \rho_{e,1}, \ \forall \ t \in [t_k, t_{k+N}),$$
if $\tilde{x}(t_k) \in \Omega_{\rho_{e,1}}$

$$\frac{\partial V(\tilde{x}(t_k))}{\partial x} (f(\tilde{x}(t_k), u(t_k)))$$

$$\leq \frac{\partial V(\tilde{x}(t_k))}{\partial x} (f(\tilde{x}(t_k), h(x(t_k))))$$
if $\tilde{x}(t_k) \in \Omega_{\rho}/\Omega_{\rho_{e,1}}$

MOTIVATION FOR HANDLING SIMULTANEOUS ACTUATOR AND SENSOR ATTACKS

• Continuous stirred tank reactor (CSTR) with second-order $A \to B$ reaction:

$$\frac{dC_A}{dt} = \frac{F}{V}(C_{A0} - C_A) - k_0 e^{\frac{-E}{RT}} C_A^2
\frac{dT}{dt} = \frac{F}{V}(T_0 - T) + \frac{-\Delta H}{\rho_L C_p} k_0 e^{\frac{-E}{RT}} C_A^2 + \frac{Q}{\rho_L C_p V}$$

- Control objective: Optimize process economics while maintaining the closed-loop state in Ω_{ρ_1}
 - ♦ Economic cost:

$$\int_{t_k}^{t_{k+N}} [k_0 e^{-\frac{E}{RT(\tau)}} C_A(\tau)^2] d\tau$$

♦ Manipulated input constraint

$$0.5 \le C_{A0} \le 7.5 \text{ kmol/m}^3$$

♦ Deviation variables:

$$x_1 = C_A - C_{As}, \quad x_2 = T - T_s$$

 \diamond Process model in input-affine form $\dot{x} = \tilde{f}(x) + gu$

MOTIVATION FOR HANDLING SIMULTANEOUS ACTUATOR AND SENSOR ATTACKS

- Lyapunov-based controller: $h(x) = -1.6x_1 0.01x_2$ (M. Heidarinejad, J. Liu, and P. D. Christofides, SCL, 2012)
 - \diamond A quadratic Lyapunov function of the form $V_1(x) = x^T P x$ with:

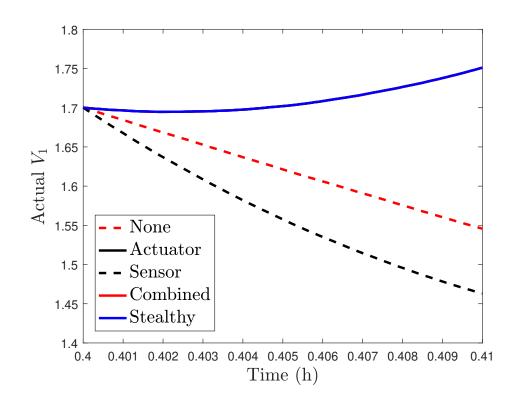
$$P = \left[\begin{array}{cc} 110.11 & 0 \\ 0 & 0.12 \end{array} \right]$$

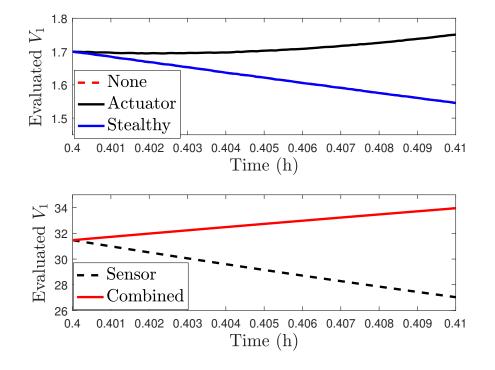
- \diamond Stability region $\rho_1 = 440$ (i.e., $\Omega_{\rho_1} = \{x \in \mathbb{R}^2 : V(x) \leq \rho_1\}$)
- $\diamond \ \Omega_{\rho_{e_1}} \subset \Omega_{\rho}, \ \rho_{e_1} = 330$
- LEMPC parameters: $N = 10, \Delta = 0.01 \text{ h}$
- Process simulated with an integration step size of 10^{-3} h
- The LEMPC receives full state feedback with the full system state $x = [x_1 \ x_2]^T$
- Attack detection policy (initialized at 0.4 h when attack begins): Check if Lyapunov function evaluated at the state measurement decreases over Δ

