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# **Inversion of Linear Systems**

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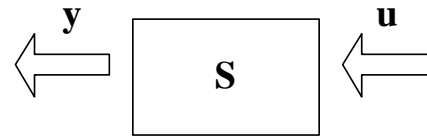
# Outline

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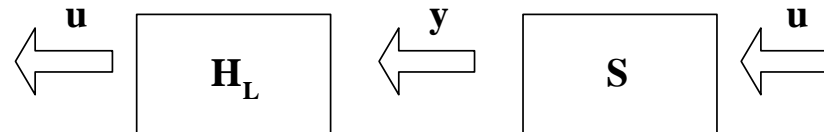
- Motivation
- Approaches to inversion
  - Massey-Sain algorithm
  - Moylan's algorithm
- Partial invertibility
- Future work

# What is Inversion?

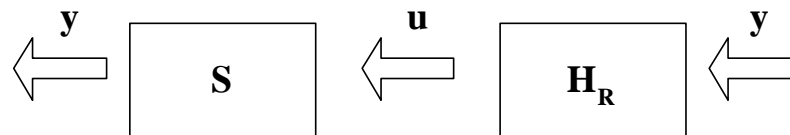
- Given the system



- Left Inverse: Outputs  $u$  when  $y$  is the input



- Right Inverse: Given a desired  $y$ , outputs the necessary  $u$



- Here, we will focus on linear time-invariant systems

# Why Study Inversion?

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Inverse systems are useful in a variety of situations

- Left inverses:
  - Communication systems: determining input to a dynamical channel
  - Fault detection and identification
  - Coding theory
- Right inverses:
  - Feed-forward control
  - Disturbance decoupling

# Approach 1: Transfer Function Matrix

- For SISO systems, inversion is easily handled through a transfer function approach

$$\frac{Y(z)}{U(z)} = \mathcal{S}(z) \Rightarrow \frac{U(z)}{Y(z)} = \frac{1}{\mathcal{S}(z)}$$

- Inverse may be non-causal!
- For MIMO systems, problem is more complicated

$$\begin{bmatrix} Y_1(z) \\ Y_2(z) \\ \vdots \\ Y_p(z) \end{bmatrix} = \underbrace{\begin{bmatrix} \mathcal{S}_{11}(z) & \mathcal{S}_{12}(z) & \cdots & \mathcal{S}_{1m}(z) \\ \mathcal{S}_{21}(z) & \mathcal{S}_{22}(z) & \cdots & \mathcal{S}_{2m}(z) \\ \vdots & \vdots & \ddots & \vdots \\ \mathcal{S}_{p1}(z) & \mathcal{S}_{p2}(z) & \cdots & \mathcal{S}_{pm}(z) \end{bmatrix}}_{\mathcal{S}(z)} \begin{bmatrix} U_1(z) \\ U_2(z) \\ \vdots \\ U_m(z) \end{bmatrix}$$

- Left inverse: find  $\mathcal{H}_L(z)$  such that  $\mathcal{H}_L(z)\mathcal{S}(z) = \mathbf{I}_m$  (requires  $p \geq m$ )
- Right inverse: find  $\mathcal{H}_R(z)$  such that  $\mathcal{S}(z)\mathcal{H}_R(z) = \mathbf{I}_p$  (requires  $m \geq p$ )

# Approach 2: State-Space Methods

State-space description of system  $\mathcal{S}$  is

$$x[k + 1] = Ax[k] + Bu[k]$$

$$y[k] = Cx[k] + Du[k]$$

where  $x \in \mathbb{R}^n$ ,  $y \in \mathbb{R}^p$  and  $u \in \mathbb{R}^m$

- Will focus on left-invertibility in rest of talk
- **Definition:** System  $\mathcal{S}$  has an inverse with delay  $L$  if  $u[k]$  can be uniquely determined from  $y[k], y[k + 1], \dots, y[k + L]$  (and perhaps  $x[k]$ )
  - Smallest such  $L$  is called “inherent delay” of system
- What are the conditions on matrices  $A, B, C, D$  for the system to be invertible?