VARIOUS ATTACK POLICIES

(H. Oyama, D. Messina, K. K. Rangan, and H. Durand, Frontiers in Chemical Engineering, 2022)

- Actuator attack ($u = 0.5 \text{ kmol/m}^3$): Discoverable
- False sensor measurement $(x_1 + 0.5 \text{ kmol/m}^3)$: Not discoverable (no safety issue)
- Combined actuator and sensor attack: Discoverable
- Stealthy actuator and sensor attack (sensor measurements follow trajectory they should have taken): Not discoverable
 - ♦ State moves closer to safe operating region boundary



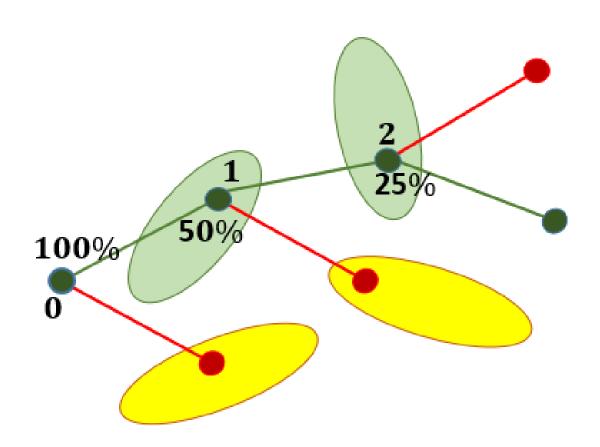


PREVENTING SAFETY ISSUES DURING SIMULTANEOUS ATTACKS

- Multiple detector types can be used to aid in cornering an attacker
 - ♦ Examples:
 - ▶ Redundant estimators and forcing the decrease of the Lyapunov function across a sampling period
 - ▶ Redundant estimators and state predictions with a redundant control law
 - Resilient under sufficient conditions
 - Closed-loop state cannot leave a safe operating region in the presence of individual or simultaneous attacks before attack detection
 - ▶ Potentially challenging to obtain reasonable control law parameters satisfying resilience theory (K. Nieman et al., J. Loss Prev. Process Ind., 2023)
 - ♦ Combining detectors may create simultaneous detection and diagnosis strategies (D. Messina and H. Durand, ADCHEM, 2024)
 - ▶ Violation of one detection metric and not the second indicates whether the attack is on actuators or sensors based on detection metric construction
- Need a strategy for detecting attacks on sensors that might flag them even with all sensors being compromised

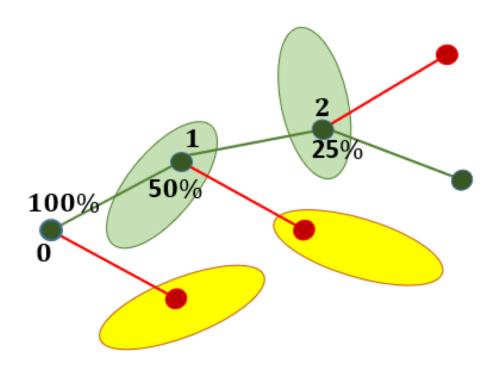
(H. Oyama et al., Digital Chemical Engineering, 2023)

- Set up expectations for measurements that would be "hard" to fake
 - ♦ At every sampling time, two control actions are available
 - ♦ Should result in non-overlapping possible sets of measured states
 - ♦ One of the two is randomly selected



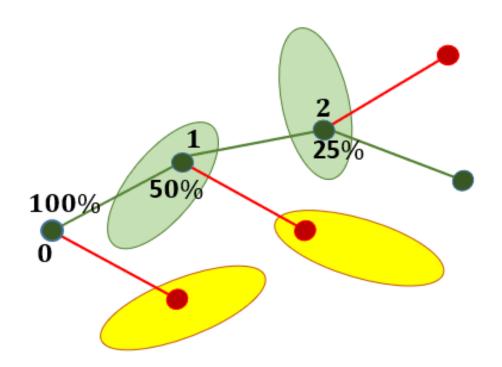
• At the first sampling time:

- ♦ Both possible control actions keep the closed-loop state in a safe operating region with or without a sensor attack
- ♦ The controller randomly applies one of the two
- ♦ The attacker is forced to guess a state measurement in one of the two possible measured state sets to avoid being caught
- \diamond 50% chance to guess correctly

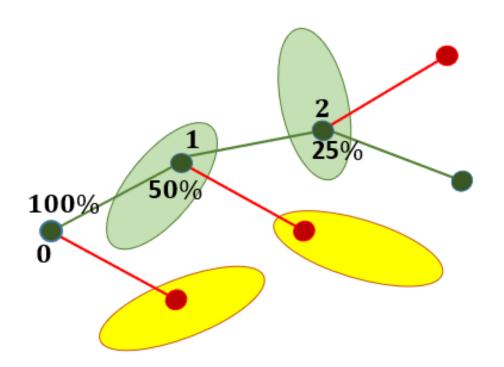


• At the second sampling time:

- ♦ The attacker is "caught" if they guessed the wrong set
- ♦ If they guessed the correct set, a false measurement is provided to a controller
- ♦ One of two known control actions will be applied that keeps the closed-loop state in a safe operating region regardless of whether there is an attack
- \diamond 50% chance to guess the next set that the detector is expecting
- \diamond 25% chance they guess correctly twice in a row



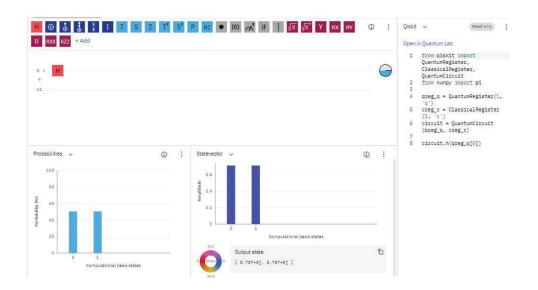
- Applying this strategy continuously makes it unlikely an attacker goes long without being detected
- Validate that the measurement p sampling times ago was correct
 - Otherwise the attacker is unlikely to have not been noticed between then and now
 - \diamond p sampling periods of safety under rogue inputs from an originally correct state measurement required



IMPLEMENTING CONTROL ON QUANTUM COMPUTERS

Quantum Computing

- Quantum computing is a technology of recent interest in chemical engineering (D.E Bernal et al., AIChE J., 2022; A. Ajagekar and F. You, CACE, 2020; A. Ajagekar, T. Humble, and F. You, CACE, 2020; A. Ajagekar, Energy, 2019)
- Quantum computers exist today of different types
- Quantum annealing
 - Solves an optimization problem
- Gate-based computers



- Control is implemented using computing
 - ♦ Advances in computing should be evaluated for control relevance

QUANTUM COMPUTING-IMPLEMENTED CONTROL

- Computational bottlenecks faced in control implementation raise the question of whether these could be overcome by next-generation computing frameworks
- Before investigating that, we must address the safety implications of quantum computing
 - ♦ Quantum computing introduces non-standard operating challenges
 - ▷ "Noise" (errors) in computation in near-term quantum devices (noisy intermediate-scale quantum (NISQ) devices)
 - ▶ Non-deterministic algorithms
 - ♦ Need to define algorithm and hardware requirements to facilitate guidance of algorithm designers and co-design principles for algorithms and control laws to promote safety

• Key needs:

- ♦ Develop strategies for simulating closed-loop systems under controllers implemented on quantum computers (K. Kasturi Rangan et al., AIP Publishing, 2023)
- ♦ Develop theory relating non-determinism due to noise and algorithms to safety principles for quantum computation (K. Nieman et al., IECR, 2022)
- ♦ Co-design control laws and quantum algorithms to pair efficiency benefits with safety (K. Nieman et al., Digital Chemical Engineering, 2024)

QUANTUM MECHANICS FOR QUANTUM COMPUTING VS. CHEMISTRY

- Reminders from chemistry:
 - ♦ "Time-independent Schrödinger equation" (eigenvalue-eigenvector relationship)

$$\hat{H}(x,t)\psi(x,t) = E\psi(x,t)$$

♦ Time-dependent Schrödinger equation

$$\hat{H}(x,t)\psi(x,t) = i\hbar \frac{\partial \psi(x,t)}{\partial t}$$

- $\hat{H}(x,t)$: Hamiltonian (total energy operator)
- \bullet Energy E
- \hbar : Reduced Planck constant
- $\psi(x,t)$: Wavefunction of the quantum system
 - ♦ Contains information about position of a quantum system
 - \diamond Example: $\psi(x,t)$ is the wavefunction of an electron
 - $\flat \psi(x_0, t_0)^* \psi(x_0, t_0) dx$ conveys the probability that the quantum particle will be found in a spatial interval with width dx around x_0 at time t_0 (T. Engel,