# Massey-Sain Algorithm (1)

- Appeared in IEEE Transactions on Automatic Control, vol. 14, 1969
- Find expressions for output in terms of input:

$$y[k] = Cx[k] + Du[k]$$

$$y[k + 1] = Cx[k + 1] + Du[k + 1]$$

$$= CAx[k] + CBu[k] + Du[k + 1]$$

- Continuing in this way, we get

$$\underbrace{\begin{bmatrix} y[k] \\ y[k + 1] \\ y[k + 2] \\ \vdots \\ y[k + L] \end{bmatrix}}_{Y_{[k,L]}} = \underbrace{\begin{bmatrix} C \\ CA \\ CA^2 \\ \vdots \\ CA^L \end{bmatrix}}_{\mathcal{O}_L} x[k] + \underbrace{\begin{bmatrix} D & 0 & 0 & \cdots & 0 \\ CB & D & 0 & \cdots & 0 \\ CAB & CB & D & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ CA^{L-1}B & CA^{L-2}B & CA^{L-3}B & \cdots & D \end{bmatrix}}_{M_L} \underbrace{\begin{bmatrix} u[k] \\ u[k + 1] \\ u[k + 2] \\ \vdots \\ u[k + L] \end{bmatrix}}_{U_{[k,L]}}$$

# Massey-Sain Algorithm (2)

$$M_L = \left[ \begin{array}{c|cccc} D & 0 & 0 & \cdots & 0 \\ CB & D & 0 & \cdots & 0 \\ CAB & CB & D & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ CA^{L-1}B & CA^{L-2}B & CA^{L-3}B & \cdots & D \end{array} \right]$$

- Notice that

$$\text{rank}(M_L) \leq \text{rank}(M_{L-1}) + m$$

- **Theorem:** System  $\mathcal{S}$  has an inverse with delay  $L$  if and only if

$$\text{rank}(M_L) - \text{rank}(M_{L-1}) = m$$



# Massey-Sain Algorithm (3)

$$M_L = \left[ \begin{array}{c|ccc} D & 0 & 0 & \cdots & 0 \\ CB & D & 0 & \cdots & 0 \\ CAB & CB & D & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ CA^{L-1}B & CA^{L-2}B & CA^{L-3}B & \cdots & D \end{array} \right]$$

## Proof of Theorem:

- If  $\text{rank}(M_L) - \text{rank}(M_{L-1}) < m$ , there exists a  $U_{[k,L]}$  in the null-space of  $M_L$ , with  $u[k] \neq 0$ , which is indistinguishable from the all-zero input
- If  $\text{rank}(M_L) - \text{rank}(M_{L-1}) = m$ , the first  $m$  columns are linearly independent of each other, and of the rest of the columns in  $M_L$ .
  - There exists a matrix  $\mathcal{K}$  such that

$$\begin{aligned} \mathcal{K}M_L &= \left[ \mathbf{I}_m \mid \mathbf{0} \right] \\ \Rightarrow \mathcal{K}Y_{[k,L]} &= \mathcal{K}O_L x[k] + u[k] \end{aligned}$$

# Massey-Sain Algorithm (4)

Construction of inverse:

- Input is given by

$$u[k] = -\mathcal{K}\mathcal{O}_L x[k] + \mathcal{K}Y_{[k,L]} \quad (1)$$

- Substitute into state-transition equation

$$\begin{aligned} x[k+1] &= Ax[k] + Bu[k] \\ &= Ax[k] - B\mathcal{K}\mathcal{O}_L x[k] + B\mathcal{K}Y_{[k,L]} \\ &= (A - B\mathcal{K}\mathcal{O}_L) x[k] + B\mathcal{K}Y_{[k,L]} \end{aligned} \quad (2)$$

- Equations (1) and (2) together form the state-space model of the inverse

# Example of Massey-Sain Algorithm (1)

$$x[k+1] = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} x[k] + \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} u[k]$$
$$y[k] = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} x[k] + \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} u[k]$$

- Test for invertibility:

$$M_0 = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}, \quad \text{rank}(M_0) = 1$$

$$M_1 = \left[ \begin{array}{cc|cc} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{array} \right], \quad \text{rank}(M_1) = 3$$

- Since  $\text{rank}(M_1) - \text{rank}(M_0) = 2$ , system is invertible with delay 1