QUANTUM MECHANICS FOR QUANTUM COMPUTING VS. CHEMISTRY

- Wavefunctions are derived from a more fundamental notion of "quantum states"
 - ♦ "Time-independent Schrödinger equation" (eigenvalue-eigenvector relationship)

$$\bar{H}(t) |\Psi(t)\rangle = E |\Psi(t)\rangle$$

♦ Time-dependent Schrödinger equation

$$\bar{H}(t) |\Psi(t)\rangle = i\hbar \frac{\partial |\Psi(t)\rangle}{\partial t}$$

- $|\Psi(t)\rangle$ is the "quantum state"
 - ♦ "Dirac notation"
- Wavefunctions are derived from the quantum state in a way that makes them particularly good for representing information about position
- Position is continuous
- Gate-based quantum computers generally stay with the binary concept of classical computing
 - ♦ We only want to have 2 possible quantum states for the system
 - ♦ Position will not work for this

QUANTUM MECHANICS FOR QUANTUM COMPUTING VS. CHEMISTRY

- Wavefunctions are derived from a more fundamental notion of "quantum states"
 - ♦ "Time-independent Schrödinger equation" (eigenvalue-eigenvector relationship)

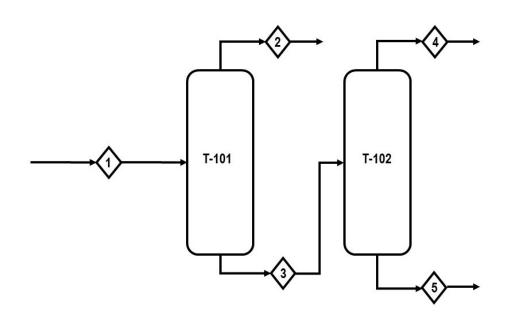
$$\bar{H}(t) |\Psi(t)\rangle = E |\Psi(t)\rangle$$

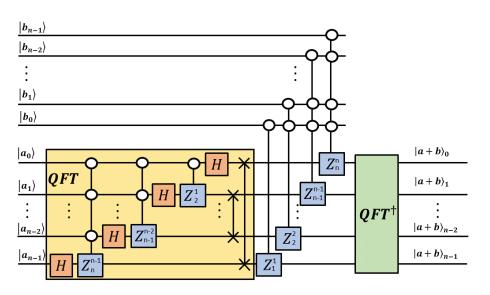
♦ Time-dependent Schrödinger equation

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- $|\Psi(t)\rangle$ is the "quantum state"
 - ♦ "Dirac notation"
- Wavefunctions are derived from the quantum state in a way that makes them particularly good for representing information about position
- Position is continuous
- Gate-based quantum computers generally stay with the binary concept of classical computing
 - ♦ Wavefunctions are not used in quantum computing
 - \diamond Two possible quantum states: $|0\rangle$ and $|1\rangle$ (regardless of actual implementation)

CONCEPTUALIZING QUANTUM CIRCUITS

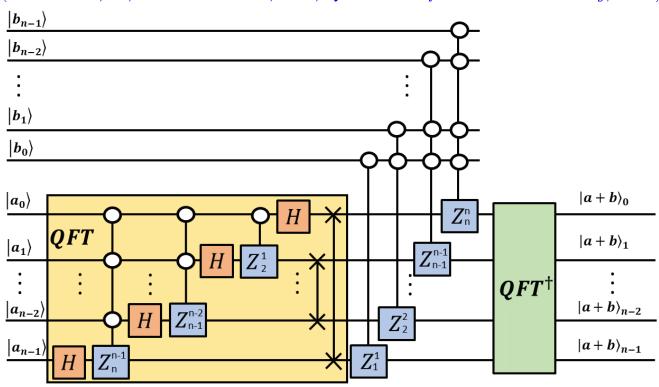




- Each unit of a chemical plant changes the state of a process stream
 - Symbols and labeling for process units create meaning for chemical engineers regarding the expected state changes
- Each block ("gate") in a quantum circuit changes the state of a quantum system
 - Symbols and labeling for the gates create meaning regarding the expected state changes
 - \diamond Example: H gate puts a qubit in an equal superposition of two states

QFT-BASED ADDITION

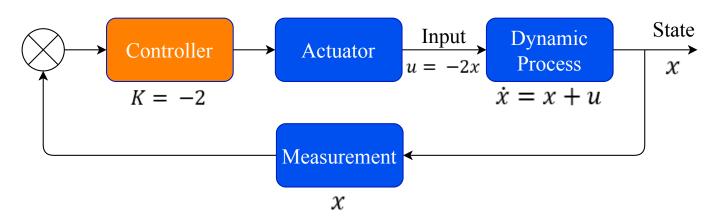
(Ruiz-Perez, L., Garcia-Escartin, J.C., Quantum Information Processing, 2017)



- QFT-based addition: Add two integers a and b (S. Anagolum, Github)
- Binary representations of both numbers are translated to qubit states
- Quantum gates are applied (including those in the inverse QFT, QFT[†]) to obtain final qubit states representative of the bits of the sum

QUANTUM COMPUTING-IMPLEMENTED CONTROL EXAMPLE

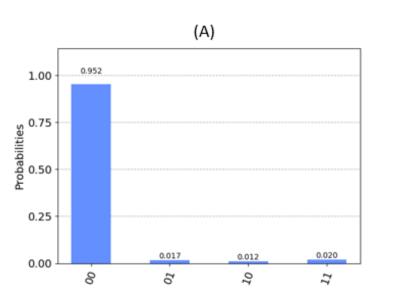
Motivation

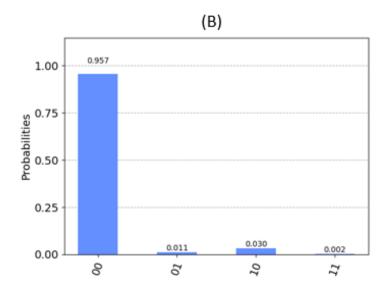


- Address the safety of running a control strategy using a quantum algorithm on a quantum device
 - ♦ Perform initial study through a simulation
 - \diamond A linear dynamic process, $\dot{x} = x + u$, classically stabilized using the control law u = -2x
 - \diamond Evaluate u = -2x as the negative of x + x
 - Addition performed using a quantum simulator (qasm_simulator) accessed via Qiskit
 - ♦ Quantum simulator does not inherently have noise
 - ▶ Required to select a noise model

QUANTUM COMPUTING-IMPLEMENTED CONTROL EXAMPLE

Noise Models

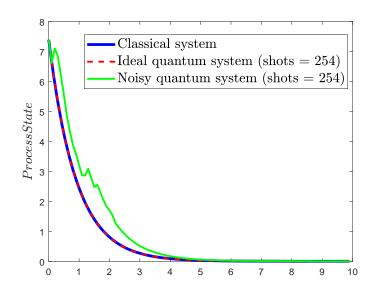


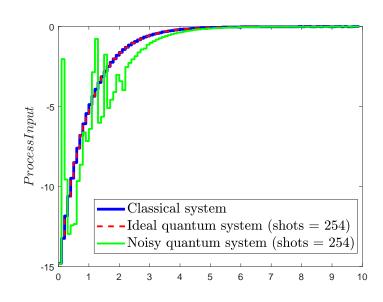


- A depolarizing error parameter for qasm_simulator was selected using command for modeling the noise from the 5-qubit quantum device, ibmq_manila, on the qasm_simulator
 - \diamond The controlled Z gate was simulated with both the qasm_simulator using this noise model from the device backend and with the depolarizing error parameter set to a fixed value on qasm_simulator
- A depolarizing error parameter of 0.05 was determined to sufficiently approximate the results from the simulations based on ibmq manila

QUANTUM COMPUTING-IMPLEMENTED CONTROL Results

- Comparison between the state trajectories (left) and input trajectories (right) when run with 254 shots for x(0) = 7.4
 - ♦ Classical computer ("Classical system"),
 - ♦ Quantum simulator with 254 shots and no noise ("Ideal quantum system")
 - ♦ Quantum simulator with 254 shots and noise ("Noisy quantum system")
- Some deviation is observed between the noisy system and the other two, related to the size (in binary) of the state measurement and number of shots