## Example of Massey-Sain Algorithm (2)

- Find  $\mathcal{K}$  such that  $\mathcal{K}M_1 = \left[ \mathbf{I}_2 \mid \mathbf{0}_2 \right]$

$$\Rightarrow \mathcal{K} = \begin{bmatrix} 0 & 0 & 1 & -1 \\ 0 & 1 & -1 & 1 \end{bmatrix}$$

- Inverse is given by

$$\begin{aligned} x[k+1] &= \left( A - BK \begin{bmatrix} C \\ CA \end{bmatrix} \right) x[k] + BK \begin{bmatrix} y[k] \\ y[k+1] \end{bmatrix} \\ &= \begin{bmatrix} 1 & 0 \\ 1 & 0 \end{bmatrix} x[k] + \begin{bmatrix} 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} y[k] \\ y[k+1] \end{bmatrix} \end{aligned}$$

$$\begin{aligned} u[k] &= -\mathcal{K} \begin{bmatrix} C \\ CA \end{bmatrix} x[k] + \mathcal{K} \begin{bmatrix} y[k] \\ y[k+1] \end{bmatrix} \\ &= \begin{bmatrix} 1 & -1 \\ -1 & 0 \end{bmatrix} x[k] + \begin{bmatrix} 0 & 0 & 1 & -1 \\ 0 & 1 & -1 & 1 \end{bmatrix} \begin{bmatrix} y[k] \\ y[k+1] \end{bmatrix} \end{aligned}$$

# Massey-Sain Algorithm: Upper Bound

- When can we stop iterating?
- **Theorem:** If system is not invertible for  $L = n$ , it is not invertible at all
- Willsky later tightened upper bound to  $L = n - q + 1$ , where  $q$  is the nullity of  $D$

# Moylan's Algorithm (1)

- Appeared in IEEE Transactions on Automatic Control, vol. 22, 1977
- Define the matrix

$$M(\lambda) = \begin{bmatrix} A - \lambda I & B \\ C & D \end{bmatrix}$$

- **Theorem:**  $\mathcal{S}$  is invertible if and only if  $\text{rank}(M(\lambda)) = n + m$  for some real  $\lambda$
- **Proof of Necessity:**
  - Assume  $\text{rank}(M(\lambda)) < n + m$  for all  $\lambda$
  - Then for any  $\lambda_i$ , there exist  $x_i$  and  $u_i$  such that

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix} \begin{bmatrix} x_i \\ u_i \end{bmatrix} = \begin{bmatrix} \lambda_i x_i \\ 0 \end{bmatrix}$$

- It is possible to find scalars  $\alpha_i$  such that  $\sum_{i=0}^n \alpha_i x_i = 0$
- The input  $u[k] = \sum_{i=0}^n \alpha_i \lambda_i^k u_i$  results in  $y[k] = 0$  for all  $k$
- Thus, the system is not invertible

# Moylan's Algorithm (2)

Proof of sufficiency follows from the following construction:

- Consider the more general system

$$x[k + 1] = Ax[k] + Bu[k] + v[k]$$

$$y[k] = Cx[k] + Du[k]$$

- Suppose  $D$  has rank  $r < p$ 
  - There exists a non-singular  $p \times p$  matrix  $Q_1$  such that

$$Q_1 D = \begin{bmatrix} D_0 \\ 0 \end{bmatrix}, \quad Q_1 C = \begin{bmatrix} C_1 \\ C_2 \end{bmatrix}$$

- There exists a non-singular  $(p - r) \times (p - r)$  matrix  $Q_2$  such that

$$Q_2 C_2 = \begin{bmatrix} \tilde{C}_2 \\ 0 \end{bmatrix}$$

# Moylan's Algorithm (3)

- Let

$$\tilde{y}[k] = \underbrace{\begin{bmatrix} I & 0 \\ 0 & Q_2 \end{bmatrix}}_Q Q_1 y[k]$$

- This gives

$$\begin{bmatrix} \tilde{y}_1[k] \\ \tilde{y}_2[k] \\ \tilde{y}_3[k] \end{bmatrix} = \begin{bmatrix} C_1 \\ \tilde{C}_2 \\ 0 \end{bmatrix} x[k] + \begin{bmatrix} D_0 \\ 0 \\ 0 \end{bmatrix} u[k]$$

- Define similarity transformation  $\tilde{x}[k] = T x[k]$  such that

$$\tilde{C}_2 T^{-1} = \begin{bmatrix} 0 & I \end{bmatrix}$$



# Moylan's Algorithm (4)

- New system:

$$\begin{bmatrix} \tilde{x}_1[k+1] \\ \tilde{x}_2[k+1] \end{bmatrix} = \begin{bmatrix} \tilde{A}_{11} & \tilde{A}_{12} \\ \tilde{A}_{21} & \tilde{A}_{22} \end{bmatrix} \begin{bmatrix} \tilde{x}_1[k] \\ \tilde{x}_2[k] \end{bmatrix} + \begin{bmatrix} \tilde{B}_1 \\ \tilde{B}_2 \end{bmatrix} u[k] + \begin{bmatrix} \tilde{v}_1[k] \\ \tilde{v}_2[k] \end{bmatrix}$$
$$\begin{bmatrix} \tilde{y}_1[k] \\ \tilde{y}_2[k] \\ \tilde{y}_3[k] \end{bmatrix} = \begin{bmatrix} \tilde{C}_{11} & \tilde{C}_{12} \\ 0 & I \\ 0 & 0 \end{bmatrix} \begin{bmatrix} \tilde{x}_1[k] \\ \tilde{x}_2[k] \end{bmatrix} + \begin{bmatrix} D_0 \\ 0 \\ 0 \end{bmatrix} u[k]$$

- Note that  $\tilde{y}_2[k] = \tilde{x}_2[k]$
- Define

$$z_1[k] = \tilde{y}_1[k] - \tilde{C}_{12}\tilde{y}_2[k]$$

$$z_2[k] = \tilde{y}_2[k+1] - \tilde{A}_{22}\tilde{y}_2[k] - \tilde{v}_2[k]$$

$$w[k] = \tilde{v}_1[k] + \tilde{A}_{12}\tilde{y}_2[k]$$

$$q[k] = \tilde{x}_1[k]$$

# Moylan's Algorithm (5)

- System can be written as

$$q[k + 1] = \hat{A}q[k] + \hat{B}u[k] + w[k]$$

$$z[k] = \hat{C}q[k] + \hat{D}u[k]$$

where

$$\hat{A} = \begin{bmatrix} \tilde{A}_{11} \end{bmatrix}, \quad \hat{B} = \begin{bmatrix} \tilde{B}_1 \end{bmatrix}$$

$$\hat{C} = \begin{bmatrix} \tilde{C}_{11} \\ \tilde{A}_{12} \end{bmatrix}, \quad \hat{D} = \begin{bmatrix} D_0 \\ \tilde{B}_2 \end{bmatrix}$$

- This system has  $\hat{p} \leq p$  outputs and  $\hat{n} \leq n$  states
- If  $\text{rank}(\hat{D}) = m$ , inverse is given by  $\hat{D}^\dagger z[k] = \hat{D}^\dagger \hat{C}q[k] + u[k]$
- If  $\text{rank}(\hat{D}) < m$ , repeat procedure on new system
  - What if  $\hat{p} < m$ ?

# Moylan's Algorithm (6)

- Define

$$\begin{aligned}\hat{M}(\lambda) &= \begin{bmatrix} \hat{A} - \lambda I & \hat{B} \\ \hat{C} & \hat{D} \end{bmatrix} \\ &= \begin{bmatrix} \tilde{A}_{11} - \lambda I & \tilde{B}_1 \\ \tilde{C}_{11} & D_0 \\ \tilde{A}_{12} & \tilde{B}_2 \end{bmatrix}\end{aligned}$$

- If  $\text{rank}(M(\lambda)) = n + m$  for  $\lambda = \lambda_i$ , then

$$\text{rank}(\hat{M}(\lambda_i)) = \hat{n} + m$$

- This implies  $\hat{p} \geq m$ !

# Moylan's Algorithm (7)

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The following can be proved in a similar manner:

- System  $\mathcal{S}$  has a stable inverse if and only if  $\text{rank}(M(\lambda)) = n + m$  for all  $|\lambda| > 1$
- System  $\mathcal{S}$  is invertible with unknown initial state if and only if  $\text{rank}(M(\lambda)) = n + m$  for all  $\lambda$

# Partial Invertibility (1)

What if we only want to invert some of the inputs?