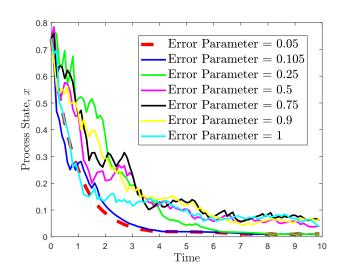


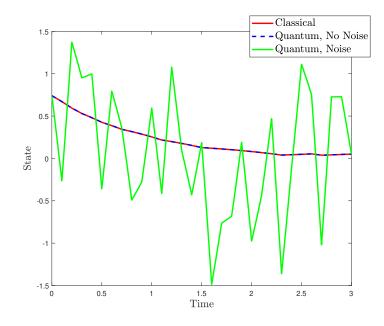
QUANTUM COMPUTING-IMPLEMENTED CONTROL

Moving Toward Theory

- Why does the closed-loop state move toward the origin despite the noise?
 - ♦ Understand potential of using NISQ devices for control
- Factors include number of stabilizing inputs and likelihood of selecting the "correct" input for a given shot count
 - ♦ Example:
 - ▶ Rounded state measurement: 74

 - \triangleright Number of inputs that would cause x to move toward the origin: 182 (71.09%)
 - Sampling period length can also be key
 - ▷ Suggest co-design of quantum algorithms and control laws





ADVANCED CONTROL AND QUANTUM COMPUTATION

- Rigorous theory for LEMPC makes it attractive for considering the implications of non-deterministic inputs on stability guarantees
 - Initial investigations of closedloop stability of quantum computing-implemented inputs should focus on simple quantum computing algorithms

State MeasurementControl Action0000111100011110

1010

0010

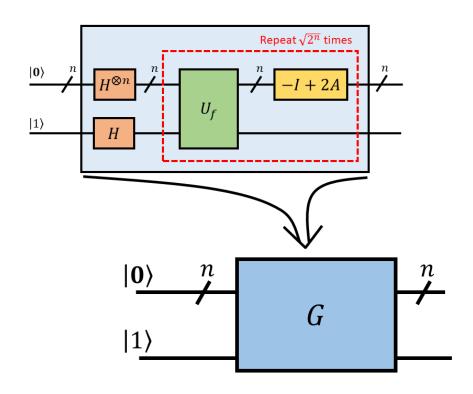
Table 1: LEMPC solution lookup table

• Consider LEMPC solutions in a look-up table

- ♦ For relating to quantum computing, must express state measurements and inputs in binary
- ⋄ Requires quantization of state measurements for LEMPC
- ♦ Also quantize control actions output by LEMPC

SEARCHING AN LEMPC LOOKUP TABLE VIA MODIFIED GROVER'S SEARCH

- Grover's search algorithm is a quantum computing algorithm for searching an unsorted list (Yanofsky and Mannucci, Cambridge University Press, 2008)
- A modified version of Grover's algorithm could be used to search the LEMPC lookup table
 - Not efficient for solving this problem
 - Show how non-deterministic inputs can be generated by a quantum computing algorithm tied to LEMPC



- Modified Grover's algorithm implementation strategy:
 - ♦ Use a series of controlled Grover blocks to represent the state/input pairings
 - \diamond Measurements return the "correct" input with probability λ

IMPLICATIONS FOR CLOSED-LOOP STABILITY

- Probability of obtaining the expected control action from Grover's algorithm: λ
- Consider x(t) and $\bar{x}(t) \in \Omega_{\rho_e}$
 - \diamond Control action computed by the LEMPC on a classical computer would maintain $x(t_k)$ and $\bar{x}(t_k)$ in Ω_{ρ} for $t \in [t_k, t_{k+1})$
 - \diamond The modified Grover algorithm would return the same control action as the classical computer with probability λ
 - ♦ Conclusion:

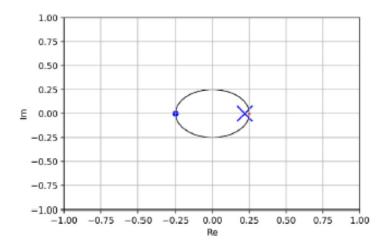
$$\triangleright \mathbf{P}(x(t), \bar{x}(t) \in \Omega_{\rho} \forall t \in [t_k, t_{k+1})) \ge \lambda$$