- Suppose  $u[k] = \begin{bmatrix} u_1[k] \\ u_2[k] \end{bmatrix}$ , with  $u_1[k] \in \mathbb{R}^{m_1}$ ,  $u_2[k] \in \mathbb{R}^{m_2}$ 
  - Partition  $B$  and  $D$  as  $B = \begin{bmatrix} B_1 & B_2 \end{bmatrix}$ ,  $D = \begin{bmatrix} D_1 & D_2 \end{bmatrix}$
- System  $\mathcal{S}$  becomes

$$x[k+1] = Ax[k] + B_1u_1[k] + B_2u_2[k]$$

$$y[k] = Cx[k] + D_1u_1[k] + D_2u_2[k]$$

- What are the conditions on the system such that  $u_1[k]$  is invertible?

# Partial Invertibility (2)

Suppose  $x[k]$  is unknown

- Can we invert  $u_1[k]$  based only on  $y[k], y[k + 1], \dots, y[k + L]$ , for some  $L$ ?
- The response of system  $\mathcal{S}$  over  $L + 1$  time units is given by

$$\begin{bmatrix} y[k] \\ y[k + 1] \\ \vdots \\ y[k + L] \end{bmatrix} = \begin{bmatrix} C \\ CA \\ \vdots \\ CA^L \end{bmatrix} x[k] + \begin{bmatrix} D_1 & 0 & \cdots & 0 \\ CB_1 & D_1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ CA^{L-1}B_1 & CA^{L-2}B_1 & \cdots & D_1 \end{bmatrix} \begin{bmatrix} u_1[k] \\ u_1[k + 1] \\ \vdots \\ u_1[k + L] \end{bmatrix} + \begin{bmatrix} D_2 & 0 & \cdots & 0 \\ CB_2 & D_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ CA^{L-1}B_2 & CA^{L-2}B_2 & \cdots & D_2 \end{bmatrix} \begin{bmatrix} u_2[k] \\ u_2[k + 1] \\ \vdots \\ u_2[k + L] \end{bmatrix}$$

# Partial Invertibility (3)

- Previous expression can be written more compactly as

$$\underbrace{\begin{bmatrix} y[k] \\ y[k+1] \\ y[k+2] \\ \vdots \\ y[k+L] \end{bmatrix}}_{Y_{[k,L]}} = \underbrace{\begin{bmatrix} D_1 & D_2 & 0 & \cdots & 0 & C \\ CB_1 & CB_2 & D & \cdots & 0 & CA \\ CAB_1 & CAB_2 & CB & \cdots & 0 & CA^2 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ CA^{L-1}B_1 & CA^{L-1}B_2 & CA^{L-2}B & \cdots & D & CA^L \end{bmatrix}}_{J_L} \underbrace{\begin{bmatrix} u_1[k] \\ u_2[k] \\ u[k+1] \\ \vdots \\ u[k+L] \\ x[k] \end{bmatrix}}_{I_{[k,L]}}$$

- Partition  $J_L$  as  $J_L = \left[ \begin{array}{c|c} \mathbf{\Gamma}_L & \mathbf{\Psi}_L \end{array} \right]$
- **Theorem:**  $u_1[k]$  is invertible with delay  $L$  and unknown state  $x[k]$  if and only if

$$\text{rank}(J_L) - \text{rank}(\mathbf{\Psi}_L) = m_1 \ .$$

# Partial Invertibility (4)

## Proof of Theorem:

- If  $\text{rank}(J_L) - \text{rank}(\Psi_L) < m_1$ , there exists a  $I_{[k,L]}$  in the null-space of  $J_L$ , with  $u_1[k] \neq 0$ , which is indistinguishable from the all-zero input
- If  $\text{rank}(J_L) - \text{rank}(\Psi_L) = m_1$ , the first  $m_1$  columns are linearly independent of each other, and of the rest of the columns in  $J_L$ .
  - There exists a matrix  $\mathcal{K}$  such that

$$\begin{aligned}\mathcal{K}J_L &= \left[ \mathbf{I}_{m_1} \mid \mathbf{0} \right] \\ \Rightarrow \mathcal{K}Y_{[k,L]} &= u_1[k]\end{aligned}$$



# Example: Fault Detection (1)

- Consider a model of the F-8 aircraft (Teneketzis et al.)

$$\dot{x}(t) = \underbrace{\begin{bmatrix} -0.01357 & -32.2 & -46.3 & 0 \\ 0.00012 & 0 & 1.214 & 0 \\ -0.0001212 & 0 & -1.214 & 1 \\ 0.00057 & 0 & -9.01 & -0.6696 \end{bmatrix}}_A x(t) + \underbrace{\begin{bmatrix} -0.433 \\ 0.1394 \\ -0.1394 \\ -0.1577 \end{bmatrix}}_{B_u} u(t) + \underbrace{\begin{bmatrix} -46.3 \\ 1.214 \\ -1.214 \\ -9.01 \end{bmatrix}}_{B_d} d(t)$$
$$y(t) = \underbrace{\begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}}_C x(t)$$

where  $u(t)$  is the elevator deflection in radians, and  $d(t)$  is the wind disturbance

- We wish to detect failures in the elevator
  - Invert  $u(t)$  and compare to the specified value

## Example: Fault Detection (2)

- Perform test for inversion:

$$J_0 = \left[ \begin{array}{c|cc} D_u & D_d & C \end{array} \right], \quad \text{rank}(J_0) - \text{rank}(\Psi_0) = 0$$

$$J_1 = \left[ \begin{array}{c|cccc} D_u & D_d & 0 & 0 & C \\ CB_u & CB_d & D_u & D_d & CA \end{array} \right], \quad \text{rank}(J_1) - \text{rank}(\Psi_1) = 1$$

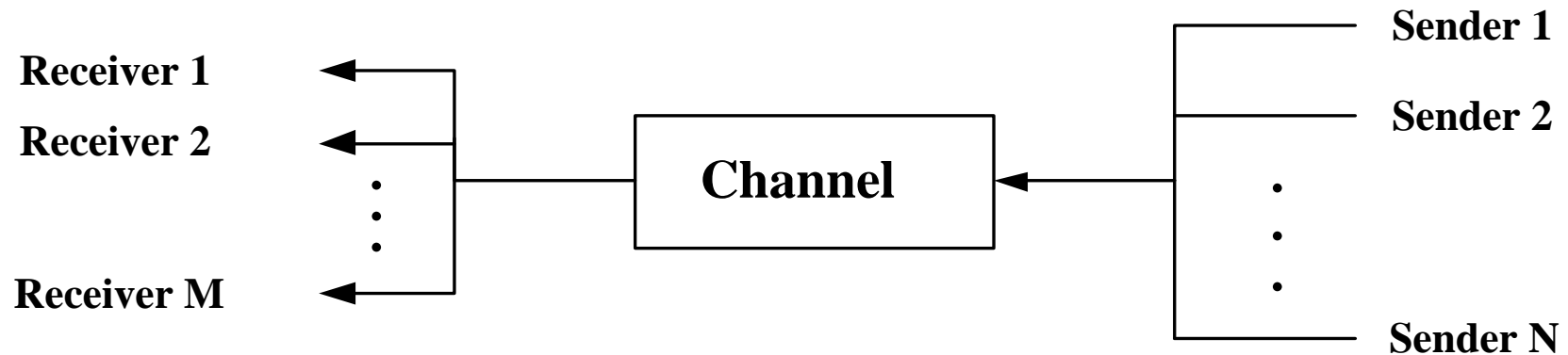
- Find  $\mathcal{K}$  such that  $\mathcal{K}J_1 = \left[ \begin{array}{c|cccc} 1 & 0 & 0 & 0 & 0 \end{array} \right]$

- Input is given by

$$u(t) = \underbrace{\begin{bmatrix} -0.0018 & -6.5936 & -0.2048 & -7.8097 \end{bmatrix}}_{\mathcal{K}} \begin{bmatrix} y(t) \\ \dot{y}(t) \end{bmatrix}$$

- Thus  $u(t)$  is invertible with “delay” 1

# Example: Multiplexing in Communication Channels



- Suppose multiple users broadcast through a dynamic channel
  - Is it possible for Receiver  $i$  to only decode messages from a certain subset of senders?

# Future Work

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- Partial inversion:
  - Finish proof of upper bound on  $L$  for partial inversion with unknown state
  - Study construction of partial inverses for systems with known initial conditions
  - Develop partial inverses of minimum dimension
- Investigate invertibility of hybrid/switched systems
- Design system inverses that are robust to parameter variations