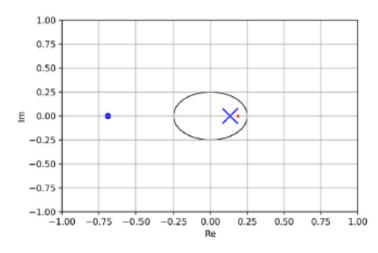
- Consider x(t) and $\bar{x}(t) \in \Omega_{\rho}/\Omega_{\rho_e}$
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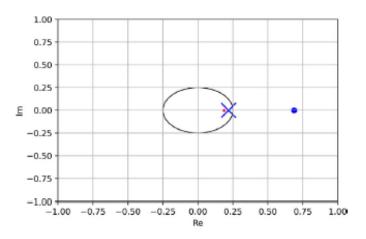
$$ightharpoonup \mathbf{P}(x(t), \bar{x}(t) \in \Omega_{\rho} \forall \ t \in [t_k, t_{k+1})) \geq \lambda$$

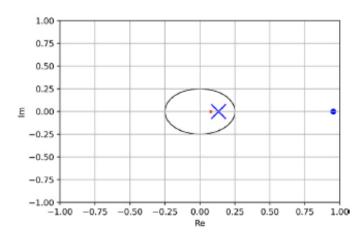
AMPLITUDE AMPLIFICATION

- Amplitudes: c_0 and c_1 in $c_0 |0\rangle + c_1 |1\rangle$
- Modify amplitude to change the probability of measuring the qubits in a certain state









AMPLITUDE AMPLIFICATION AND MPC

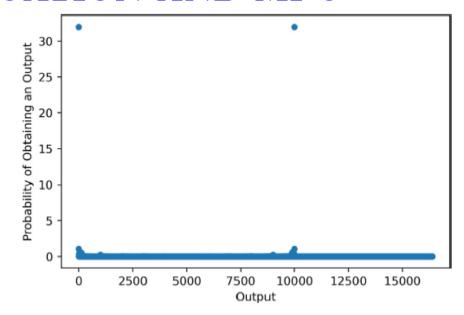
• MPC-like optimization problem

$$\max_{u \in S(\Delta)} \quad \tilde{x}(t_{k+N})$$
s.t.
$$\dot{\tilde{x}}(t) = \tilde{x}(t) + u(t)$$

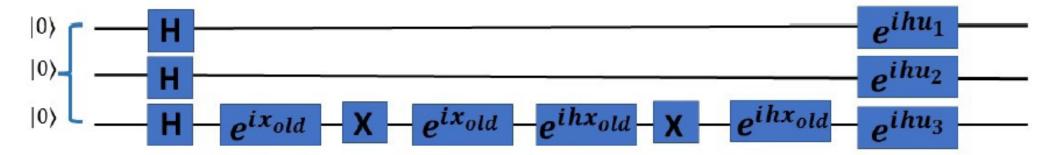
$$\tilde{x}(t_k) = x(t_k)$$

$$u(t) \in [-5, -4, \dots, 3, 4]$$

- Solve with amplitude amplification strategy (D. Koch et al., Entropy, 2022)
 - ♦ Solutions to the integer program in the phases of a quantum operator
 - Raise probability of measuring qubits in a state related to optimization solution



- Superposition-based parallelization of numerical integration
 - \diamond One integration step: $x_{new} = x_{old} + h(x_{old} + u)$



CONCLUSIONS

- Next-generation manufacturing values flexibility and profitability
 - ♦ Facilitated by automation advances such as economic model predictive control
 - ♦ Flexible and profitable systems may not be secure
 - ▶ Attacks on control systems may undermine process safety
- Integrated detection and control policies geared toward nonlinear systems have potential to enable attacks of various types to be detected before causing safety issues
 - ⋄ Requires sufficient control-theoretic conditions
 - ♦ May require at least some sensors to be secure
 - ▶ Handling attacks after detection likely requires some actuators to be secure
- Fundamental notions of cyberattack-resilience and discoverability for nonlinear systems provide insights into potential future directions for securing controllers
- Quantum computing provides another interesting potential direction for the future of next-generation manufacturing
 - ♦ Control theory and practice require further exploration to determine if benefit exists for quantum computing-implemented control

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